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HYPERBOLIC POLYNOMIALS

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# Combinatorial and algorithmic aspects of hyperbolic polynomials

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## Abstract

Let  $p(x_1, \dots, x_n) = \sum_{(r_1, \dots, r_n) \in I_{n,n}} a_{(r_1, \dots, r_n)} \prod_{1 \leq i \leq n} x_i^{r_i}$  be homogeneous polynomial of degree  $n$  in  $n$  real variables with integer nonnegative coefficients. The support of such polynomial  $p(x_1, \dots, x_n)$  is defined as  $supp(p) = \{(r_1, \dots, r_n) \in I_{n,n} : a_{(r_1, \dots, r_n)} \neq 0\}$ . The convex hull  $CO(supp(p))$  of  $supp(p)$  is called the Newton polytope of  $p$ . We study the following decision problems , which are far-reaching generalizations of the classical perfect matching problem , :

- **Problem 1** . Consider a homogeneous polynomial  $p(x_1, \dots, x_n)$  of degree  $n$  in  $n$  real variables with nonnegative integer coefficients given as a black box (oracle ) . *Is it true that  $(1, 1, \dots, 1) \in supp(p)$  ?*
- **Problem 2** . Consider a homogeneous polynomial  $p(x_1, \dots, x_n)$  of degree  $n$  in  $n$  real variables with nonnegative integer coefficients given as a black box (oracle ) . *Is it true that  $(1, 1, \dots, 1) \in CO(supp(p))$  ?*

We prove that for hyperbolic polynomials these two problems are equivalent and can be solved by deterministic polynomial-time oracle algorithms . This result is based on a "hyperbolic" generalization of Rado theorem .

## 1 Introduction and motivating examples

Let  $p(x_1, \dots, x_n) = \sum_{(r_1, \dots, r_n) \in I_{n,n}} a_{(r_1, \dots, r_n)} \prod_{1 \leq i \leq n} x_i^{r_i}$  be homogeneous polynomial of degree  $n$  in  $n$  real variables. Here  $I_{k,n}$  stands for the set of vectors  $r = (r_1, \dots, r_k)$  with nonnegative integer components and  $\sum_{1 \leq i \leq k} r_i = n$ . We mainly study in this paper homogeneous polynomials with nonnegative integer coefficients .

**Definition 1.1:** The support of polynomial  $p(x_1, \dots, x_n)$  as above is defined as  $supp(p) = \{(r_1, \dots, r_n) \in I_{n,n} : a_{(r_1, \dots, r_n)} \neq 0\}$  . The convex hull  $CO(supp(p))$  of  $supp(p)$  is called the Newton polytope of  $p$  . ■

We will study the following decision problems :

- **Problem 1** . Consider a homogeneous polynomial  $p(x_1, \dots, x_n)$  of degree  $n$  in  $n$  real variables with nonnegative integer coefficients given as a black box (oracle ) . *Is it true that  $(1, 1, \dots, 1) \in supp(p)$  ?*
- **Problem 2** . Consider a homogeneous polynomial  $p(x_1, \dots, x_n)$  of degree  $n$  in  $n$  real variables with nonnegative integer coefficients given as a black box (oracle ) . *Is it true that  $(1, 1, \dots, 1) \in CO(supp(p))$  ?*

Our goal is solve these decision problems using deterministic polynomial-time oracle algorithms , i.e algorithms which evaluate the given  $p(x_1, \dots, x_n)$  at polynomial in  $n$  and  $\log(p(1, 1, \dots, 1))$  number of rational vectors  $(q_1, \dots, q_n)$  ; these rational vectors  $(q_1, \dots, q_n)$  have polynomial in  $n$  and  $\log(p(1, 1, \dots, 1))$  bit-wise complexity ; and the additional auxilary arithmetic computations also take polynomial in  $n$  and  $\log(p(1, 1, \dots, 1))$  number of steps .

The next example explains some (well known ) origins of the both problems .

**Example 1.2:** Consider first a multilinear polynomial  $mul(x_1, \dots, x_n) = \prod_{1 \leq i \leq n} \sum_{1 \leq j \leq n} a_{i,j} x_j$  , where  $A = (a_{i,j} : 1 \leq i, j \leq n)$  is a square integer matrix . Then

$$p(x_1, \dots, x_n) = \sum_{r=(r_1, \dots, r_n) \in I_{n,n}} \prod_{1 \leq i \leq n} \alpha_i^{r_i} per(A_r) \frac{1}{\prod_{1 \leq i \leq n} r_i!}, \quad (1)$$

where  $A_r$  is a square integer matrix consisting of  $r_i$  copies of the  $i$ th column of  $A$  ,  $1 \leq i \leq n$  ; and  $per(A)$  is the usual permanent of  $A$  . Notice that in this multilinear case the polynomial  $p(.)$  can be evaluated in  $O(N^2)$  arithmetic operations and  $(1, 1, \dots, 1) \in supp(p)$  iff  $per(A) \neq 0$  . Therefore unless  $P = NP$  there is no hope to design deterministic polynomial oracle algorithm solving Problem 1 in this case . Next consider even more general case of determinantal polynomials :

$$q(x_1, \dots, x_n) = \det\left(\sum_{1 \leq i \leq n} A_i x_i\right),$$

where  $\mathbf{A} = (A_1, \dots, A_n)$  is a  $n$ -tuple of positive semidefinite  $n \times n$  hermitian matrices , i.e.  $A_i \succeq 0$  , with integer entries . Recall that the mixed discriminant

$$D(\mathbf{A}) = \frac{\partial^n}{\partial \alpha_1 \dots \partial \alpha_n} \det\left(\sum_{1 \leq i \leq n} A_i x_i\right).$$

It is well known (see , for instance , [19] ) that a determinantal polynomial  $q()$  can be represented as

$$q(x_1, \dots, x_n) = \sum_{r \in I_{n,n}} \prod_{1 \leq i \leq n} x_i^{r_i} D(\mathbf{A}_r), \frac{1}{\prod_{1 \leq i \leq n} r_i!} \quad (2)$$

where a  $n$ -tuple of square matrices consists of  $r_i$  copies of  $A_i$ ,  $1 \leq i \leq k$  . One of the equivalent formulations [28] of the classical Rado theorem states that  $D(\mathbf{A}_r) > 0$  iff

$$Rank\left(\sum_{i \in S} A_i\right) \geq |S| \text{ for all } S \subset \{1, 2, \dots, n\} \quad (3)$$

One important corollary of the Rado conditions (3) is that

$$supp(q) = CO(supp(q)) \cap I_{n,n}. \quad (4)$$

I.e. if integer vectors  $r, r(1), r(2), \dots, r(k) \in I(n, n)$  and

$$r = \sum_{1 \leq i \leq k} a(i)r(i), a(i) \geq 0, 1 \leq i \leq k; \sum_{1 \leq i \leq k} a(i),$$

and  $D(\mathbf{A}_{r(i)}) > 0, 1 \leq i \leq k$  then also  $D(\mathbf{A}_r) > 0$ . Notice that in this case Problem 1 and Problem 2 are equivalent.

We can rewrite Rado conditions (3) as follows :

$$\max_{r \in \text{supp}(q)} \sum_{i \in S} r_i \geq |S| \text{ for all } S \subset \{1, 2, \dots, n\} \quad (5)$$

Putting things together we get the following Fact .

**Fact 1.3:** The following properties of determinantal polynomial  $q((x_1, \dots, x_n)) = \det(\sum_{1 \leq i \leq n} A_i x_i)$  with  $n \times n$  hermitian matrices  $A_i \succeq 0, 1 \leq i \leq n$  are equivalent .

1.  $(1, 1, \dots, 1) \notin \text{supp}(q)$ .
2.  $(1, 1, \dots, 1) \notin CO(\text{supp}(q))$ .
3. There exists nonempty  $S \subset \{1, 2, \dots, n\}$  such that

$$\sum_{1 \leq i \leq n} r_i s_i < \sum_{1 \leq i \leq n} s_i = |S| \text{ for all } (r_1, \dots, r_n) \in \text{supp}(q), \quad (6)$$

, where  $(s_1, \dots, s_n)$  is a characteristic function of the subset  $S$  , i.e.  $s_i = 1$  if  $i \in S$  , and  $s_i = 0$  otherwise .

Notice that if (6) holds then the distance  $\text{dist}(e, CO(\text{supp}(q)))$  from the vector  $e = (1, \dots, 1)$  to the Newton polytope  $CO(\text{supp}(q))$  is at least  $\sqrt{\frac{n}{|S|(n-|S|)}} \geq \frac{2}{\sqrt{n}}$  .

■

We will show that for any class of polynomials satisfying Fact () there exists a deterministic polynomial-time oracle algorithm solving both Problem 1 and Problem 2 , which are , of course , equivalent in this case . Our algorithm is based on the reduction to some convex programming problem and the consequent use of the Ellipsoids method .

The next fact about determinantal polynomials , namely their hyperbolicity , is happened to be the most important .

**Fact 1.4:** Consider a determinantal polynomial  $q((x_1, \dots, x_n)) = \det(\sum_{1 \leq i \leq n} A_i x_i)$  with  $A_i \succeq 0, 1 \leq i \leq n$  . Assume that  $q$  is not identically zero , i.e. that  $B =: \sum_{1 \leq i \leq n} A_i \succ 0$  (the sum is strictly positive definite ). For a real vector  $(x_1, \dots, x_n) \in R^n$  consider the following polynomial equation of degree  $n$  in one variable :

$$P(t) = q(x_1 - t, x_2 - t, \dots, x_n - t) = \det(\sum_{1 \leq i \leq n} A_i x_i - t \sum_{1 \leq i \leq n} A_i) = 0. \quad (7)$$

The equation (7) has  $n$  real roots counting the multiplicities ; if the real vector  $(x_1, \dots, x_n) \in R^n$  has nonnegative entries then all roots of (7) are nonnegative real numbers . ■

**Proof:** First , the matrix  $A =: \sum_{1 \leq i \leq n} A_i x_i$  is hermitian . Second ,  $\det(A - tB) = 0$  iff  $\det(B^{-\frac{1}{2}}AB^{-\frac{1}{2}} - tI) = 0$  , where  $B^{-\frac{1}{2}}$  is the unique positive definite operator square root of positive definite matrix  $B^{-1}$  . As , clearly ,  $B^{-\frac{1}{2}}AB^{-\frac{1}{2}}$  is also hermitian hence its eigenvalues , which are the roots of () , are real . If  $x_i \geq 0, 1 \leq i \leq n$  , then the matrix  $B^{-\frac{1}{2}}AB^{-\frac{1}{2}} \succeq 0$  . Therefore in this case the roots of (7) are nonnegative real numbers . ■

The main result of this paper that this hyperbolicity , which we will describe formally in Section 1.1 , is sufficient for Fact 1.3 ; i.e. Fact 1.4 implies Fact 1.4 . ■

## 1.1 Hyperbolic polynomials

The following concept of hyperbolic polynomials was originated in the theory of partial differential equations [15] .

A homogeneous polynomial  $p(x), x \in R^m$  of degree  $n$  in  $m$  real variables is called hyperbolic in the direction  $e \in R^m$  (or  $e$ - hyperbolic) if for any  $x \in R^m$  the polynomial  $p(x - \lambda e)$  in the one variable  $\lambda$  has exactly  $n$  real roots counting their multiplicities. We assume in this paper that  $p(e) > 0$  . Denote an ordered vector of roots of  $p(x - \lambda e)$  as  $\lambda(x) = (\lambda_1(x) \geq \lambda_2(x) \geq \dots \lambda_n(x))$ . It is well known that the product of roots is equal to  $p(x)$ . Call  $x \in R^m$   $e$ -positive ( $e$ -nonnegative) if  $\lambda_n(x) > 0$  ( $\lambda_n(x) \geq 0$ ). The fundamental result [15] in the theory of hyperbolic polynomials states that the set of  $e$ -nonnegative vectors is a closed convex cone. A  $k$ -tuple of vectors  $(x_1, \dots, x_k)$  is called  $e$ -positive ( $e$ -nonnegative) if  $x_i, 1 \leq i \leq k$  are  $e$ -positive ( $e$ -nonnegative).

We denote the closed convex cone of  $e$ -nonnegative vectors as  $N_e(p)$ , and the open convex cone of  $e$ -positive vectors as  $C_e(p)$ . It has been shown in [15] (see also [21]) that an  $e$ - hyperbolic polynomial  $p$  is also  $d$ - hyperbolic for all  $e$ -positive vectors  $d \in C_e(p)$ .

Let us fix  $n$  real vectors  $x_i \in R^m, 1 \leq i \leq n$  and define the following homogeneous polynomial:

$$P_{x_1, \dots, x_n}(\alpha_1, \dots, \alpha_n) = p\left(\sum_{1 \leq i \leq n} \alpha_i x_i\right) \quad (8)$$

Following [21] , we define the  $p$ -mixed value of an  $n$ -vector tuple  $\mathbf{X} = (x_1, \dots, x_n)$  as

$$M_p(\mathbf{X}) =: M_p(x_1, \dots, x_n) = \frac{\partial^n}{\partial \alpha_1 \dots \partial \alpha_n} p\left(\sum_{1 \leq i \leq n} \alpha_i x_i\right) \quad (9)$$

Equivalently, the  $p$ -mixed value  $M_p(x_1, \dots, x_n)$  can be defined by the polarization (see [21]) :

$$M_p(x_1, \dots, x_n) = 2^{-n} \sum_{b_i \in \{-1, +1\}, 1 \leq i \leq n} p\left(\sum_{1 \leq i \leq n} b_i x_i\right) \prod_{1 \leq i \leq n} b_i \quad (10)$$

Associate with any vector  $r = (r_1, \dots, r_n) \in I_{n,n}$  an  $n$ -tuple of  $m$ -dimensional vectors  $\mathbf{X}_r$  consisting of  $r_i$  copies of  $x_i (1 \leq i \leq n)$ . It follows, for instance from the polarization identity (10), that

$$P_{x_1, \dots, x_n}(\alpha_1, \dots, \alpha_n) = \sum_{r \in I_{n,n}} \prod_{1 \leq i \leq n} \alpha_i^{r_i} M_p(\mathbf{X}_r) \frac{1}{\prod_{1 \leq i \leq n} r_i!} \quad (11)$$

For  $e$ -nonnegative tuple  $\mathbf{X} = (x_1, \dots, x_n)$ , define its capacity as:

$$Cap(\mathbf{X}) = \inf_{\alpha_i > 0, \prod_{1 \leq i \leq n} \alpha_i = 1} P_{x_1, \dots, x_n}(\alpha_1, \dots, \alpha_n) \quad (12)$$

Probably the best known example of a hyperbolic polynomial is

$$P(\alpha_0, \dots, \alpha_k) = \text{Det} \left( \sum_{0 \leq i \leq k} \alpha_i A_i \right) \quad (13)$$

where  $A_i, 0 \leq i \leq k$  are hermitian matrices and the linear space spanned by  $A_i, 0 \leq i \leq k$  contains a strictly positive definite matrix:  $\sum_{0 \leq i \leq k} \beta_i A_i = B > 0$ . This polynomial is hyperbolic in the direction  $\beta = (\beta_1, \dots, \beta_k)$ . We can assume wlog that  $B = I$  and that  $\beta = (1, 0, 0, \dots, 0)$ . In other words, after a nonsingular linear change of variables

$$P(\alpha_0, \dots, \alpha_k) = \text{Det} \left( \sum_{0 \leq i \leq k} \alpha_i B_i \right) \quad (14)$$

where the matrices  $B_i, 1 \leq i \leq k$  are hermitian and  $B_0 = I$ .

In this case mixed forms are just mixed discriminants.

We make a substantial use of the following very recent result [22] , which is a rather direct corollary of [1] , [30] and even much older [10] .

**Theorem 1.5:** Consider a homogeneous polynomial  $p(y_1, y_2, y_3))$  of degree  $n$  in 3 real variables which is hyperbolic in the direction  $(0, 0, 1)$ . Assume that  $p(0, 0, 1) = 1$  . Then there exists two  $n \times n$  real symmetric matrices  $A, B$  such that

$$p(y_1, y_2, y_3)) = \det(y_1 A + y_2 B + y_3 I).$$

It has been shown in [16] that most of known facts (and some opened problems ) about hyperbolic polynomials follow from Theorem 1.5 .

## 2 A hyperbolic analogue of the Rado theorem

**Definition 2.1:** Consider a homogeneous polynomial  $p(x), x \in R^m$  of degree  $n$  in  $m$  real variables which is hyperbolic in the direction  $e$ .Denote an ordered vector of roots of  $p(x - \lambda e)$  as  $\lambda(x) = (\lambda_1(x) \geq \lambda_2(x) \geq \dots \lambda_n(x))$  . We define the  $p$ -rank of  $x \in R^m$  in direction  $e$  as  $\text{Rank}_p(x) = |\{i : \lambda_i(x) \neq 0\}|$ . It follows from Theorem 1.5 that the  $p$ -rank of  $x \in R^m$  in any direction  $d \in C_e$  is equal to the  $p$ -rank of  $x \in R^m$  in direction  $e$  , which we call the  $p$ -rank of  $x \in R^m$  . ■

Consider the following polynomial in one variable  $D(t) = p(td + x) = \sum_{0 \leq i \leq n} c_i t^i$ . It follows from the identity (11) that

$$c_n = M_p(d, \dots, d)(n!)^{-1} = p(d), c_{n-1} = M_p(x, d, \dots, d)(1!(n-1)!)^{-1}, \dots, c_0 = M_p(x, \dots, x)(n!)^{-1} = p(x). \quad (15)$$

Let  $(\lambda_1^{(d)}(x) \geq \lambda_2^{(d)}(x) \geq \dots \geq \lambda_n^{(d)}(x))$  be the (real) roots of  $x$  in the  $e$ -positive direction  $d$ , i.e. the roots of the equation  $p(td - x) = 0$  . Define (canonical symmetric functions) :

$$S_{k,d}(x) = \sum_{1 \leq i_1 < i_2 < \dots < i_k \leq n} \lambda_{i_1}(x) \lambda_{i_2}(x) \dots \lambda_{i_k}(x).$$

Then  $S_{k,d}(x) = \frac{c_{n-k}}{c_n}$ . Clearly if  $x$  is  $e$ -nonnegative then for any  $e$ -positive  $d$  the  $p$ -rank  $Rank_p(x) = \max\{k : S_{k,d}(x) > 0\}$ . The next theorem, which we prove in Appendix A, is the main mathematical result of this paper.

**Theorem 2.2:** Consider a homogeneous polynomial  $p(x), x \in R^m$  of degree  $n$  in  $m$  real variables which is hyperbolic in the direction  $e$ . Let  $(\mathbf{X}) = (x_1, \dots, x_k), x_i \in R^m$  be  $e$ -nonnegative)  $n$ -tuple of  $m$ -dimensional vectors, i.e.  $x_i, 1 \leq i \leq k$  are  $e$ -nonnegative. Then the  $p$ -mixed form  $M_p(\mathbf{X}) =: M_p(x_1, \dots, x_n)$  is positive iff the following generalized Rado conditions hold :

$$Rank_p(\sum_{i \in S} x_i) \geq |S| \text{ for all } S \subset \{1, 2, \dots, n\}. \quad (16)$$

**Definition 2.3:** Call a homogeneous polynomial  $p(x), x \in R^n$  of degree  $n$  in  $n$  real variables  $P$ -hyperbolic if it is hyperbolic in direction  $e = (1, 1, \dots, 1)$  (vector of all ones) and all canonical orts  $e_i, 1 \leq i \leq n$  (rows of the identity matrix  $I$ ) are  $e$ -nonnegative. Call a homogeneous polynomial  $q(x), x \in R^n$  of degree  $n$  in  $n$  real variables with nonnegative coefficients  $S$ -hyperbolic if there exists a  $P$ -hyperbolic polynomial  $p$  such that  $supp(p) = supp(q)$ . ■

**Corollary 2.4:** Let  $q(x), x \in R^n$  be  $S$ -hyperbolic polynomial of degree  $n$ . Then  $CO(supp(q)) \cap I_{n,n}$ .

**Corollary 2.5:** Let  $q(x), x \in R^n$  be  $S$ -hyperbolic polynomial of degree  $n$ . Then the following conditions are equivalent

1.  $e \in CO(supp(q))$ .
2.  $e \in supp(q)$ , i.e.  $\frac{\partial^n}{\partial \alpha_1 \dots \partial \alpha_n} q(x) > 0$ .
3.  $Cap(p) =: \inf_{\alpha_i > 0, \prod_{1 \leq i \leq n} \alpha_i = 1} q(\alpha_1, \dots, \alpha_n) > 0$ .
4. For all  $\epsilon > 0$  there exists a vector  $(\alpha_1, \dots, \alpha_n)$  with positive entries such that the following inequality holds :

$$\sum_{1 \leq i \leq n} \left| \frac{\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n)}{q(\alpha_1, \dots, \alpha_n)} - 1 \right|^2 \leq \epsilon. \quad (17)$$

5. There exists a vector  $(\alpha_1, \dots, \alpha_n)$  with positive entries such that the following inequality holds :

$$\sum_{1 \leq i \leq n} \left| \frac{\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n)}{q(\alpha_1, \dots, \alpha_n)} - 1 \right|^2 \leq \frac{1}{n}. \quad (18)$$

6. For all subsets  $S \subset \{1, 2, \dots, n\}$  the following inequality holds :

$$\sum_{i \in S} r_i \geq |S| \text{ for all } (r_1, \dots, r_n) \in supp(q). \quad (19)$$

(We sketch a proof in Appendix C . )

The following result , which we prove in Appendix B , is a "polynomial" generalization of Lemma 4.2 in [17] .

**Proposition 2.6:** *The condition (18) implies the condition (19) for all homogeneous polynomial  $q(x)$ ,  $x \in R^n$  of degree  $n$  in  $n$  real variables with nonnegative coefficients .*

### 3 The ellipsoid algorithm

Consider a homogeneous polynomial  $q(x)$ ,  $x \in R^n$  of degree  $n$  in  $n$  real variables with nonnegative integer coefficients . Associate with such  $q$  the following convex functional

$$f(y_1, \dots, y_n) = \log(q(e^{y_1}, e^{y_2}, \dots, e^{y_n})).$$

**Proposition 3.1:** *The following conditions are equivalent*

1.  $e = (1, 1, \dots, 1) \in CO(supp(q))$  .
2.  $\inf_{y_1+\dots+y_n=0} f(y_1, \dots, y_n) \geq 0$ .

If  $e = (1, 1, \dots, 1) \notin CO(supp(q))$  then  $\inf_{y_1+\dots+y_n=0} f(y_1, \dots, y_n) = -\infty$ .  
Let  $dist(e, CO(supp(q))) = \Delta^{-1} > 0$  and  $Q = \log(q(e))$  . Define  $\gamma = (Q + 1)\Delta$  . Then

$$\inf_{y_1+\dots+y_n=0, (|y_1|^2+\dots+|y_n|^2)^{\frac{1}{2}} \leq \gamma} f(y_1, \dots, y_n) = \min_{y_1+\dots+y_n=0, |y_1|^2+\dots+|y_n|^2 \leq \gamma} f(y_1, \dots, y_n) \leq -1. \quad (20)$$

**Proof:** Our proof is a strighforward application of the concavity of the logarithm on the positive semi-axis and of Hanh-Banach separation theorem . It will be included in the full version . ■

Proposition 3.1 suggests the following natural approach to solve Problem 2 , i.e. to decide whether  $e = (1, 1, \dots, 1) \in CO(supp(q))$  or not :

find  $\min_{y_1+\dots+y_n=0, |y_1|^2+\dots+|y_n|^2 \leq \gamma} f(y_1, \dots, y_n)$  with absolute accuracy  $\frac{1}{3}$  . If the resulting value is greater or equal  $-\frac{1}{3}$  then  $e = (1, 1, \dots, 1) \in CO(supp(q))$  ; if the resulting value is less or equal  $-\frac{2}{3}$  then  $e = (1, 1, \dots, 1) \notin CO(supp(q))$  . And , of course , it is natural to use the ellipsoid method . Our main tool is the following property of the ellipsoid algorithm [26]: For a prescribed accuracy  $\delta > 0$ , it finds a  $\delta$ -minimizer of a differentiable convex function  $f$  in a ball  $B$ , that is a point  $x_\delta \in B$  with  $f(x_\delta) \leq \min_B f + \delta$ , in no more than

$$O\left(n^2 \ln\left(\frac{2\delta + \text{Var}_B(f)}{\delta}\right)\right), \quad \text{Var}_B(f) = \max_B f - \min_B f \quad (21)$$

iterations. Each iteration requires a single computation of the value and of the gradient of  $f$  at a given point, plus  $O(n^2)$  elementary operations to run the algorithm itself. In our case, this is

easily seen to cost at most  $O(n^2)$  oracle calls and  $O(n)$  elementary arithmetic operations . We have  $n - 1$  dimensional ball  $B_\gamma = \{(y_1, \dots, y_n) : y_1 + \dots + y_n = 0, |y_1|^2 + \dots + |y_n|^2 \leq \gamma\}$ . A straighforward computations show that

$$Var_B(f) \leq \log(q(1, 1, \dots, 1)e^{\gamma n}) - \log(q(1, 1, \dots, 1)e^{-\gamma n}) \leq 2\gamma n.$$

Which gives  $O(n^2(\ln(n) + \ln(\gamma)))$  iterations of the ellipsoid method needed to solve Problem 2 , it amounts to  $O(n^4(\ln(n) + \ln(\gamma)))$  oracle calls . And  $O(n^4(\ln(n) + \ln(\gamma)))$  is polynomial in  $n$  even if  $\gamma$  is exponentially large ( $dist(e, CO(supp(q)))$  is exponentially small ). The problem is that if  $\gamma$  is exponentially large ( and it can happened ) then we need to call oracles on inputs with exponential bit-size .

Putting things together , we get the following conclusion :

*If it is promised that either  $e = (1, 1, \dots, 1) \in CO(supp(q))$  or  $dist(e, CO(supp(q))) \geq poly(n)^{-1}$  for some fixed polynomial  $poly(n)$  then Problem 1 can be solved by a deterministic polynomial-time oracle algorithm based on the ellipsoid method .*

And at this point we can say nothing about Problem 1 , i.e. deciding whether  $e = (1, 1, \dots, 1) \in supp(q)$  or not . Corollary 2.5 says that if  $q$  is  $S$ -hyperbolic polynomial then Problem 1 and Problem 2 are equivalent ; moreover if  $e = (1, 1, \dots, 1) \notin supp(q)$  then here exists nonempty  $S \subset \{1, 2, \dots, n\}$  such that

$$\sum_{1 \leq i \leq n} r_i s_i < \sum_{1 \leq i \leq n} s_i = |S| \text{ for all } (r_1, \dots, r_n) \in supp(q), \quad (22)$$

, where  $(s_1, \dots, s_n)$  is a characteristic function of the subset  $S$  , i.e.  $s_i = 1$  if  $i \in S$  , and  $s_i = 0$  otherwise .

Notice that if (22) holds then the distance  $dist(e, CO(supp(q)))$  from the vector  $e = (1, \dots, 1)$  to the Newton polytope  $CO(supp(q))$  is at least  $\sqrt{\frac{n}{|S|(n-|S|)}} \geq \frac{2}{\sqrt{n}}$  . Thus we have the next theorem .

**Theorem 3.2:** *Problem 1 and Problem 2 are equivalent for  $S$ -hyperbolic polynomials . There exists a deterministic polynomial-time oracle algorithm solving Problem 1 for a given  $S$ -hyperbolic polynomial  $q(\alpha_1, \dots, \alpha_n)$  with integer coefficients . It requires  $O(n^4(\ln(n) + \ln(\ln(q(1, 1, \dots, 1))))$  oracle calls and it bit-wise complexity (which roughly the radius of the ball  $B_\gamma$  ) is  $O(n^{\frac{1}{2}} \ln(q(1, 1, \dots, 1)))$*

## 4 Hyperbolic Sinkhorn scaling

We will discuss briefly in this section another method , which is essentially a large step version of the gradient descent .

**Definition 4.1:** Consider an  $e$ -nonnegative tuple  $\mathbf{X} = (x_1, \dots, x_n)$  such that the sum of its components  $S(\mathbf{X}) = d = \sum_{1 \leq i \leq k} x_i$  is  $e$ -positive. Define  $tr_d(x)$  as a sum of roots of the univariate polynomial equation  $p(x - td) = 0$ .

Define the following map (Hyperbolic Sinkhorn Scaling) acting on such tuples:

$$HS(\mathbf{X}) = \mathbf{Y} = \left( \frac{x_1}{tr_d(x_1)}, \dots, \frac{x_n}{tr_d(x_n)} \right)$$

Hyperbolic Sinkhorn Iteration (**HSI**) is a recursive procedure:

$$\mathbf{X}_{j+1} = HS(\mathbf{X}_j), j \geq 0, \mathbf{X}_0 \text{ is an } e\text{-nonnegative tuple with } \sum_{1 \leq i \leq k} x_i \in C_e.$$

We also define the doubly-stochastic defect of  $e$ -nonnegative tuples with  $e$ -positive sums as

$$DS(\mathbf{X}) = \sum_{1 \leq i \leq k} (tr_d(x_i) - 1)^2; \sum_{1 \leq i \leq k} x_i = d \in C_e$$

■

We can define the map  $HS(\cdot)$  directly in terms of the  $P$ -hyperbolic polynomial

$$Q(\alpha_1, \dots, \alpha_n) = P_{x_1, \dots, x_n}(\alpha_1, \dots, \alpha_n) = p\left(\sum_{1 \leq i \leq n} \alpha_i x_i\right).$$

Indeed, if  $\sum_{1 \leq i \leq n} \alpha_i x_i = d \in C_e$  then

$$tr_d(\alpha_i x_i) = \frac{\alpha_i \frac{\partial}{\partial \alpha_i} Q(\alpha_1, \dots, \alpha_n)}{Q(\alpha_1, \dots, \alpha_n)} \quad (23)$$

This gives the following way to redefine the map  $HS(\mathbf{X})$ :

$$HS(\alpha_1, \dots, \alpha_n) = \left( \frac{Q(\alpha_1, \dots, \alpha_n)}{\frac{\partial}{\partial \alpha_1} Q(\alpha_1, \dots, \alpha_n)}, \dots, \frac{Q(\alpha_1, \dots, \alpha_n)}{\frac{\partial}{\partial \alpha_n} Q(\alpha_1, \dots, \alpha_n)} \right); \alpha_i > 0, 1 \leq i \leq n.$$

And correspondingly the doubly-stochastic defect of  $(\alpha_1, \dots, \alpha_n)$  is equal to

$$\sum_{1 \leq i \leq n} \left| \frac{\alpha_i \frac{\partial}{\partial \alpha_i} Q(\alpha_1, \dots, \alpha_n)}{Q(\alpha_1, \dots, \alpha_n)} - 1 \right|^2,$$

the same as the left side of (18). Notice that  $\sum_{1 \leq i \leq n} tr_d(x_i) = n$  by the Euler's identity.

**Example 4.2:** Consider the following hyperbolic polynomial in  $n$  variables:  $p(z_1, \dots, z_n) = \prod_{1 \leq i \leq n} z_i$ . It is  $e$ -hyperbolic for  $e = (1, 1, \dots, 1)$ . And  $N_e$  is a nonnegative orthant,  $C_e$  is a positive orthant. An  $e$ -nonnegative tuple  $\mathbf{X} = (x_1, \dots, x_n)$  can be represented by an  $n \times n$  matrix  $A_{\mathbf{X}}$  with nonnegative entries: the  $i$ th column of  $A$  is a vector  $x_i \in R^n$ . If  $Z = (z_1, \dots, z_n) \in R^n$  and  $d = (d_1, \dots, d_n) \in R^n$ ;  $z_i > 0, 1 \leq i \leq n$ , then  $tr_d(Z) = \sum_{1 \leq i \leq n} \frac{z_i}{d_i}$ . Recall that for a square matrix  $A = \{a_{ij} : 1 \leq i, j \leq N\}$  row scaling is defined as

$$R(A) = \left\{ \frac{a_{ij}}{\sum_j a_{ij}} \right\},$$

column scaling as  $C(A) = \left\{ \frac{a_{ij}}{\sum_i a_{ij}} \right\}$  assuming that all denominators are nonzero. The iterative process  $\dots CRCA(A)$  is called *Sinkhorn's iterative scaling* (SI). In terms of the matrix  $A_{\mathbf{X}}$  the map  $HS(\mathbf{X})$  can be realized as follows:

$$A_{HS(\mathbf{X})} = C(R(A_{\mathbf{X}}))$$

So, the map  $HS(\mathbf{X})$  is indeed a (rather far-reaching) generalization of Sinkhorn's scaling. Other generalizations (not all hyperbolic) can be found in [20], [3], [2]. ■

The following result , proved in [16] , allows to use **(HSI)** to solve Problem 1 for  $P$ -hyperbolic polynomials  $q$  in the same way as it was done for the perfect matching problem in [20] , [17] ; and for the Edmonds' problem in [3] . The corresponding complexity is  $O(n \log(q(e)))$  iterations of **(HSI)** , which can be done in  $O(n^3 \log(q(e)))$  oracle calls . The algorithm works in the following way :

*Run  $K = O(n \log(q(e)))$  Hyperbolic Sinkhorn Iterations  $\mathbf{X}_{j+1} = HS(\mathbf{X}_j)$  ; if  $DS(\mathbf{X}_i) \leq \frac{1}{n}$  for some  $i \leq K$  then the  $p$ -mixed form  $M_p(\mathbf{X}_0) > 0$  , and  $M_p(\mathbf{X}_0) = 0$  otherwise .*

**Proposition 4.3:** *Let  $y_i = \frac{x_i}{tr_d(x_1)}$  , where  $x_i$  is  $e$ -nonnegative ,  $1 \leq i \leq n$  , and  $d = \sum_{1 \leq i \leq n} x_i$  is  $e$ - positive . Then (clearly)  $w = \sum_{1 \leq i \leq n} y_i$  is  $e$ - positive . Let positive real numbers  $\lambda_1 \geq \dots \geq \lambda_n$  be the roots of the equation  $p(w - td) = 0$ . Then  $\sum_{1 \leq i \leq n} \lambda_i = n$  and thus  $p(w) = p(d) \prod_{1 \leq i \leq n} \lambda_i \leq p(d)$  .*

*In terms of the corresponding  $P$ -hyperbolic polynomial  $Q$  , the following inequality holds :*

$$Q\left(\left(\frac{\partial}{\partial \alpha_1} Q(\alpha_1, \dots, \alpha_n)\right)^{-1}, \dots, \left(\frac{\partial}{\partial \alpha_n} Q(\alpha_1, \dots, \alpha_n)\right)^{-1}\right) \leq Q(\alpha_1, \dots, \alpha_n)^{-(n-1)}; \alpha_i > 0. \quad (24)$$

## 5 Conclusion and Acknowledgments

Univariate polynomials with real roots appear quite often in modern combinatorics , especially in the context of integer polytopes . We discovered in this paper rather unexpected and very likely far-reaching connections between hyperbolic polynomials and many classical combinatorial and algorithmic problems . There are still several open problems . The most interesting is a hyperbolic generalization of the van der Waerden conjecture for permanents of doubly stochastic matrices .

**Question 5.1:** Define the van der Waerden constant of a hyperbolic in direction  $e$  polynomial  $p(y_1, \dots, y_m)$  of degree  $n$  in  $m$  real variables as

$$VDW(p) = \inf \frac{M_p(x_1, \dots, x_n)}{Cap(x_1, \dots, x_n)}$$

where the infimum is taken over the set of tuples  $(x_1, \dots, x_n)$  of  $e$ -positive vectors. It is easy to see that  $VDW(p) \leq \frac{n!}{n^n}$ . Is  $VDW(p) = \frac{n!}{n^n}$  ? Is it positive ? ■

For a hyperbolic in direction  $(1, 1, \dots, 1)$  polynomial  $Mul(y_1, \dots, y_n) = y_1 y_2 \dots y_n$  this question is equivalent to the famous van der Waerden conjecture for permanents of doubly stochastic matrices , proved in [12] , [13] . For a hyperbolic in direction  $I$  polynomial  $\det(X)$  ,  $X$  is  $n \times n$  hermitian matrix , it is equivalent to Bapat's conjecture [5] (it was also hinted in [12] ) , proved by the author in [18] , [29] .

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## A Proof of the (main ) Theorem 2.2

Before proving Theorem 2.2 , we will recall some basic properties of  $p$ -mixed forms and prove a few auxillary results . The following fact was proved in [21]

**Fact A.1:** Consider a homogeneous polynomial  $p(x), x \in R^m$  of degree  $n$  in  $m$  real variables which is hyperbolic in the direction  $e$ . Then the following properties hold .

1. The  $p$ -mixed form  $M_p(x_1, \dots, x_n)$  is linear in each  $x_i, 1 \leq i \leq n$ .
2. If  $x_1, x_2, \dots, x_{n-1}$  are  $e$ -nonnegative then the linear functional  $l(x) = M_p(x_1, \dots, x_{n-1}, x)$  is nonnegative on the closed cone  $N_e$  of  $e$ -nonnegative vectors .
3. If the tuples  $(x_1, \dots, x_n), (y_1, \dots, y_n), (x_1 - y_1, \dots, x_n - y_n)$  are  $e$ -nonnegative then

$$0 \leq M_p(y_1, \dots, y_n) \leq M_p(x_1, \dots, x_n).$$

4. Fix  $e$ -positive vector  $d$  and consider the following homogeneous polynomial  $p_d(x), x \in R^m$  of degree  $n - 1$  in  $m$  real variables :  $p_d(x) =: M_p(x, x, \dots, x, d)$  . Then  $p_d$  is hyperbolic in any  $e$ -positive direction  $v \in C_e(p)$  . If  $g \in C_e(p)$  (  $e$ -positive respect to the polynomial  $p$  ) then also  $g \in C_v(p_d)$  for all  $v \in C_e(p)$  .

■

The next fact is well known .

**Fact A.2:** Consider a sequence of univariate polynomials of the same degree  $n$  :  $P_k(t) = \sum_{0 \leq i \leq n} a_{i,k} t^i$  . suppose that  $\lim_{k \rightarrow \infty} a_{i,k} = a_i, 0 \leq i \leq n$  and  $a_n \neq 0$  . Define  $P(t) = \sum_{0 \leq i \leq n} a_i t^i$  . Then roots of  $P_k$  converge to roots of  $P$  . In particular if roots of all polynomials  $P_k$  are real then also roots of  $P$  are real ; if roots of all polynomials  $P_k$  are real nonnegative numbers then also roots of  $P$  are real nonnegative numbers . ■

The following corollary of Theorem 1.5 plays crucial role in our proof of Theorem 2.2 .

**Corollary A.3:**

1. Consider a homogeneous polynomial  $p(x), x \in R^m$  of degree  $n$  in  $m$  real variables which is hyperbolic in the direction  $e$ . Let  $x_1, x_2, x_3$  be three  $e$ -nonnegative vectors and  $d = x_1 + x_2 + x_3$  is  $e$ -positive . Assume wlog that  $p(x_1 + x_2 + x_3) = 1$  . Then there exists three symmetric positive semidefinite matrices  $A, B, C$  such that  $p(a_1 x_1 + a_2 x_2 + a_3 x_3) = \det(a_1 A + a_2 B + a_3 C)$  for all real  $a_1, a_2, a_3$ . Additionally , the roots of  $a_1 x_1 + a_2 x_2 + a_3 x_3$  in the direction  $d$  , i.e. the roots of the equation  $p(a_1 x_1 + a_2 x_2 + a_3 x_3 - td) = 0$  , coincide with the eigenvalues of  $a_1 A + a_2 B + a_3 C$  .
2. Theorem 2.2 is true for  $e$ -nonnegative tuples  $(\mathbf{X}) = (x_1, \dots, x_n), x_i \in R^m$  consisting of at most three distinct components , i.e the cardinality of the set  $\{x_1, \dots, x_n\}$  is at most three .

**Proof:**

1. Consider the following homogeneous polynomial  $L(b_1, b_2, b_3) = P(b_1x_1 + b_2x_2 + b_3(x_1 + x_2 + x_3))$  of degree  $n$  in 3 real variables. It follows from Theorem 1.5 that there exists two real symmetric matrices  $A$  and  $B$  such that  $L(b_1, b_2, b_3) = \det(b_1A + b_2B + b_3I)$ . It follows that they both positive semidefinite, and  $C = I - A - B$  is also positive semidefinite. Take a real linear combination  $z = a_1x_1 + a_2x_2 + a_3x_3$ . Then

$$p(z - t(x_1 + x_2 + x_3)) = \det((a_1 - a_3)A + (a_2 - a_3)B + a_3I - tI) = \det(a_1A + a_2B + a_3C - tI).$$

This proves that  $p(a_1x_1 + a_2x_2 + a_3x_3) = \det(a_1A + a_2B + a_3C)$  for all real  $a_1, a_2, a_3$  by putting  $t = 0$ . And it also proves the "eigenvalues" statement.

2. Consider  $e$ -nonnegative tuple  $(\mathbf{X})$  consisting of  $r_i$  copies of  $x_i$ ,  $1 \leq i \leq 3$ ;  $r_1 + r_2 + r_3 = n$ . Assume that  $d = x_1 + x_2 + x_3$  is  $e$ -positive (if it is not then  $M_p(\mathbf{X}) = 0$  by a simple argument based on the monotonicity of  $p$ -mixed forms). It follows from the polarization formula (10), that

$$M_p(\mathbf{X}) = \sum_{1 \leq i \leq k < \infty} d_i p(t_{1,i}x_1 + t_{2,i}x_2 + t_{3,i}x_3),$$

and this formula is universal, i.e. holds for all homogeneous polynomial of degree  $n$ , in particular for  $\det(X)$ ,  $X$  is  $n \times n$  symmetric matrix. Therefore, using the first part of this Corollary we get that the  $p$ -mixed form  $M_p(\mathbf{X}) = D(\mathbf{A})$ , where the matrix tuple  $\mathbf{A}$  consists of  $r_1$  copies of  $A$ ,  $r_2$  copies of  $B$  and  $r_3$  copies of  $C$  and  $D(\mathbf{A})$  is the mixed discriminant. Using Rado theorem for mixed discriminants we get that  $D(\mathbf{A}) > 0$  iff

$$\text{Rank}(\sum_{i \in S} A_i) \geq \sum_{i \in S} r_i \text{ for all } S \subset \{1, 2, 3\}.$$

But from the first part we get that  $\text{Rank}(\sum_{i \in S} A_i)$  is equal to  $p$ -rank  $\text{Rank}_p(\sum_{i \in S} x_i)$  of  $\sum_{i \in S} x_i$  for all  $S \subset \{1, 2, 3\}$ .

**Proposition A.4:** Consider similarly to part 4 of Fact A.1 the polynomial  $p_d(x) =: M_p(x, x, \dots, x, d)$  where  $d$  is  $e$ -nonnegative and  $\text{Rank}_p(d) \geq 1$ . Then  $p_d$  is hyperbolic in any direction  $z \in N_e(p)$  which is  $e$ -nonnegative and satisfies the following inequalities:

$$\text{Rank}_p(z) \geq n - 1; \text{ Rank}_p(z + d) = n. \quad (25)$$

Also, if  $y \in N_e(p)$  is  $e$ -nonnegative then also  $y \in N_z(p_d)$ , i.e. is  $z$ -nonnegative respect to the polynomial  $p_d$ .

**Proof:** Let  $z \in N_e(p)$  be  $e$ -nonnegative satisfying (25). Consider univariate polynomial  $P(t) = M_p(tz + x, tz + x, \dots, tz + x, d)$ . Then  $P(t) = \sum_{0 \leq i \leq n-1} a_i t^i$  and  $a_{n-1} = M_p(z, z, \dots, z, d)$ . It follows from Corollary A.3 that  $a_{n-1} > 0$ . Consider now a sequence of univariate polynomials  $P_k(t) = M_p(tz_k + x, tz_k + x, \dots, tz_k + x, d_k)$ . Where  $z_k, d_k$  are  $e$ -positive and  $\lim_{k \rightarrow \infty} z_k = z$ ,  $\lim_{k \rightarrow \infty} d_k = d$ . Then the coefficients of polynomials  $P_k$  converge to the coefficients of the

polynomial  $P$ . It follows from part 4 of Fact A.1 that the roots of  $P_k$  are real. Since  $a_{n-1} > 0$  hence using Fact A.2 we get that the roots of  $P$  are also real. This exactly means that the polynomial  $p_d$  is hyperbolic in direction  $z$ . The  $d$ -nonnegativity statement follows from the nonnegativity part of Fact A.2. ■

We are ready now to present our proof of Theorem 2.2.

**Proof:** (Proof of Theorem 2.2). The "only if" part is simple. Indeed supposed that there exists a subset  $S \subset \{1, 2, \dots, n\}$  such that  $\text{Rank}_p(\sum_{i \in S} x_i) < |S|$ , i.e. using the identities (15)  $M_p(k, k, \dots, k, d, \dots, d) = 0$ , where  $k = \sum_{i \in S} x_i$ ,  $d \in C_e(p)$  is  $e$ -positive and the  $n$ -tuple  $(k, k, \dots, k, d, \dots, d)$  consists of  $|S|$  copies of  $k = \sum_{i \in S} x_i$ . Let  $d$  be any  $e$ -positive positive vector such that  $d - x_i$  is  $e$ -nonnegative,  $1 \leq i \leq n$ . Using the monotonicity of  $p$ -mixed forms we get that

$$M_p(x_1, \dots, x_n) \leq M_p(k, k, \dots, k, d, \dots, d) = 0.$$

Our proof of the "if" part is by induction in the degree  $n$ . Suppose that the generalized Rado conditions (16) hold. Then at least  $\text{Rank}_p(x_n) \geq 1$ . Consider the following homogeneous polynomial of degree  $n-1$ :

$$p_d(x) = M_p(x, x, \dots, x, d), \quad d = x_n.$$

We get from Proposition a.4 the following assertion:

The polynomial  $p_d(x)$  is hyperbolic in direction  $z = \sum_{1 \leq i \leq n-1} x_i$  and the vectors  $x_i \in N_z(p_d)$ ,  $1 \leq i \leq n-1$ , i.e. are  $z$ -nonnegative respect to the polynomial  $p_d$ .

Indeed, it follows from the generalized Rado conditions (16) that  $\text{Rank}_p(z) \geq n-1$  and  $\text{Rank}_p(z + d) = \text{Rank}_p(\sum_{1 \leq i \leq n} x_i) = n$ .

Next we show that the  $n-1$ -tuple  $\mathbf{Y} = (x_1, \dots, x_{n-1})$  satisfies the generalized Rado conditions

$$\text{Rank}_{p_d}(\sum_{i \in S} x_i) \geq |S| \quad \text{for all } S \subset \{1, 2, \dots, n-1\}.$$

Or equivalently, that

$$M_p(k, \dots, k, z, \dots, z, d) > 0; \quad k = \sum_{i \in S} x_i, \quad z = \sum_{1 \leq i \leq n-1} x_i, \quad d = x_n, \quad S \subset \{1, \dots, n-1\}, \quad (26)$$

where the  $n$ -tuple  $\mathbf{T} = (k, \dots, k, z, \dots, z, d)$  consists of  $|S|$  copies of  $k$ ,  $n-1-|S|$  copies of  $z$  and one copy of  $d$ .

It is easy to see that the generalized Rado conditions for the  $n$ -tuple  $\mathbf{T}$  are implied by the generalized Rado conditions for the original  $n$ -tuple  $\mathbf{X} = (x_1, \dots, x_{n-1}, x_n)$ . Since the  $n$ -tuple  $(k, \dots, k, z, \dots, z, d)$  consists of at most three distinct components hence we can apply part 2 of Corollary A.3. Therefore we get that indeed

$$\text{Rank}_{p_d}(\sum_{i \in S} x_i) \geq |S| \quad \text{for all } S \subset \{1, 2, \dots, n-1\}. \quad (27)$$

Thus, by induction in the degree, we get that  $p_d$ -mixed form  $M_{p_d}(x_1, \dots, x_{n-1}) > 0$ : the polynomial  $p_d$  of degree  $n-1$  in  $m$  real variables is  $z$ -hyperbolic. But

$$M_{p_d}(x_1, \dots, x_{n-1}) = \frac{\partial^{n-1}}{\partial \alpha_1 \dots \partial \alpha_{n-1}} p_d(\sum_{1 \leq i \leq n-1} \alpha_i x_i) = \\ = \frac{\partial^{n-1}}{\partial \alpha_1 \dots \partial \alpha_{n-1}} M_p(\sum_{1 \leq i \leq n-1} \alpha_i x_i, \dots, \sum_{1 \leq i \leq n-1} \alpha_i x_i, x_n) = (n-1)! M_p(x_1, \dots, x_n).$$

We conclude that if Theorem 2.2 is true for  $n-1$  then it is also true for  $n$ , and the case " $n=1$ " is trivially true. ■

## B Proof of Proposition 2.6

**Proof:** Assume wlog that  $q(\alpha_1, \dots, \alpha_n) = 1$ . It follows from the Euler's identity that

$$\sum_{1 \leq i \leq n} \alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) = n.$$

Let  $q(\alpha_1, \dots, \alpha_n) = \sum_{(r_1, \dots, r_n) \in \text{supp}(q)} a_{(r_1, \dots, r_n)} \prod_{1 \leq i \leq n} \alpha_i^{r_i}$ .

Define (positive numbers)  $b_{(r_1, \dots, r_n)} = a_{(r_1, \dots, r_n)} \prod_{1 \leq i \leq n} \alpha_i^{r_i}$ ,  $(r_1, \dots, r_n) \in \text{supp}(q)$ .

Then  $\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) = \sum_{(r_1, \dots, r_n) \in \text{supp}(q)} r_i b_{(r_1, \dots, r_n)}$ .

Suppose that for some subset  $S \subset \{1, 2, \dots, n\}$ ,  $1 \leq |S| < n$  we have the inequality  $\sum_{i \in S} r_i < |S|$  for all  $(r_1, \dots, r_n) \in \text{supp}(q)$ . Then  $\sum_{i \in S} \alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) \leq |S| - 1$ . But the condition (18) says that  $\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) = 1 + \delta_i$  and  $\sum_{1 \leq i \leq n} |\delta_i|^2 \leq \frac{1}{n}$ . By the Cauchy-Schwarz inequality,  $\sum_{i \in S} |\delta_i| \leq \sqrt{\frac{|S|}{n}} < 1$ . Therefore,

$$\sum_{i \in S} \alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) \geq |S| - \sum_{i \in S} |\delta_i| > |S| - 1.$$

The last inequality gives a contradiction. ■

## C A sketch of a proof of Corollary 2.4

**Proof:** By Theorem 2.2 the conditions (1) and (2) are equivalent. (2) implies (3) for any homogeneous polynomial with nonnegative coefficients.

Let  $\alpha_i = e^{y_i}$ ,  $1 \leq i \leq n$ ;  $\sum_{1 \leq i \leq n} y_i = 0$ . Consider the following convex functional

$$f(y_1, \dots, y_n) = \log(q(e^{y_1}, e^{y_2}, \dots, e^{y_n})).$$

Here  $q(x)$ ,  $x \in R^n$  is a homogeneous polynomial of degree  $n$  in  $n$  real variables with nonnegative coefficients. Then

$$\frac{\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n)}{q(\alpha_1, \dots, \alpha_n)} = \frac{\partial}{\partial y_i} f(y_1, \dots, y_n), \quad 1 \leq i \leq n.$$

Notice the condition (3) is equivalent to the following condition:

$$\inf_{y_1 + \dots + y_n = 0} f(y_1, \dots, y_n) = L > -\infty.$$

Consider the anti-gradient flow , i.e. the system of differential equations

$$y_i(t)' = -\left(\frac{\partial}{\partial y_i} f(y_1, \dots, y_n) - 1\right), y_i(0) = 0; 1 \leq i \leq n.$$

It is well known that in this convex case the gradient flow is defined for all  $t \geq 0$  . Using the Euler's identity we get that

$$\frac{d}{dt} f(y_1(t), \dots, y_n(t)) = -\beta(t) =: -\sum_{1 \leq i \leq n} \left| \frac{\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n)}{q(\alpha_1, \dots, \alpha_n)} - 1 \right|^2$$

It is easy to see that , because of the convexity of  $f$  , a nonnegative function  $\beta(t)$  is non-increasing on  $[0, \infty)$  .

As  $\inf_{y_1+\dots+y_n=0} f(y_1, \dots, y_n) = L > -\infty$  thus  $\int_0^\infty \beta(t) dt < \infty$  . Thus  $\lim_{t \rightarrow \infty} \beta(t) = 0$  . This proves the implication (3)  $\rightarrow$  (4) for all homogeneous polynomials of degree  $n$  in  $n$  real variables with nonnegative coefficients .

The implication (4)  $\rightarrow$  (5) is obvious . The implication (5)  $\rightarrow$  (6) for general homogeneous polynomials of degree  $n$  in  $n$  real variables with nonnegative coefficients is Proposition 2.6 .

Finally , the condition (6) is equivalent to (2) . ■