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HYPERBOLIC POLYNOMIALS

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Combinatorial and algorithmic aspects of hyperbolic polynomials

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Abstract

Let $p(x_1, \dots, x_n) = \sum_{(r_1, \dots, r_n) \in I_{n,n}} a_{(r_1, \dots, r_n)} \prod_{1 \leq i \leq n} x_i^{r_i}$ be homogeneous polynomial of degree n in n real variables with integer nonnegative coefficients. The support of such polynomial $p(x_1, \dots, x_n)$ is defined as $\text{supp}(p) = \{(r_1, \dots, r_n) \in I_{n,n} : a_{(r_1, \dots, r_n)} \neq 0\}$. The convex hull $CO(\text{supp}(p))$ of $\text{supp}(p)$ is called the Newton polytope of p . We study the following decision problems, which are far-reaching generalizations of the classical perfect matching problem, :

- **Problem 1** . Consider a homogeneous polynomial $p(x_1, \dots, x_n)$ of degree n in n real variables with nonnegative integer coefficients given as a black box (oracle) . *Is it true that $(1, 1, \dots, 1) \in \text{supp}(p)$?*
- **Problem 2** . Consider a homogeneous polynomial $p(x_1, \dots, x_n)$ of degree n in n real variables with nonnegative integer coefficients given as a black box (oracle) . *Is it true that $(1, 1, \dots, 1) \in CO(\text{supp}(p))$?*

We prove that for hyperbolic polynomials these two problems are equivalent and can be solved by deterministic polynomial-time oracle algorithms . This result is based on a "hyperbolic" generalization of Rado theorem .

1 Introduction and motivating examples

Let $p(x_1, \dots, x_n) = \sum_{(r_1, \dots, r_n) \in I_{n,n}} a_{(r_1, \dots, r_n)} \prod_{1 \leq i \leq n} x_i^{r_i}$ be homogeneous polynomial of degree n in n real variables. Here $I_{k,n}$ stands for the set of vectors $r = (r_1, \dots, r_k)$ with nonnegative integer components and $\sum_{1 \leq i \leq k} r_i = n$. We mainly study in this paper homogeneous polynomials with nonnegative integer coefficients .

Definition 1.1: The support of polynomial $p(x_1, \dots, x_n)$ as above is defined as $\text{supp}(p) = \{(r_1, \dots, r_n) \in I_{n,n} : a_{(r_1, \dots, r_n)} \neq 0\}$. The convex hull $CO(\text{supp}(p))$ of $\text{supp}(p)$ is called the Newton polytope of p . ■

We will study the following decision problems :

- **Problem 1** . Consider a homogeneous polynomial $p(x_1, \dots, x_n)$ of degree n in n real variables with nonnegative integer coefficients given as a black box (oracle) . *Is it true that $(1, 1, \dots, 1) \in \text{supp}(p)$?*
- **Problem 2** . Consider a homogeneous polynomial $p(x_1, \dots, x_n)$ of degree n in n real variables with nonnegative integer coefficients given as a black box (oracle) . *Is it true that $(1, 1, \dots, 1) \in CO(\text{supp}(p))$?*

Our goal is solve these decision problems using deterministic polynomial-time oracle algorithms , i.e algorithms which evaluate the given $p(x_1, \dots, x_n)$ at polynomial in n and $\log(p(1, 1, \dots, 1))$ number of rational vectors (q_1, \dots, q_n) ; these rational vectors (q_1, \dots, q_n) have polynomial in n and $\log(p(1, 1, \dots, 1))$ bit-wise complexity ; and the additional auxiliary arithmetic computations also take polynomial in n and $\log(p(1, 1, \dots, 1))$ number of steps .
The next example explains some (well known) origins of the both problems .

Example 1.2: Consider first a multilinear polynomial $mul(x_1, \dots, x_n) = \prod_{1 \leq i \leq n} \sum_{1 \leq j \leq n} a_{i,j} x_j$, where $A = (a_{i,j} : 1 \leq i, j \leq n)$ is a square integer matrix . Then

$$p(x_1, \dots, x_n) = \sum_{r=(r_1, \dots, r_n) \in I_{n,n}} \prod_{1 \leq i \leq n} \alpha_i^{r_i} per(A_r) \frac{1}{\prod_{1 \leq i \leq n} r_i!}, \quad (1)$$

where A_r is a square integer matrix consisting of r_i copies of the i th column of A , $1 \leq i \leq n$; and $per(A)$ is the usual permanent of A . Notice that in this multilinear case the polynomial $p(\cdot)$ can be evaluated in $O(N^2)$ arithmetic operations and $(1, 1, \dots, 1) \in supp(p)$ iff $per(A) \neq 0$. Therefore unless $P = NP$ there is no hope to design deterministic polynomial oracle algorithm solving Problem 1 in this case . Next consider even more general case of determinantal polynomials :

$$q(x_1, \dots, x_n) = \det\left(\sum_{1 \leq i \leq n} A_i x_i\right),$$

where $\mathbf{A} = (A_1, \dots, A_n)$ is a n -tuple of positive semidefinite $n \times n$ hermitian matrices , i.e. $A_i \succeq 0$, with integer entries . Recall that the mixed discriminant

$$D(\mathbf{A}) = \frac{\partial^n}{\partial \alpha_1 \dots \partial \alpha_n} \det\left(\sum_{1 \leq i \leq n} A_i x_i\right).$$

It is well known (see , for instance , [19]) that a determinantal polynomial $q(\cdot)$ can be represented as

$$q(x_1, \dots, x_n) = \sum_{r \in I_{n,n}} \prod_{1 \leq i \leq n} x_i^{r_i} D(\mathbf{A}_r), \frac{1}{\prod_{1 \leq i \leq n} r_i!} \quad (2)$$

where a n -tuple of square matrices consists of r_i copies of A_i , $1 \leq i \leq k$. One of the equivalent formulations [28] of the classical Rado theorem states that $D(\mathbf{A}_r) > 0$ iff

$$Rank\left(\sum_{i \in S} A_i\right) \geq |S| \text{ for all } S \subset \{1, 2, \dots, n\} \quad (3)$$

One important corollary of the Rado conditions (3) is that

$$supp(q) = CO(supp(q)) \cap I_{n,n}. \quad (4)$$

I.e. if integer vectors $r, r(1), r(2), \dots, r(k) \in I(n, n)$ and

$$r = \sum_{1 \leq i \leq k} a(i)r(i), a(i) \geq 0, 1 \leq i \leq k; \sum_{1 \leq i \leq k} a(i),$$

and $D(\mathbf{A}_{r(i)}) > 0, 1 \leq i \leq k$ then also $D(\mathbf{A}_r) > 0$. Notice that in this case Problem 1 and Problem 2 are equivalent.

We can rewrite Rado conditions (3) as follows :

$$\max_{r \in \text{supp}(q)} \sum_{i \in S} r_i \geq |S| \text{ for all } S \subset \{1, 2, \dots, n\} \quad (5)$$

Putting things together we get the following Fact .

Fact 1.3: The following properties of determinantal polynomial $q((x_1, \dots, x_n) = \det(\sum_{1 \leq i \leq n} A_i x_i)$ with $n \times n$ hermitian matrices $A_i \succeq 0, 1 \leq i \leq n$ are equivalent .

1. $(1, 1, \dots, 1) \notin \text{supp}(q)$.
2. $(1, 1, \dots, 1) \notin CO(\text{supp}(q))$.
3. There exists nonempty $S \subset \{1, 2, \dots, n\}$ such that

$$\sum_{1 \leq i \leq n} r_i s_i < \sum_{1 \leq i \leq n} s_i = |S| \text{ for all } (r_1, \dots, r_n) \in \text{supp}(q), \quad (6)$$

, where (s_1, \dots, s_n) is a characteristic function of the subset S , i.e. $s_i = 1$ if $i \in S$, and $s_i = 0$ otherwise .

Notice that if (6) holds then the distance $\text{dist}(e, CO(\text{supp}(q)))$ from the vector $e = (1, \dots, 1)$ to the Newton polytope $CO(\text{supp}(q))$ is at least $\sqrt{\frac{n}{|S|(n-|S|)}} \geq \frac{2}{\sqrt{n}}$.

■

We will show that for any class of polynomials satisfying Fact () there exists a deterministic polynomial-time oracle algorithm solving both Problem 1 and Problem 2, which are, of course, equivalent in this case. Our algorithm is based on the reduction to some convex programming problem and the consequent use of the Ellipsoids method.

The next fact about determinantal polynomials, namely their hyperbolicity, is happened to be the most important.

Fact 1.4: Consider a determinantal polynomial $q((x_1, \dots, x_n) = \det(\sum_{1 \leq i \leq n} A_i x_i)$ with $A_i \succeq 0, 1 \leq i \leq n$. Assume that q is not identically zero, i.e. that $B =: \sum_{1 \leq i \leq n} A_i \succ 0$ (the sum is strictly positive definite). For a real vector $(x_1, \dots, x_n) \in \mathbb{R}^n$ consider the following polynomial equation of degree n in one variable :

$$P(t) = q(x_1 - t, x_2 - t, \dots, x_n - t) = \det\left(\sum_{1 \leq i \leq n} A_i x_i - t \sum_{1 \leq i \leq n} A_i\right) = 0. \quad (7)$$

The equation (7) has n real roots counting the multiplicities; if the real vector $(x_1, \dots, x_n) \in \mathbb{R}^n$ has nonnegative entries then all roots of (7) are nonnegative real numbers. ■

Proof: First , the matrix $A =: \sum_{1 \leq i \leq n} A_i x_i$ is hermitian . Second , $\det(A - tB) = 0$ iff $\det(B^{-\frac{1}{2}}AB^{-\frac{1}{2}} - tI) = 0$, where $B^{-\frac{1}{2}}$ is the unique positive definite operator square root of positive definite matrix B^{-1} . As , clearly , $B^{-\frac{1}{2}}AB^{-\frac{1}{2}}$ is also hermitian hence its eigenvalues , which are the roots of () , are real . If $x_i \geq 0, 1 \leq i \leq n$, then the matrix $B^{-\frac{1}{2}}AB^{-\frac{1}{2}} \succeq 0$. Therefore in this case the roots of (7) are nonnegative real numbers . ■

The main result of this paper that this hyperbolicity , which we will describe formally in Section 1.1 , is sufficient for Fact 1.3 ; i.e. Fact 1.4 implies Fact 1.4 . ■

1.1 Hyperbolic polynomials

The following concept of hyperbolic polynomials was originated in the theory of partial differential equations [15] .

A homogeneous polynomial $p(x), x \in R^m$ of degree n in m real variables is called hyperbolic in the direction $e \in R^m$ (or e -hyperbolic) if for any $x \in R^m$ the polynomial $p(x - \lambda e)$ in the one variable λ has exactly n real roots counting their multiplicities. We assume in this paper that $p(e) > 0$. Denote an ordered vector of roots of $p(x - \lambda e)$ as $\lambda(x) = (\lambda_1(x) \geq \lambda_2(x) \geq \dots \lambda_n(x))$. It is well known that the product of roots is equal to $p(x)$. Call $x \in R^m$ e -positive (e -nonnegative) if $\lambda_n(x) > 0$ ($\lambda_n(x) \geq 0$). The fundamental result [15] in the theory of hyperbolic polynomials states that the set of e -nonnegative vectors is a closed convex cone. A k -tuple of vectors (x_1, \dots, x_k) is called e -positive (e -nonnegative) if $x_i, 1 \leq i \leq k$ are e -positive (e -nonnegative).

We denote the closed convex cone of e -nonnegative vectors as $N_e(p)$, and the open convex cone of e -positive vectors as $C_e(p)$. It has been shown in [15] (see also [21]) that an e -hyperbolic polynomial p is also d -hyperbolic for all e -positive vectors $d \in C_e(p)$.

Let us fix n real vectors $x_i \in R^m, 1 \leq i \leq n$ and define the following homogeneous polynomial:

$$P_{x_1, \dots, x_n}(\alpha_1, \dots, \alpha_n) = p\left(\sum_{1 \leq i \leq n} \alpha_i x_i\right) \quad (8)$$

Following [21] , we define the p -mixed value of an n -vector tuple $\mathbf{X} = (x_1, \dots, x_n)$ as

$$M_p(\mathbf{X}) =: M_p(x_1, \dots, x_n) = \frac{\partial^n}{\partial \alpha_1 \dots \partial \alpha_n} p\left(\sum_{1 \leq i \leq n} \alpha_i x_i\right) \quad (9)$$

Equivalently, the p -mixed value $M_p(x_1, \dots, x_n)$ can be defined by the polarization (see [21]) :

$$M_p(x_1, \dots, x_n) = 2^{-n} \sum_{b_i \in \{-1, +1\}, 1 \leq i \leq n} p\left(\sum_{1 \leq i \leq n} b_i x_i\right) \prod_{1 \leq i \leq n} b_i \quad (10)$$

Associate with any vector $r = (r_1, \dots, r_n) \in I_{n,n}$ an n -tuple of m -dimensional vectors \mathbf{X}_r consisting of r_i copies of $x_i (1 \leq i \leq n)$. It follows, for instance from the polarization identity (10), that

$$P_{x_1, \dots, x_n}(\alpha_1, \dots, \alpha_n) = \sum_{r \in I_{n,n}} \prod_{1 \leq i \leq n} \alpha_i^{r_i} M_p(\mathbf{X}_r) \frac{1}{\prod_{1 \leq i \leq n} r_i!} \quad (11)$$

For e -nonnegative tuple $\mathbf{X} = (x_1, \dots, x_n)$, define its capacity as:

$$Cap(\mathbf{X}) = \inf_{\alpha_i > 0, \prod_{1 \leq i \leq n} \alpha_i = 1} P_{x_1, \dots, x_n}(\alpha_1, \dots, \alpha_n) \quad (12)$$

Probably the best known example of a hyperbolic polynomial is

$$P(\alpha_0, \dots, \alpha_k) = \text{Det} \left(\sum_{0 \leq i \leq k} \alpha_i A_i \right) \quad (13)$$

where $A_i, 0 \leq i \leq k$ are hermitian matrices and the linear space spanned by $A_i, 0 \leq i \leq k$ contains a strictly positive definite matrix: $\sum_{0 \leq i \leq k} \beta_i A_i = B \succ 0$. This polynomial is hyperbolic in the direction $\beta = (\beta_1, \dots, \beta_k)$. We can assume wlog that $B = I$ and that $\beta = (1, 0, 0, \dots, 0)$. In other words, after a nonsingular linear change of variables

$$P(\alpha_0, \dots, \alpha_k) = \text{Det} \left(\sum_{0 \leq i \leq k} \alpha_i B_i \right) \quad (14)$$

where the matrices $B_i, 1 \leq i \leq k$ are hermitian and $B_0 = I$.

In this case mixed forms are just mixed discriminants.

We make a substantial use of the following very recent result [22], which is a rather direct corollary of [1], [30] and even much older [10].

Theorem 1.5: *Consider a homogeneous polynomial $p(y_1, y_2, y_3)$ of degree n in 3 real variables which is hyperbolic in the direction $(0, 0, 1)$. Assume that $p(0, 0, 1) = 1$. Then there exists two $n \times n$ real symmetric matrices A, B such that*

$$p(y_1, y_2, y_3) = \det(y_1 A + y_2 B + y_3 I).$$

It has been shown in [16] that most of known facts (and some opened problems) about hyperbolic polynomials follow from Theorem 1.5.

2 A hyperbolic analogue of the Rado theorem

Definition 2.1: Consider a homogeneous polynomial $p(x), x \in R^m$ of degree n in m real variables which is hyperbolic in the direction e . Denote an ordered vector of roots of $p(x - \lambda e)$ as $\lambda(x) = (\lambda_1(x) \geq \lambda_2(x) \geq \dots \geq \lambda_n(x))$. We define the p -rank of $x \in R^m$ in direction e as $\text{Rank}_p(x) = |\{i : \lambda_i(x) \neq 0\}|$. It follows from Theorem 1.5 that the p -rank of $x \in R^m$ in any direction $d \in C_e$ is equal to the p -rank of $x \in R^m$ in direction e , which we call the p -rank of $x \in R^m$. ■

Consider the following polynomial in one variable $D(t) = p(td + x) = \sum_{0 \leq i \leq n} c_i t^i$. It follows from the identity (11) that

$$c_n = M_p(d, \dots, d)(n!)^{-1} = p(d), c_{n-1} = M_p(x, d, \dots, d)(1!(n-1)!)^{-1}, \dots, c_0 = M_p(x, \dots, x)(n!)^{-1} = p(x). \quad (15)$$

Let $(\lambda_1^{(d)}(x) \geq \lambda_2^{(d)}(x) \geq \dots \geq \lambda_n^{(d)}(x))$ be the (real) roots of x in the e -positive direction d , i.e. the roots of the equation $p(td - x) = 0$. Define (canonical symmetric functions) :

$$S_{k,d}(x) = \sum_{1 \leq i_1 < i_2 < \dots < i_k \leq n} \lambda_{i_1}(x) \lambda_{i_2}(x) \dots \lambda_{i_k}(x).$$

Then $S_{k,d}(x) = \frac{c_{n-k}}{c_n}$. Clearly if x is e -nonnegative then for any e -positive d the p -rank $\text{Rank}_p(x) = \max\{k : S_{k,d}(x) > 0\}$. The next theorem, which we prove in Appendix A, is the main mathematical result of this paper.

Theorem 2.2: Consider a homogeneous polynomial $p(x), x \in R^m$ of degree n in m real variables which is hyperbolic in the direction e . Let $(\mathbf{X}) = (x_1, \dots, x_k), x_i \in R^m$ be e -nonnegative n -tuple of m -dimensional vectors, i.e. $x_i, 1 \leq i \leq k$ are e -nonnegative. Then the p -mixed form $M_p(\mathbf{X}) =: M_p(x_1, \dots, x_n)$ is positive iff the following generalized Rado conditions hold:

$$\text{Rank}_p\left(\sum_{i \in S} x_i\right) \geq |S| \text{ for all } S \subset \{1, 2, \dots, n\}. \quad (16)$$

Definition 2.3: Call a homogeneous polynomial $p(x), x \in R^n$ of degree n in n real variables P -hyperbolic if it is hyperbolic in direction $e = (1, 1, \dots, 1)$ (vector of all ones) and all canonical orts $e_i, 1 \leq i \leq n$ (rows of the identity matrix I) are e -nonnegative. Call a homogeneous polynomial $q(x), x \in R^n$ of degree n in n real variables with nonnegative coefficients S -hyperbolic if there exists a P -hyperbolic polynomial p such that $\text{supp}(p) = \text{supp}(q)$. ■

Corollary 2.4: Let $q(x), x \in R^n$ be S -hyperbolic polynomial of degree n . Then $\text{CO}(\text{supp}(q)) \cap I_{n,n}$.

Corollary 2.5: Let $q(x), x \in R^n$ be S -hyperbolic polynomial of degree n . Then the following conditions are equivalent

1. $e \in \text{CO}(\text{supp}(q))$.
2. $e \in \text{supp}(q)$, i.e. $\frac{\partial^n}{\partial \alpha_1 \dots \partial \alpha_n} q(x) > 0$.
3. $\text{Cap}(p) =: \inf_{\alpha_i > 0, \prod_{1 \leq i \leq n} \alpha_i = 1} q(\alpha_1, \dots, \alpha_n) > 0$.
4. For all $\epsilon > 0$ there exists a vector $(\alpha_1, \dots, \alpha_n)$ with positive entries such that the following inequality holds:

$$\sum_{1 \leq i \leq n} \left| \frac{\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n)}{q(\alpha_1, \dots, \alpha_n)} - 1 \right|^2 \leq \epsilon. \quad (17)$$

5. There exists a vector $(\alpha_1, \dots, \alpha_n)$ with positive entries such that the following inequality holds:

$$\sum_{1 \leq i \leq n} \left| \frac{\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n)}{q(\alpha_1, \dots, \alpha_n)} - 1 \right|^2 \leq \frac{1}{n}. \quad (18)$$

6. For all subsets $S \subset \{1, 2, \dots, n\}$ the following inequality holds:

$$\sum_{i \in S} r_i \geq |S| \text{ for all } (r_1, \dots, r_n) \in \text{supp}(q). \quad (19)$$

(We sketch a proof in Appendix C .)

The following result , which we prove in Appendix B , is a "polynomial" generalization of Lemma 4.2 in [17] .

Proposition 2.6: *The condition (18) implies the condition (19) for all homogeneous polynomial $q(x), x \in R^n$ of degree n in n real variables with nonnegative coefficients .*

3 The ellipsoid algorithm

Consider a homogeneous polynomial $q(x), x \in R^n$ of degree n in n real variables with nonnegative integer coefficients . Associate with such q the following convex functional

$$f(y_1, \dots, y_n) = \log(q(e^{y_1}, e^{y_2}, \dots, e^{y_n})).$$

Proposition 3.1: *The following conditions are equivalent*

1. $e = (1, 1, \dots, 1) \in CO(\text{supp}(q))$.
2. $\inf_{y_1 + \dots + y_n = 0} f(y_1, \dots, y_n) \geq 0$.

If $e = (1, 1, \dots, 1) \notin CO(\text{supp}(q))$ then $\inf_{y_1 + \dots + y_n = 0} f(y_1, \dots, y_n) = -\infty$. Let $\text{dist}(e, CO(\text{supp}(q))) = \Delta^{-1} > 0$ and $Q = \log(q(e))$. Define $\gamma = (Q + 1)\Delta$. Then

$$\inf_{y_1 + \dots + y_n = 0, (|y_1|^2 + \dots + |y_n|^2)^{\frac{1}{2}} \leq \gamma} f(y_1, \dots, y_n) = \min_{y_1 + \dots + y_n = 0, |y_1|^2 + \dots + |y_n|^2 \leq \gamma} f(y_1, \dots, y_n) \leq -1. \quad (20)$$

Proof: Our proof is a strigthforward application of the concavity of the logarithm on the positive semi-axis and of Hanh-Banach separation theorem . It will be included in the full version . ■

Proposition 3.1 suggests the following natural approach to solve Problem 2 , i.e. to decide whether $e = (1, 1, \dots, 1) \in CO(\text{supp}(q))$ or not :

find $\min_{y_1 + \dots + y_n = 0, |y_1|^2 + \dots + |y_n|^2 \leq \gamma} f(y_1, \dots, y_n)$ with absolute accuracy $\frac{1}{3}$. If the resulting value is greater or equal $-\frac{1}{3}$ then $e = (1, 1, \dots, 1) \in CO(\text{supp}(q))$; if the resulting value is less or equal $-\frac{2}{3}$ then $e = (1, 1, \dots, 1) \notin CO(\text{supp}(q))$. And , of course , it is natural to use the ellipsoid method . Our main tool is the following property of the ellipsoid algorithm [26]: For a prescribed accuracy $\delta > 0$, it finds a δ -minimizer of a differentiable convex function f in a ball B , that is a point $x_\delta \in B$ with $f(x_\delta) \leq \min_B f + \delta$, in no more than

$$O\left(n^2 \ln\left(\frac{2\delta + \text{Var}_B(f)}{\delta}\right)\right), \quad \text{Var}_B(f) = \max_B f - \min_B f \quad (21)$$

iterations. Each iteration requires a single computation of the value and of the gradient of f at a given point, plus $O(n^2)$ elementary operations to run the algorithm itself. In our case, this is

easily seen to cost at most $O(n^2)$ oracle calls and $O(n)$ elementary arithmetic operations . We have $n - 1$ dimensional ball $B_\gamma = \{(y_1, \dots, y_n) : y_1 + \dots + y_n = 0, |y_1|^2 + \dots + |y_n|^2 \leq \gamma\}$. A straightforward computations show that

$$\text{Var}_B(f) \leq \log(q(1, 1, \dots, 1)e^{\gamma n}) - \log(q(1, 1, \dots, 1)e^{-\gamma n}) \leq 2\gamma n.$$

Which gives $O(n^2(\ln(n) + \ln(\gamma)))$ iterations of the ellipsoid method needed to solve Problem 2 , it amounts to $O(n^4(\ln(n) + \ln(\gamma)))$ oracle calls . And $O(n^4(\ln(n) + \ln(\gamma)))$ is polynomial in n even if γ is exponentially large ($\text{dist}(e, CO(\text{supp}(q)))$ is exponentially small). The problem is that if γ is exponentially large (and it can happened) then we need to call oracles on inputs with exponential bit-size .

Putting things together , we get the following conclusion :

If it is promised that either $e = (1, 1, \dots, 1) \in CO(\text{supp}(q))$ or $\text{dist}(e, CO(\text{supp}(q))) \geq \text{poly}(n)^{-1}$ for some fixed polynomial $\text{poly}(n)$ then Problem 1 can be solved by a deterministic polynomial-time oracle algorithm based on the ellipsoid method .

And at this point we can say nothing about Problem 1 , i.e. deciding whether $e = (1, 1, \dots, 1) \in \text{supp}(q)$ or not . Corollary 2.5 says that if q is S -hyperbolic polynomial then Problem 1 and Problem 2 are equivalent ; moreover if $e = (1, 1, \dots, 1) \notin \text{supp}(q)$ then there exists nonempty $S \subset \{1, 2, \dots, n\}$ such that

$$\sum_{1 \leq i \leq n} r_i s_i < \sum_{1 \leq i \leq n} s_i = |S| \text{ for all } (r_1, \dots, r_n) \in \text{supp}(q), \quad (22)$$

, where (s_1, \dots, s_n) is a characteristic function of the subset S , i.e. $s_i = 1$ if $i \in S$, and $s_i = 0$ otherwise .

Notice that if (22) holds then the distance $\text{dist}(e, CO(\text{supp}(q)))$ from the vector $e = (1, \dots, 1)$ to the Newton polytope $CO(\text{supp}(q))$ is at least $\sqrt{\frac{n}{|S|(n-|S|)}} \geq \frac{2}{\sqrt{n}}$. Thus we have the next theorem .

Theorem 3.2: *Problem 1 and Problem 2 are equivalent for S -hyperbolic polynomials . There exists a deterministic polynomial-time oracle algorithm solving Problem 1 for a given S -hyperbolic polynomial $q(\alpha_1, \dots, \alpha_n)$ with integer coefficients . It requires $O(n^4(\ln(n) + \ln(\ln(q(1, 1, \dots, 1))))$ oracle calls and its bit-wise complexity (which roughly the radius of the ball B_γ) is $O(n^{\frac{1}{2}} \ln(q(1, 1, \dots, 1)))$*

4 Hyperbolic Sinkhorn scaling

We will discuss briefly in this section another method , which is essentially a large step version of the gradient descent .

Definition 4.1: Consider an e -nonnegative tuple $\mathbf{X} = (x_1, \dots, x_n)$ such that the sum of its components $S(\mathbf{X}) = d = \sum_{1 \leq i \leq k} x_i$ is e -positive. Define $\text{tr}_d(x)$ as a sum of roots of the univariate polynomial equation $p(x - td) = 0$.

Define the following map (Hyperbolic Sinkhorn Scaling) acting on such tuples:

$$HS(\mathbf{X}) = \mathbf{Y} = \left(\frac{x_1}{\text{tr}_d(x_1)}, \dots, \frac{x_n}{\text{tr}_d(x_n)} \right)$$

Hyperbolic Sinkhorn Iteration (**HSI**) is a recursive procedure:

$$\mathbf{X}_{j+1} = HS(\mathbf{X}_j), j \geq 0, \mathbf{X}_0 \text{ is an } e\text{-nonnegative tuple with } \sum_{1 \leq i \leq k} x_i \in C_e.$$

We also define the doubly-stochastic defect of e -nonnegative tuples with e -positive sums as

$$DS(\mathbf{X}) = \sum_{1 \leq i \leq k} (tr_d(x_i) - 1)^2; \sum_{1 \leq i \leq k} x_i = d \in C_e$$

■

We can define the map $HS(\cdot)$ directly in terms of the P -hyperbolic polynomial

$$Q(\alpha_1, \dots, \alpha_n) = P_{x_1, \dots, x_n}(\alpha_1, \dots, \alpha_n) = p\left(\sum_{1 \leq i \leq n} \alpha_i x_i\right).$$

Indeed, if $\sum_{1 \leq i \leq n} \alpha_i x_i = d \in C_e$ then

$$tr_d(\alpha_i x_i) = \frac{\alpha_i \frac{\partial}{\partial \alpha_i} Q(\alpha_1, \dots, \alpha_n)}{Q(\alpha_1, \dots, \alpha_n)} \quad (23)$$

This gives the following way to redefine the map $HS(\mathbf{X})$:

$$HS(\alpha_1, \dots, \alpha_n) = \left(\frac{Q(\alpha_1, \dots, \alpha_n)}{\frac{\partial}{\partial \alpha_1} Q(\alpha_1, \dots, \alpha_n)}, \dots, \frac{Q(\alpha_1, \dots, \alpha_n)}{\frac{\partial}{\partial \alpha_n} Q(\alpha_1, \dots, \alpha_n)} \right); \alpha_i > 0, 1 \leq i \leq n.$$

And correspondingly the doubly-stochastic defect of $(\alpha_1, \dots, \alpha_n)$ is equal to

$$\sum_{1 \leq i \leq n} \left| \frac{\alpha_i \frac{\partial}{\partial \alpha_i} Q(\alpha_1, \dots, \alpha_n)}{Q(\alpha_1, \dots, \alpha_n)} - 1 \right|^2,$$

the same as the left side of (18) . Notice that $\sum_{1 \leq i \leq n} tr_d(x_i) = n$ by the Euler's identity .

Example 4.2: Consider the following hyperbolic polynomial in n variables: $p(z_1, \dots, z_n) = \prod_{1 \leq i \leq n} z_i$. It is e -hyperbolic for $e = (1, 1, \dots, 1)$. And N_e is a nonnegative orthant, C_e is a positive orthant. An e -nonnegative tuple $\mathbf{X} = (x_1, \dots, x_n)$ can be represented by an $n \times n$ matrix $A_{\mathbf{X}}$ with nonnegative entries: the i th column of A is a vector $x_i \in R^n$. If $Z = (z_1, \dots, z_n) \in R^n$ and $d = (d_1, \dots, d_n) \in R^n; z_i > 0, 1 \leq i \leq n$, then $tr_d(Z) = \sum_{1 \leq i \leq n} \frac{z_i}{d_i}$. Recall that for a square matrix $A = \{a_{ij} : 1 \leq i, j \leq N\}$ row scaling is defined as

$$R(A) = \left\{ \frac{a_{ij}}{\sum_j a_{ij}} \right\},$$

column scaling as $C(A) = \left\{ \frac{a_{ij}}{\sum_i a_{ij}} \right\}$ assuming that all denominators are nonzero. The iterative process $\dots C R C R(A)$ is called *Sinkhorn's iterative scaling* (SI). In terms of the matrix $A_{\mathbf{X}}$ the map $HS(\mathbf{X})$ can be realized as follows:

$$A_{HS(\mathbf{X})} = C(R(A_{\mathbf{X}}))$$

So, the map $HS(\mathbf{X})$ is indeed a (rather far-reaching) generalization of Sinkhorn's scaling. Other generalizations (not all hyperbolic) can be found in [20], [3], [2]. ■

The following result , proved in [16] , allows to use **(HSI)** to solve Problem 1 for P -hyperbolic polynomials q in the same way as it was done for the perfect matching problem in [20] , [17] ; and for the Edmonds' problem in [3] . The corresponding complexity is $O(n \log(q(e)))$ iterations of **(HSI)** , which can be done in $O(n^3 \log(q(e)))$ oracle calls . The algorithm works in the following way :

Run $K = O(n \log(q(e)))$ Hyperbolic Sinkhorn Iterations $\mathbf{X}_{j+1} = HS(\mathbf{X}_j)$; if $DS(\mathbf{X}_i) \leq \frac{1}{n}$ for some $i \leq K$ then the p -mixed form $M_p(\mathbf{X}_0) > 0$, and $M_p(\mathbf{X}_0) = 0$ otherwise .

Proposition 4.3: Let $y_i = \frac{x_i}{\text{tr}_d(x_1)}$, where x_i is e -nonnegative , $1 \leq i \leq n$, and $d = \sum_{1 \leq i \leq n} x_i$ is e - positive . Then (clearly) $w = \sum_{1 \leq i \leq n} y_i$ is e - positive . Let positive real numbers $\lambda_1 \geq \dots \geq \lambda_n$ be the roots of the equation $p(w - td) = 0$. Then $\sum_{1 \leq i \leq n} \lambda_i = n$ and thus $p(w) = p(d) \prod_{1 \leq i \leq n} \lambda_i \leq p(d)$.

In terms of the corresponding P -hyperbolic polynomial Q , the following inequality holds :

$$Q((\frac{\partial}{\partial \alpha_1} Q(\alpha_1, \dots, \alpha_n))^{-1}, \dots, (\frac{\partial}{\partial \alpha_n} Q(\alpha_1, \dots, \alpha_n))^{-1}) \leq Q(\alpha_1, \dots, \alpha_n)^{-(n-1)}; \alpha_i > 0. \quad (24)$$

5 Conclusion and Acknowledgments

Univariate polynomials with real roots appear quite often in modern combinatorics , especially in the context of integer polytopes . We discovered in this paper rather unexpected and very likely far-reaching connections between hyperbolic polynomials and many classical combinatorial and algorithmic problems . There are still several open problems . The most interesting is a hyperbolic generalization of the van der Waerden conjecture for permanents of doubly stochastic matrices .

Question 5.1: Define the van der Waerden constant of a hyperbolic in direction e polynomial $p(y_1, \dots, y_m)$ of degree n in m real variables as

$$VDW(p) = \inf \frac{M_p(x_1, \dots, x_n)}{Cap(x_1, \dots, x_n)}$$

where the infimum is taken over the set of tuples (x_1, \dots, x_n) of e -positive vectors. It is easy to see that $VDW(p) \leq \frac{n!}{n^n}$. Is $VDW(p) = \frac{n!}{n^n}$? Is it positive ? ■

For a hyperbolic in direction $(1, 1, \dots, 1)$ polynomial $Mul(y_1, \dots, y_n) = y_1 y_2 \dots y_n$ this question is equivalent to the famous van der Waerden conjecture for permanents of doubly stochastic matrices , proved in [12] , [13] . For a hyperbolic in direction I polynomial $\det(X)$, X is $n \times n$ hermitian matrix , it is equivalent to Bapat's conjecture [5] (it was also hinted in [12]) , proved by the author in [18] , [29] .

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A Proof of the (main) Theorem 2.2

Before proving Theorem 2.2 , we will recall some basic properties of p -mixed forms and prove a few auxillary results . The following fact was proved in [21]

Fact A.1: Consider a homogeneous polynomial $p(x), x \in R^m$ of degree n in m real variables which is hyperbolic in the direction e . Then the following properties hold .

1. The p -mixed form $M_p(x_1, \dots, x_n)$ is linear in each $x_i, 1 \leq i \leq n$.
2. If x_1, x_2, \dots, x_{n-1} are e -nonnegative then the linear functional $l(x) = M_p(x_1, \dots, x_{n-1}, x)$ is nonnegative on the closed cone N_e of e -nonnegative vectors .
3. If the tuples $(x_1, \dots, x_n), (y_1, \dots, y_n), (x_1 - y_1, \dots, x_n - y_n)$ are e -nonnegative then

$$0 \leq M_p(y_1, \dots, y_n) \leq M_p(x_1, \dots, x_n).$$

4. Fix e -positive vector d and consider the following homogeneous polynomial $p_d(x), x \in R^m$ of degree $n - 1$ in m real variables : $p_d(x) =: M_p(x, x, \dots, x, d)$. Then p_d is hyperbolic in any e -positive direction $v \in C_e(p)$. If $g \in C_e(p)$ (e -positive respect to the polynomial p) then also $q \in C_v(p_d)$ for all $v \in C_e(p)$.

■

The next fact is well known .

Fact A.2: Consider a sequence of univariate polynomials of the same degree n : $P_k(t) = \sum_{0 \leq i \leq n} a_{i,k} t^i$. suppose that $\lim_{k \rightarrow \infty} a_{i,k} = a_i, 0 \leq i \leq n$ and $a_n \neq 0$.

Define $P(t) = \sum_{0 \leq i \leq n} a_i t^i$. Then roots of P_k converge to roots of P . In particular if roots of all polynomials P_k are real then also roots of P are real ; if roots of all polynomials P_k are real nonnegative numbers then also roots of P are real nonnegative numbers . ■

The following corollary of Theorem 1.5 plays crucial role in our proof of Theorem 2.2 .

Corollary A.3:

1. Consider a homogeneous polynomial $p(x), x \in R^m$ of degree n in m real variables which is hyperbolic in the direction e . Let x_1, x_2, x_3 be three e -nonnegative vectors and $d = x_1 + x_2 + x_3$ is e -positive . Assume wlog that $p(x_1 + x_2 + x_3) = 1$. Then there exists three symmetric positive semidefinite matrices A, B, C such that $p(a_1 x_1 + a_2 x_2 + a_3 x_3) = \det(a_1 A + a_2 B + a_3 C)$ for all real a_1, a_2, a_3 . Additionally , the roots of $a_1 x_1 + a_2 x_2 + a_3 x_3$ in the direction d , i.e. the roots of the equation $p(a_1 x_1 + a_2 x_2 + a_3 x_3 - t d) = 0$, coincide with the eigenvalues of $a_1 A + a_2 B + a_3 C$.
2. Theorem 2.2 is true for e -nonnegative tuples $(\mathbf{X}) = (x_1, \dots, x_n), x_i \in R^m$ consisting of at most three distinct components , i.e the cardinality of the set $\{x_1, \dots, x_n\}$ is at most three .

Proof:

1. Consider the following homogeneous polynomial $L(b_1, b_2, b_3) = P(b_1x_1 + b_2x_2 + b_3(x_1 + x_2 + x_3))$ of degree n in 3 real variables . It follows from Theorem 1.5 that there exists two real symmetric matrices A and B such that $L(b_1, b_2, b_3) = \det(b_1A + b_2B + b_3I)$. It follows that they both positive semidefinite , and $C = I - A - B$ is also positive semidefinite . Take a real linear combination $z = a_1x_1 + a_2x_2 + a_3x_3$. Then

$$p(z - t(x_1 + x_2 + x_3)) = \det((a_1 - a_3)A + (a_2 - a_3)B + a_3I - tI) = \det(a_1A + a_2B + a_3C - tI).$$

This proves that $p(a_1x_1 + a_2x_2 + a_3x_3) = \det(a_1A + a_2B + a_3C)$ for all real a_1, a_2, a_3 by putting $t = 0$. And it also proves the "eigenvalues " statement .

2. Consider e -nonnegative tuple (\mathbf{X}) consisting of r_i copies of x_i , $1 \leq i \leq 3$; $r_1 + r_2 + r_3 = n$. Assume that $d = x_1 + x_2 + x_3$ is e -positive (if it is not then $M_p(\mathbf{X}) = 0$ by a simple argument based on the monotonicity of p -mixed forms). It follows from the polarization formula (10) , that

$$M_p(\mathbf{X}) = \sum_{1 \leq i \leq k < \infty} d_i p(t_{1,i}x_1 + t_{2,i}x_2 + t_{3,i}x_3),$$

and this formula is universal , i.e. holds for all homogeneous polynomial of degree n , in particular for $\det(X)$, X is $n \times n$ symmetric matrix . Therefore , using the first part of this Corollary we get that the p -mixed form $M_p(\mathbf{X}) = D(\mathbf{A})$, where the matrix tuple \mathbf{A} consists of r_1 copies of A , r_2 copies of B and r_3 copies of C and $D(\mathbf{A})$ is the mixed discriminant . Using Rado theorem for mixed discriminants we get that $D(\mathbf{A}) > 0$ iff

$$\text{Rank}(\sum_{i \in S} A_i) \geq \sum_{i \in S} r_i \text{ for all } S \subset \{1, 2, 3\}.$$

But from the first part we get that $\text{Rank}(\sum_{i \in S} A_i)$ is equal to p -rank $\text{Rank}_p(\sum_{i \in S} x_i)$ of $\sum_{i \in S} x_i$ for all $S \subset \{1, 2, 3\}$.

■

Proposition A.4: Consider similarly to part 4 of Fact A.1 the polynomial $p_d(x) =: M_p(x, x, \dots, x, d)$ where d is e -nonnegative and $\text{Rank}_p(d) \geq 1$. Then p_d is hyperbolic in any direction $z \in N_e(p)$ which is e -nonnegative and satisfies the following inequalities :

$$\text{Rank}_p(z) \geq n - 1; \text{Rank}_p(z + d) = n. \quad (25)$$

Also , if $y \in N_e(p)$ is e -nonnegative then also $y \in N_z(p_d)$, i.e. is z -nonnegative respect to the polynomial p_d .

Proof: Let $z \in N_e(p)$ be e -nonnegative satisfying (25) . Consider univariate polynomial $P(t) = M_p(tz + x, tz + x, \dots, tz + x, d)$. Then $P(t) = \sum_{0 \leq i \leq n-1} a_i t^i$ and $a_{n-1} = M_p(z, z, \dots, z, d)$. It follows from Corollary A.3 that $a_{n-1} > 0$. Consider now a sequence of univariate polynomials $P_k(t) = M_p(tz_k + x, tz_k + x, \dots, tz_k + x, d_k)$. Where z_k, d_k are e -positive and $\lim_{k \rightarrow \infty} z_k = z$, $\lim_{k \rightarrow \infty} d_k = d$. Then the coefficients of polynomials P_k converge to the coefficients of the

polynomial P . It follows from part 4 of Fact A.1 that the roots of P_k are real . Since $a_{n-1} > 0$ hence using Fact A.2 we get that the roots of P are also real . This exactly means that the polynomial p_d is hyperbolic in direction z . The d -nonnegativity statement follows from the nonnegativity part of Fact A.2 . ■

We are ready now to present our proof of Theorem 2.2 .

Proof: (Proof of Theorem 2.2) . The "only if" part is simple . Indeed supposed that there exists a subset $S \subset \{1, 2, \dots, n\}$ such that $\text{Rank}_p(\sum_{i \in S} x_i) < |S|$, i.e. using the identities (15) $M_p(k, k, \dots, k, d, \dots, d) = 0$, where $k = \sum_{i \in S} x_i$, $d \in C_e(p)$ is e -positive and the n -tuple $(k, k, \dots, k, d, \dots, d)$ consists of $|S|$ copies of $k = \sum_{i \in S} x_i$. Let d be any e -positive positive vector such that $d - x_i$ is e -nonnegative , $1 \leq i \leq n$. Using the monotonicity of p -mixed forms we get that

$$M_p(x_1, \dots, x_n) \leq M_p(k, k, \dots, k, d, \dots, d) = 0.$$

Our proof of the "if" part is by induction in the degree n . Suppose that the generalized Rado conditions (16) hold . Then at least $\text{Rank}_p(x_n) \geq 1$. Consider the following homogeneous polynomial of degree $n - 1$:

$$p_d(x) = M_p(x, x, \dots, x, d), \quad d = x_n.$$

We get from Proposition a.4 the following assertion :

The polynomial $p_d(x)$ is hyperbolic in direction $z = \sum_{1 \leq i \leq n-1} x_i$ and the vectors $x_i \in N_z(p_d)$, $1 \leq i \leq n - 1$, i.e. are z -nonnegative respect to the polynomial p_d .

Indeed , it follows from the generalized Rado conditions (16) that $\text{Rank}_p(z) \geq n - 1$ and $\text{Rank}_p(z + d) = \text{Rank}_p(\sum_{1 \leq i \leq n} x_i) = n$.

Next we show that the $n-1$ -tuple $\mathbf{Y} = (x_1, \dots, x_{n-1})$ satisfies the generalized Rado conditions

:

$$\text{Rank}_{p_d}(\sum_{i \in S} x_i) \geq |S| \quad \text{for all } S \subset \{1, 2, \dots, n - 1\}.$$

Or equivalently , that

$$M_p(k, \dots, k, z, \dots, z, d) > 0; k = \sum_{i \in S} x_i, z = \sum_{1 \leq i \leq n-1} x_i, d = x_n, S \subset \{1, \dots, n - 1\}, \quad (26)$$

where the n -tuple $\mathbf{T} = (k, \dots, k, z, \dots, z, d)$ consists of $|S|$ copies of k , $n - 1 - |S|$ copies of z and one copy of d .

It is easy to see that the generalized Rado conditions for the n -tuple \mathbf{T} are implied by the generalized Rado conditions for the original n -tuple $\mathbf{X} = (x_1, \dots, x_{n-1}, x_n)$. Since the n -tuple $(k, \dots, k, z, \dots, z, d)$ consists of at most three distinct components hence we can apply part 2 of Corollary A.3 . Therefore we get that indeed

$$\text{Rank}_{p_d}(\sum_{i \in S} x_i) \geq |S| \quad \text{for all } S \subset \{1, 2, \dots, n - 1\}. \quad (27)$$

Thus , by induction in the degree , we get that p_d -mixed form $M_{p_d}(x_1, \dots, x_{n-1}) > 0$: the polynomial p_d of degree $n - 1$ in m real variables is z -hyperbolic . But

$$M_{p_d}(x_1, \dots, x_{n-1}) = \frac{\partial^{n-1}}{\partial \alpha_1 \dots \partial \alpha_{n-1}} p_d(\sum_{1 \leq i \leq n-1} \alpha_i x_i) = \\ = \frac{\partial^{n-1}}{\partial \alpha_1 \dots \partial \alpha_{n-1}} M_p(\sum_{1 \leq i \leq n-1} \alpha_i x_i, \dots, \sum_{1 \leq i \leq n-1} \alpha_i x_i, x_n) = (n-1)! M_p(x_1, \dots, x_n).$$

We conclude that if Theorem 2.2 is true for $n-1$ then it is also true for n , and the case " $n=1$ " is trivially true. ■

B Proof of Proposition 2.6

Proof: Assume wlog that $q(\alpha_1, \dots, \alpha_n) = 1$. It follows from the Euler's identity that

$$\sum_{1 \leq i \leq n} \alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) = n.$$

Let $q(\alpha_1, \dots, \alpha_n) = \sum_{(r_1, \dots, r_n) \in \text{supp}(q)} a_{(r_1, \dots, r_n)} \prod_{1 \leq i \leq n} \alpha_i^{r_i}$. Define (positive numbers) $b_{(r_1, \dots, r_n)} = a_{(r_1, \dots, r_n)} \prod_{1 \leq i \leq n} \alpha_i^{r_i}$, $(r_1, \dots, r_n) \in \text{supp}(q)$. Then $\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) = \sum_{(r_1, \dots, r_n) \in \text{supp}(q)} r_i b_{(r_1, \dots, r_n)}$.

Suppose that for some subset $S \subset \{1, 2, \dots, n\}$, $1 \leq |S| < n$ we have the inequality $\sum_{i \in S} r_i < |S|$ for all $(r_1, \dots, r_n) \in \text{supp}(q)$. Then $\sum_{i \in S} \alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) \leq |S| - 1$. But the condition (18) says that $\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) = 1 + \delta_i$ and $\sum_{1 \leq i \leq n} |\delta_i|^2 \leq \frac{1}{n}$. By the Cauchy-Schwarz inequality, $\sum_{i \in S} |\delta_i| \leq \sqrt{|S|} < 1$. Therefore,

$$\sum_{i \in S} \alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n) \geq |S| - \sum_{i \in S} |\delta_i| > |S| - 1.$$

The last inequality gives a contradiction. ■

C A sketch of a proof of Corollary 2.4

Proof: By Theorem 2.2 the conditions (1) and (2) are equivalent. (2) implies (3) for any homogeneous polynomial with nonnegative coefficients.

Let $\alpha_i = e^{y_i}$, $1 \leq i \leq n$, $\sum_{1 \leq i \leq n} y_i = 0$. Consider the following convex functional

$$f(y_1, \dots, y_n) = \log(q(e^{y_1}, e^{y_2}, \dots, e^{y_n})).$$

Here $q(x)$, $x \in R^n$ is a homogeneous polynomial of degree n in n real variables with nonnegative coefficients. Then

$$\frac{\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n)}{q(\alpha_1, \dots, \alpha_n)} = \frac{\partial}{\partial y_i} f(y_1, \dots, y_n), 1 \leq i \leq n.$$

Notice the condition (3) is equivalent to the following condition:

$$\inf_{y_1 + \dots + y_n = 0} f(y_1, \dots, y_n) = L > -\infty.$$

Consider the anti-gradient flow , i.e. the system of differential equations

$$y_i(t)' = -(\frac{\partial}{\partial y_i} f(y_1, \dots, y_n) - 1), y_i(0) = 0; 1 \leq i \leq n.$$

It is well known that in this convex case the gradient flow is defined for all $t \geq 0$. Using the Euler's identity we get that

$$\frac{d}{dt} f(y_1(t), \dots, y_n(t)) = -\beta(t) =: - \sum_{1 \leq i \leq n} \left| \frac{\alpha_i \frac{\partial}{\partial \alpha_i} q(\alpha_1, \dots, \alpha_n)}{q(\alpha_1, \dots, \alpha_n)} - 1 \right|^2$$

It is easy to see that , because of the convexity of f , a nonnegative function $\beta(t)$ is non-increasing on $[0, \infty)$.

As $\inf_{y_1+\dots+y_n=0} f(y_1, \dots, y_n) = L > -\infty$ thus $\int_0^\infty \beta(t) dt < \infty$. Thus $\lim_{t \rightarrow \infty} \beta(t) = 0$. This proves the implication (3) \rightarrow (4) for all homogeneous polynomials of degree n in n real variables with nonnegative coefficients .

The implication (4) \rightarrow (5) is obvious . The implication (5) \rightarrow (6) for general homogeneous polynomials of degree n in n real variables with nonnegative coefficients is Proposition 2.6 .

Finally , the condition (6) is equivalent to (2) . ■