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2

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TITLE: MULTIDOMAIN, MULTIRECORD-TYPE DATATRIEVE-11 DATABASES

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MULTIDOMAIN, MULTIRECORD-TYPE DATATRIEVE-11 DATA BASES

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ABSTRACT

Data bases consisting of multiple domains and multirecord-type domains present special problems in design, loading, maintenance, and report generation. The logical association of records is a fundamental concern in all these problem areas. This paper describes techniques for dealing with this and other specifics using Datatrieve-11, Sort-11, FORTRAN 77, and the RSX-11M Indirect Command File Processor.

INTRODUCTION

Hierarchical data bases can be simulated using the OCCURS* clause and its variant the OCCURS DEPENDING. However, these will result in much wasted disk space if the number of OCCURRENCES has a wide range and if the lengths and numbers of fields are large. This difficulty can be largely overcome by defining data bases consisting of single domains with several record redefinitions or multiple domains whose records are linked by pointers.

Retrieval of single records from moderate-sized files (5000-10000 records) using Datatrieve-11 is generally rapid enough to avoid operator fatigue, especially if keyed access is used. Similarly, reports made from these files usually can be produced within several minutes. However, creating comprehensive reports from multiple files or multirecord-type files may take unacceptable amounts of time --- ranging into many hours. (We have had some reports that took more than a day.) Compiled programs to replace Datatrieve procedures can provide multiple-order-of-magnitude reductions in the time required to create reports from such data bases and also overcome the Datatrieve-11 pool problems associated with simultaneously reading several domains.

MULTIRECORD DOMAINS

A hierarchical data base can be simulated using the OCCURS and OCCURS DEPENDING clauses to store the "children" data in the same record as the "parent". If the number of children fields actually used is fairly consistent from record to record, this method provides a convenient and efficient data organization.

However, if the number of children varies considerably from parent to parent, but the number of bytes required for each child is about the same as for the parent, more efficient use of mass storage devices can be achieved using the REDIFINES clause to create separate children records for each parent within the same file. Now the number of children can vary from zero to an arbitrarily large number with little or no penalty in wasted disk space.

Figures 1 and 2 illustrate the use of the multirecord type domain. Figure 1 shows the front and back of a worksheet used to record data needed to install digital data communications services for one of our computer users. The worksheet is broken down into several sections. Not all sections need be used, and some of those that are may occur more than once.

Figure 2 shows the complete domain definition used to specify the organization of the worksheet's data in a file. This domain is called CWR for Communications Work Request. Note first that all records associated with a specific worksheet are assigned the same value in one field, CWR NO, thus associating them with one another. To speed access to all records for a particular worksheet we make this field the primary key for the file corresponding to this domain.

The GENERAL record may be regarded as the parent in this hierarchical arrangement; all other records are children. Note that all children do not share the same record definition. They are all more or less similar in length; some may occur more than 20 times for one parent.

* Words appearing in UPPER CASE are either Datatrieve key words or field names within the data bases discussed in this paper.

A method of associating a record redefinition with a specific record is needed. The REC-TYP field performs this function. The record associated with the GENERAL section of the worksheet is assigned a REC-TYP value of 1; other sections are also given unique numeric REC-TYP entries.

Also needed is a way to assure that all fields in a new record are initialized appropriately even though the user may not explicitly enter data in them. Datatrieve will initialize fields according to the first record definition it finds, not according to the remaining field definitions in the current REDIFINES clause. The DIMMYR field defined in figure 2 assures that all unused fields beyond S11F will be set to ASCII blanks. Of course this technique is inappropriate for numeric fields, but it provides a reasonable compromise.

Other, more effective, ways to initialize all fields include creating separate domain definitions for each record type within the file and using these exclusively for data entry, and/or defining data entry procedures that explicitly initialize all fields for a given record type. Either of these methods is, by itself, sufficient to guarantee appropriate initialization.

Separate domain definitions for each record in a multirecord type file can also help alleviate the well-known Datatrieve-11 pool problem because, individually, they are much shorter than the complete domain definition. They can be used individually, along with corresponding data entry procedures and appropriate data validation tables, where the complete domain definition would cause difficulty. Figure 3 shows the domain definition for the GENERAL section alone, and Figure 4 illustrates one data entry procedure used with this domain.

It is prudent to let Datatrieve, and not the user, assign values to fields that are used to maintain the internal organization of the data. In Figures 2 and 3, these fields are CWR NO, which identifies the parent and its children, and REC TYP, which identifies the redefinition to be used for the specific record. Each set consisting of parent and children must be distinguished from all other sets. One way to maintain a unique identifier in this case CWR NO, is to store a number in a separate file with its corresponding domain definition. Each time a new parent is created, the number can be fetched, used, incremented, and replaced. The first few lines of the procedure in figure 4 do this. The record type identifier can simply be set in the data entry procedure, as in line 27 of figure 4.

Users often find it helpful to have a minimum number of procedures with which to deal while entering and updating data. One way of achieving this end is to combine S11RF and MODIFY processes in the same procedure. Figure 4 illustrates how this can be done for the parent record. In this case, the user is assumed to know whether he wishes to modify an existing record or create a new one. However, he may not know whether a specific record exists and may not wish to bother finding out. Figure 4 illustrates one way of letting Datatrieve search for a record and, if the record is found, display its contents so the user can modify it; otherwise the record is created. Note the use of nested FOR clauses; they are very handy in applications such as this.

Often a user wishes to enter data for a parent and several children in one session. Ideally, one would include all the necessary code in a single procedure. Unfortunately, the Datatrieve 11 pool problem is likely

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COMMUNICATION WORK REQUEST[illegible]

DATE	Day	Month	Year	Time	Page	Operator	Site	System	City										
	1986			12:00		1234567													
	1986			12:00		1234567													
	1986			12:00		1234567													
TIME	1986			12:00		1234567													
	1986			12:00		1234567													
	1986			12:00		1234567													
	1986			12:00		1234567													

[illegible][illegible][illegible][illegible]

Figure 1. Communications Work Request

[illegible]

Figure 2. The communications work request data base

```

001110 NEWTEMP1;
001120 DO WHILE NEWTEMP1;
001130   SETTING NEWTEMP1 REC ON BUS1278.31CUR DAT;
001140 NEWTEMP1 REC;
001150 WRITE RECORD NEWTEMP1 REC
001160   USING
001170 1 NEWTEMP1 REC
001180   5 CUR NO          PIC IS 9(6)  USAGE IS INTEGER
001190   5 REC-TYP         PIC IS 9(2)  USAGE IS INTEGER
001200   5 SITE            PIC IS 9(2)
001210   5 GENERAL
001220   10 DEV-PTH        PIC IS 9(6)  USAGE IS INTEGER
001230   10 FWD-PTH        PIC IS 9(6)  USAGE IS INTEGER
001240   10 PROJECT-NO     PIC IS 9(6)
001250   10 MEMO-DATE      USAGE IS DATE
001260   10 LOG-DATE       USAGE IS DATE
001270   10 COMP-DATE      USAGE IS DATE
001280   10 INPUT-DATE     USAGE IS DATE
001290   10 TOPO-UPDATE    USAGE IS DATE
001300   10 HTR            PIC IS 9
001310   10 FILLER        PIC IS 9(22)

```

Figure 3. The domain and revised definitions for revised type 1

```

IDU1: (270,3)LOGGENRAL.CMD
DELETE LOGGENRAL;
DEFINE PROCEDURE LOGGENRAL
FINISH
SET ABORT
DECLARE BCK-PTR PIC IS 9(6) USAGE IS INTEGER.
DECLARE CUR-CWR PIC IS 9(6) USAGE IS INTEGER.
DECLARE CUR-REV PIC IS 9(6) USAGE IS INTEGER.
DECLARE CUR-FWD PIC IS 9(6) USAGE IS INTEGER.
DECLARE PRV-FWD PIC IS 9(6) USAGE IS INTEGER.
DECLARE NEW-REC PIC IS X.
NEW-REC = "F"
BCK-PTR = #
READY NEXTPSR WRITE
IF *. "C if you wish to create a new CWR or M to modify an old one" EQ "C" THEN
  FOR NEXTPSR BEGIN
    NEW-REC = "Y"
    CUR-CWR = NEXT-PSR
    PRINT COL 2#, "This CWR's number is", CUR-CWR(-)
    MODIFY USING NEXT-PSR = CUR-CWR + 1
  END ELSE
    CUR-CWR = *. "number of the CWR you wish to modify"
FINISH
READY CWR SHARED WRITE
BCK-PTR = *. "previous CWR, if any (6 digits)"
IF NEW-REC = "Y" THEN STORE CWR USING BEGIN
  CWR-NO = CUR-CWR
  REC-TYP = 1
  REV-PTR = CUR-CWR
  FWD-PTR = CUR-CWR
  COMPL-DATE = " "
  INPUT-DATE = "TODAY"
  PROJECT-NO = *. "project number (6 chars.)"
  SITE = *. "site (2 chars.)"
  MEMO-DATE = *. "requestor's memo date (mm/dd/yy)"
  LOG-DATE = *. "log book date (mm/dd/yy)"
  TOPO-UPDATE = " "
  XFER = " "
END ELSE
IF *. "Y if you wish to modify the GENERAL section" EQ "Y" THEN BEGIN
  FOR CWR WITH (CWR-NO=CUR-CWR) AND (REC-TYP=1) MODIFY USING BEGIN
    PRINT SKIP, COMPL-DATE, XFER, SKIP
    COMPL-DATE = *. "completed date (mm/dd/yy)"
    XFER = *. "Y if this CWR has been verified for transfer to TOPO"
  END
  FOR CWR WITH (CWR-NO=CUR-CWR) AND (REC-TYP=1)
    MODIFY USING BEGIN
    PRINT SKIP, PROJECT-NO, SITE, MEMO-DATE, LOG-DATE, SKIP
    PROJECT-NO = *. "project number (6 chars.) or <tab>"
    SITE = *. "site (2 chars.) or <tab>"
    MEMO-DATE = *. "requestor's memo date (mm/dd/yy) or <tab>"
    LOG-DATE = *. "Log Book Date (mm/dd/yy) or <tab>"
  END
END
FOR CWR WITH (CWR-NO=CUR-CWR) AND (REC-TYP=1)
  WHILE SITE NOT IN SITETBL MODIFY USING BEGIN
    PRINT SKIP, "Bad site code ('. SITE(-), '); please try again.", SKIP
    SITE = *. "site (2 chars.)"
  END
END
RELEASE SITETBL
END-PROCEDURE;

```

Figure 4. Data entry procedure for record type 1

```

DELETE LOGREQSTR;
DEFINE PROCEDURE LOGREQSTR
IF *. "Y if you wish to enter data for REQUESTOR" EQ "Y" THEN BEGIN
  NEW-REC = "Y"
  FOR CWR WITH (CWR-NO = CUR-CWR) AND (REC-TYP E) 2) MODIFY USING BEGIN
    NEW-REC = "Y"
    PRINT SKIP, NAME, ORG, MAIL-STOP, PHONE, CHARGE, SKIP
    FIRST-NAME = *. "requestor's first name (13 chars.) or <tab>"
    LAST-NAME = *. "requestor's last name (15 chars.) or <tab>"
    DIV = *. "requestor's division (4 chars.) or <tab>"
    GRP = *. "requestor's group (3 chars.) or <tab>"
    MAIL-STOP = *. "requestor's mail stop (4 chars.) or <tab>"
    PHONE = *. "requestor's phone (11 digits) or <tab>"
    COST-CTR = *. "cost center (4 digits) or <tab>"
    PROG-COD = *. "program code (4 chars.) or <tab>"
  END
  IF NEW-REC NE "Y" THEN STORE CWR USING BEGIN
    CWR-NO = CUR-CWR
    REC-TYP = 2
    FIRST-NAME = *. "requestor's first name (13 chars.)"
    LAST-NAME = *. "requestor's last name (15 chars.)"
    DIV = *. "requestor's division (4 chars.)"
    GRP = *. "requestor's group (3 chars.)"
    TA = " "
    BLDG = " "
    ROOM = " "
    MAIL-STOP = *. "requestor's mail stop (4 chars.)"
    PHONE = *. "requestor's phone (11 digits)"
    COST-CTR = *. "cost center (4 digits)"
    PROG-COD = *. "program code (4 chars.)"
  END
END
END-PROCEDURE

```

Figure 5. Data entry procedure for record type 2

to prevent this. A straightforward "workaround" for this is to write an indirect command file that creates and passes instructions to Datatrieve for execution. Figure 6 shows one such command file. Unfortunately, this technique is slow and tedious because of the time required to initialize Datatrieve with each call. We are experimenting with ways to reduce this initialization time.

Data retrieval and report generation from multirecord-type domains are considerably more complex than from non-hierarchical single-record-type domains. Because no user-written logic can be included in Datatrieve Report Writer procedures, the data to be contained in a report must first be transferred to a "flat" file. A procedure to do this is shown in Figure 7, and the domain and record definitions are shown in Figure 8. Note the use of nested FOR clauses to produce the work file. These can cause Datatrieve to make millions of disk accesses taking many hours if a unique key is not used in describing the record to be selected in the inner FOR loop. For each record in the outer FOR loop, about one-half the domain indicated in the inner FOR loop will be searched for the correct record. Consider the 5000-plus record multirecord-type file used in the above illustrations. When the work file was created with no keys, the procedure of Figure 7 took over 20 h of clock time. This was true for Datatrieve running under PSX11M or RSX11M+ on a PDP11/70 with no other tasks active, except system overhead. After we modified the procedure of Figure 7 to create the file PRELOG with CWR-NO as the primary key with no duplications, the procedure took about 1 h of clock time.

Interactive use of similar nested FOR clauses to retrieve single-parent children sets from within a single multirecord-type domain takes only seconds or tens of seconds. If keys are used in specifying the desired record, thus, operator fatigue is not excessive, provided the multiline nested FOR statements are stored in "command" procedures and need not be entered by the operator.

```

.DUI:(270,3)UPDATECWR.CMD
.ENABLE QUIET
.ENABLE SUBSTITUTION
.SETS CLI <CLI>
.IF CLI EQ "DCL" SET MCR=T1;
PIP (270,3)*.DAT/UN/NM
.1001
.ASKS (213,1305) SECT Enter section you wish to
update
.IF SECT EQ "GEN" .GOTO 200
.IF SECT EQ "REQ" .GOTO 200
.IF SECT EQ "WRK" .GOTO 200
.IF SECT EQ "NEW" .GOTO 200
.IF SECT EQ "OLD" .GOTO 200
.IF SECT EQ "FRM" .GOTO 200
.IF SECT EQ "TO" .GOTO 200
.IF SECT EQ "REM" .GOTO 200
.IF SECT EQ "INS" .GOTO 200
.IF SECT EQ "COM" .GOTO 200
.DISABLE QUIET
[Bad abbreviation; please try again.
Valid abbreviations are:
GEN
REQ
WRK
NEW
OLD
FRM
TO
REM
INS
COM
.ENABLE QUIET
.GOTO 100
.2001
.SETS UPDATE "UPD"+"SECT"
.OPEN UPDATECWR.DTR
.ENABLE DATA
.SET DICTIONARY DUI:(270,3)PSR.DIC
['UPDATE'
.DISABLE DATA
.CLOSE
DTR UPDATECWR.DTR
DTR UPDATECWR.DTR /*DE/NM
.ASK CONTIN Have you more CWR's to process
.IF CONTIN .GOTO 100
.IF CLI NE "MCR" SET /'CLI'=T1;
.EXIT

```

Figure 6. Indirect command file for generalized data entry

```

1001:(270,3)PRELOGRPT.CMD
DELETE PRELOGRPT;
DEFINE PROCEDURE PRELOGRPT
FINISH
DEFINE FILE FOR PRELOG SUPERSEDE
READY PRELOG WRITE
READY CWR SHARED READ
DECLARE CWR-TMP PIC IS 9(16);
DECLARE PORT-TMP PIC IS 9(16);
DECLARE ASS-TMP USAGE IS DATE;
DECLARE COMP-TMP USAGE IS DATE;
DECLARE BY-TMP PIC IS 9(16);
DECLARE STAT-TMP PIC IS 9(16);
FOR AA IN CWR WITH REC-TYP=1 AND SITE NE "EP"
STORE Z IN PRELOG USING BEGIN
Z.CWR-NO = AA.CWR-NO
Z.REC-TYP = AA.REC-TYP
Z.INPUT-DATE = AA.INPUT-DATE
Z.LOG-DATE = AA.LOG-DATE
Z.SITE = AA.SITE
Z.PROJECT-NO = AA.PROJECT-NO
Z.MEMO-DATE = AA.MEMO-DATE
Z.ASSIGNED-DATE = #
Z.COMPLETION-DATE = #
Z.BY = #
Z.STATUS = #
END
FOR A IN CWR WITH REC-TYP=2 BEGIN
CWR-TMP = CWR-NO
PORT-TMP = C.PR-NO
FOR Z IN PRELOG WITH CWR-NO=CWR-TMP
MODIFY USING BEGIN
Z.REQ-FIRST = A.FIRST-NAME
Z.REQ-LAST = A.LAST-NAME
Z.REQ-ORG = A.DIVISION.GRP
Z.MAIL-STOP = A.MAIL-STOP
Z.PHONE = A.PHONE
END
END
FOR B IN CWR WITH REC-TYP=4 BEGIN
CWR-TMP = CWR-NO
PORT-TMP = B.PR-NO
FOR Z IN PRELOG WITH CWR-NO=CWR-TMP
MODIFY USING BEGIN
Z.NEW-PORT-NO = PORT-TMP
Z.NEW-PART = B.PAR
Z.NEW-STATUS = B.PR-STA
Z.NEW-SPEED = B.PR-SPD
Z.NEW-TYPE = B.PR-TYP
Z.TO-DEV-TYPE = B.DEV-TYP
Z.TO-DEV-NO = B.DEV-NO
END
END
FOR C IN CWR WITH REC-TYP=5 BEGIN
CWR-TMP = CWR-NO
PORT-TMP = C.PR-NO
FOR Z IN PRELOG WITH CWR-NO=CWR-TMP
MODIFY USING BEGIN
Z.OLD-PORT-NO = PORT-TMP
Z.OLD-PART = C.PAR
Z.OLD-STATUS = C.PR-STA
Z.OLD-SPEED = C.PR-SPD
Z.OLD-TYPE = C.PR-TYP
Z.FROM-DEV-TYPE = C.DEV-TYP
Z.FROM-DEV-NO = C.DEV-NO
END
END
FOR D IN CWR WITH REC-TYP=6 BEGIN
CWR-TMP = CWR-NO
FOR Z IN PRELOG WITH CWR-NO=CWR-TMP
MODIFY USING BEGIN
Z.FROM-TA = D.TA
Z.FROM-BLDG = D.BLDG
Z.FROM-ROOM = D.ROOM
Z.FROM-BOX-NO = D.TRM-BOX
END
END
FOR E IN CWR WITH REC-TYP=7 BEGIN
CWR-TMP = CWR-NO
FOR Z IN PRELOG WITH CWR-NO=CWR-TMP
MODIFY USING BEGIN
Z.TO-TA = E.TA
Z.TO-BLDG = E.BLDG
Z.TO-ROOM = E.ROOM
Z.TO-BOX-NO = E.TRM-BOX
END
END
FOR F IN CWR WITH REC-TYP=9 BEGIN
CWR-TMP = CWR-NO
ASS-TMP = F.ASSIGNED-DATE
COMP-TMP = F.COMPLETION-DATE
BY-TMP = F.CUT-BY
STAT-TMP = F.CUT-STATUS
FOR Z IN PRELOG WITH CWR-NO=CWR-TMP
MODIFY USING BEGIN
Z.ASSIGNED-DATE = ASS-TMP
Z.COMPLETION-DATE = COMP-TMP
Z.BY = BY-TMP
Z.STATUS = STAT-TMP
END
END
FOR G IN CWR WITH REC-TYP=11 BEGIN
CWR-TMP = CWR-NO
ASS-TMP = G.ASSIGNED-DATE
COMP-TMP = G.COMPLETION-DATE
BY-TMP = G.CUT-BY
STAT-TMP = G.CUT-STATUS
FOR Z IN PRELOG WITH CWR-NO=CWR-TMP
MODIFY USING BEGIN
Z.ASSIGNED-DATE = ASS-TMP
Z.COMPLETION-DATE = COMP-TMP
Z.BY = BY-TMP
Z.STATUS = STAT-TMP
END
END
END-PROCEDURE

```

Figure 7. Procedure to create flat work file for the Report writer

```

DELETE PRELOG;
DEFINE DOMAIN PRELOG
  USING PRELOGREC ON DU1:(270,3)PRELOG.DAT;
DELETE PRELOGREC;
DEFINE RECORD PRELOGREC
  USING
1 PRELOG-REC.
  5 CWR-NO PIC IS X(4).
  5 REC-TYP PIC IS X(2).
  5 SITE PIC IS X(2).
  5 PROJECT-NO PIC IS X(6).
  5 REQ-ORG PIC IS X(7).
  5 OLD-PORT-NO PIC IS X(8).
  5 FROM-TA PIC IS X(2).
  5 FROM-BLOG PIC IS X(5).
  5 FROM-ROOM PIC IS X(6).
  5 FROM-BOX-NO PIC IS X(7).
  5 FROM-DEV-TYPE PIC IS X(3).
  5 FROM-DEV-NO PIC IS X(8).
  5 NEW-PORT-NO PIC IS X(8).
  5 TO-TA PIC IS X(2).
  5 TO-BLDG PIC IS X(5).
  5 TO-ROOM PIC IS X(6).
  5 TO-BOX-NO PIC IS X(7).
  5 TO-DEV-TYPE PIC IS X(3).
  5 TO-DEV-NO PIC IS X(8).
  5 BY PIC IS X(12).
  5 STATUS PIC IS X(9).
  5 REQ-FIRST PIC IS X(6).
  5 REQ-LAST PIC IS X(10).
  5 MAIL-STOP PIC IS X(4).
  5 OLD-PART PIC IS X(2).
  5 OLD-STATUS PIC IS X(1).
  5 OLD-SPEED PIC IS X(5).
  5 OLD-TYPE PIC IS X(3).
  5 NEW-PART PIC IS X(2).
  5 NEW-STATUS PIC IS X(1).
  5 NEW-SPEED PIC IS X(5).
  5 NEW-TYPE PIC IS X(3).
  5 PHONE PIC IS X(4) EDIT-STRING IS XXXX.
  5 INPUT-DATE USAGE IS DATE.
  5 LOG-DATE USAGE IS DATE.
  5 ASSIGNED-DATE USAGE IS DATE.
  5 COMPLETION-DATE USAGE IS DATE.
  5 MEMO-DATE USAGE IS DATE.

```

Figure 8. Domain and record definitions for the flat work file

MULTIDOMAIN DATA BASES

When the lengths of the children records are substantially different from one another or from the parent record, and the number of occurrences of children varies markedly from parent to parent, a multidomain data base may be appropriate. Similarly, if relationships other than the single parent-child association are to be shown, a multidomain data base may be required.

As with the multirecord-type domain, a method of linking records is needed. One or more sets of pointers may be satisfactory for this purpose. Something as simple as the common-valued field, like CWR-NO in the multirecord-type domain above, may suffice. All the data entry and retrieval problems described above for multirecord-type domains apply here, but most are more pronounced. Using VIFWS helps to solve these problems, but the limitations of the Datatrieve II pool are very noticeable. Therefore using other software devices, including indirect command files and compiled programs, becomes highly desirable.

USE OF SUPERVISOR MODE RMS WITH DATATRIEVE II

Another paper scheduled for presentation here at the symposium discusses the linking of Datatrieve-II to an RMS supervisor mode library. With help from our local DEC office, we are now running DTRSUP on our PDP11/70 under RSX11M+. Figures 9, 10, and 11 show the task builder command file and ODL files used to create the supervisor mode Datatrieve-II. We have experienced no noticeable improvement in execution time; however, the roughly 40% increase in pool space has resulted in much greater convenience when creating procedures and in ad hoc queries. Of particular note is the new ability to READY five or six domains at once and still be able to perform useful searches, and so forth. We had found that with overlaid RMS two domains were about the maximum that could be readied at once and do much else.

```

DTRSUP/PP/CP,DTRMP/-SP-DTRIMP/MP
1 LINK TO SUPERVISOR MODE RMSRES
REC'DUP=13,641RMSRES/SV/B
1 FOLLOWING LINE APPLIES ONLY TO NON-DBMS-II DATATRIEVE
PAR-GEN:1774#0
ASC-T1:0
UNIT-S=10
TAS=...DTR
//

```

Figure 9. Task builder command file for building Datatrieve II with Supervisor mode RMS*

```

1 THE BUILD ROOTS ARE:
1
1 DTRDBM -- DBMS SUPPORT
1 DTRRMS -- NO DBMS SUPPORT (RMS ONLY)
1 DTRRDB -- RMS DEBUG VERSION
1 DTRFPE -- FLOATING POINT EMULATION INCLUDED
1 (RMS ONLY)
1 DTRDBG -- DBMS DEBUG VERSION
1
1 .ROOT DTRRMS
0DTRMP.ODL
.END
1
1 Figure 10. First ODL file for building Datatrieve-II with Supervisor-mode RMS*
1
1 DTRMP.ODL - LINKS TO SUPERVISOR MODE RMSRES
1
1 FACTORS ENDING IN DOUBLE DIGITS INDICATE CONDITIONAL BUILD FOR
1 THE VARIOUS DATATRIEVE VERSIONS. THE CODES ARE:
1
1 XXX#01 -- FPP PRESENT OR NOT NEEDED
1 XXX#02 -- FLOATING EMULATION CODE NEEDED
1
1 .PSECT 100V,GBL
1 .PSECT INIT,GBL
1
1 FOR RMS V2.0 AS SUPER MODE, REMOVE REFERENCES TO "TREE#2"
1
1 DTRRMS: .FCTR RT01-101-1(DD1,COMPLR,PAR#01,EXEC#01)
1 DTRFMS: .FCTR RT01-101-1(DD1,COMPLR,PAR#01,EXEC#01)
1 DTRHVB: .FCTR RT01-FP1-103-1(DD1,COMPLR,PAR#01,EXEC#01)
1
1 DTRRDB: .FCTR RT01-101-DEBUG-1(DD1,COMPLR,PAR#01,EXEC#01),BDTTR
1 DTRHVB: .FCTR RT01-FP1-103-DEBUG-1(DD1,COMPLR,PAR#01,EXEC#01),BDTTR
1
1 ROOT SECTIONS
1
1 RT#1: .NAME DTR
1 .FCTR DTR-RT2-RT3-DTRLB/LB:RMSDB
1
1 RT2: .FCTR DTRLB/LB:TEST:AL:BL:CK:JTD:PS:MS:DA
1 RT3: .FCTR DTRLB/LB:OD:SAVREG
1
1 1 ADD BASE RMS SUPPORT, REQUIRED ROOT SEGMENTS, AND FP STUFF
1 1 FOR RMS V2.0 SUPPORT.
1
1 RMSROT: .FCTR LB:11,1RMSL10/LB:R0AUTS:R0IMPA:R0EXSY:R0SSYM-FP2
1
1 1/O GUYS
1
1 101: .FCTR DTRLB/LB:RM-10/ RMSROT
1 102: .FCTR DTRLB/LB:10:RM:RA
1 103: .FCTR 102-DTRLB/LB:RMS21
1
1 1 INITIALIZATION STUFF.
1
1 INI1: .FCTR INIT-INI
1 INI1: .FCTR DTRLB/LB:IN
1
1 1 DEBUGGING STUFF
1
1 DEBUG: .FCTR DTRLB/LB:DBPRX:AUTO:PATCH-DTRLB/LB:BDT/DA
1 .NAME BDTX
1 BDTTRE: .FCTR BDTX-1(DD2-BD3-BD4,BDB1)
1 BDB1: .FCTR DTRLB/LB:BDT2
1 BDB2: .FCTR DTRLB/LB:BDT3
1 BDB3: .FCTR DTRLB/LB:BDT4
1 BDB4: .FCTR DTRLB/LB:DBPR
1
1 1 DDDO-TRIEVE
1
1 DD1: .FCTR DTRLB/LB:DD
1
1 1 PARSER
1
1 PAR#01: .FCTR PARSE-100V-FP1-FP2,ADT1,EDIT,CHND#01
1 P1: .FCTR DTRLB/LB:DI:HT:LM:CD:PR:PU
1 P2: .FCTR DTRLB/LB:ME:SM
1 ADT1: .FCTR DTRLB/LB:AD
1
1 1 COMMANDS
1
1 CHND#01: .FCTR C1-C1C2,INJ1-C3,RUF1#01
1 C1: .FCTR DTRLB/LB:RP:PH
1 C2: .FCTR DTRLB/LB:PA
1 C3: .FCTR DTRLB/LB:LM:TL
1
1 1 READY/FINISH
1
1 RFI#01: .FCTR RFI-RF2
1
1 RF1: .FCTR DTRLB/LB:RF
1 RF2: .FCTR DTRLB/LB:UP:ML:OF
1
1 1 EDITOR
1
1 EDIT1: .FCTR DTRLB/LB:TE:1,TPARS
1
1 1 COMPILER-ISH STUFF
1
1 NAME CMPLR
1 COMPLR: .FCTR CMPLR-DTRLB/LB:CA:RC:AS:PF:SC:TC
1
1 1 EXECUTION SYSTEM.
1
1 NAME RUNTIM
1
1 EXEC#01: .FCTR RUNTIM-EX1-EX2-EX3-EX6-1(EN4,EN5)
1 EX1: .FCTR DTRLB/LB:AP:CH:EX:RS:RU:US
1 EX2: .FCTR DTRLB/LB:NUMTIN:GETTIN
1 EX3: .FCTR DTRLB/LB:RM:LU:CO:HOPL-SORT
1 EX4: .FCTR DTRLB/LB:ME
1 EX5: .FCTR DTRLB/LB:DE
1 EX6: .FCTR DTRLB/LB:ROFM
1 EX7: .FCTR DTRLB/LB:FM:FDV:FDVPR:FDVERR:FDVDT:FDVTIO:FMSSAV
1
1 EXEC#02: .FCTR RUNTIM-EX1-EX2-EX3-1(EN3-EN4,EX7)
1
1 1 SORT
1
1 SORT: .FCTR DTRLB/LB:SORTS:CIORMS:ST
1
1 1 THE CO-TREE PLUS OUR STUFF, (ELIMINATED FROM SUPER MODE RMS, FPP MOVED
1 TO ROOT)
1
1 11R000: .FCTR RM011-11FPP,RMSFIL,RMSDEC,000*RE1
1 11R001: .FCTR RM011-11FPP,RMSFIL,RMSDEC1
1
1 P1: .FCTR DTRLB/LB:R0FP:MISC
1 P2: .FCTR DTRLB/LB:FP:MISC
1 .PSECT 0TRVSE
1 .PSECT 000FOV,GBL,AV

```

Figure 11. Second ODL file for building Datatrieve II with Supervisor mode RMS*

One would like to take advantage of the added pool space to speed up procedures by reducing disk I/O. Figure 12 illustrates how one might approach this to speed up the writing of the flat work file created by the procedure of Figure 7. The idea is to DEFINE variables for temporary storage of field contents from the hierarchical domain(s), then use a single STORE statement to write to the flat file instead of the initial STORE followed by several MODIFYS. Alas, there are too many variables AFFINED by the procedure of Figure 12, even with the increased pool space. Nevertheless, the principle is valid and should be used where possible.

COMPILED PROGRAMS

When Datatrieve's performance proves inadequate, it may be appropriate to turn to compiled programs. Using these can result in very large savings in execution time and can overcome Datatrieve-11 pool limitations associated with reading more than two or three domains simultaneously. As an example of the improvement in execution time that can be expected, refer again to the procedure shown in Figure 7. As mentioned above, this procedure required about an hour of execution time to produce a flat, keyed, work file for the Datatrieve Report Writer. We created a FORTRAN 77 program that produced the same work file (but without keys) in 2 minutes, 15 seconds, using the indexed sequential access method on the CWR domain. This corresponded to about a factor of 25 improvement in execution time. This particular program was linked to an RMS supervisor mode library under RSX11M+, V2.1. We have experienced similar improvements with overlaid RMS libraries under RSX11M, V4.0.

A major disadvantage of writing compiled programs to access data files is that file organization must be included in each program's source code, rather than be generally available, as in a Datatrieve dictionary. One way to minimize the labor required to maintain several programs in the face of file organization changes is to create a single source code file that contains a

description of the data file's organization. For the FORTRAN language, this might include a set of TYPE declaration statements, DIMENSION statements, and FORMAT statements for use with READ and WRITE statements. This source code file can then be appended to each of the compiled programs that access the data file. If the source code is written in FORTRAN, the INCLUDE statement is a convenient tool for appending the data file's organization description to the main program and its subroutines.

SORT-11

The Sort-11 software tool has proved especially useful in preparing large reports from Datatrieve-11 data bases. One method of creating large, sorted reports is to prepare a flat work file either with Datatrieve or, if necessary, a compiled program. Then use Sort 11 to order the records appropriately before starting Datatrieve's Report Writer. In this way one can create sorted reports that would exhaust the Datatrieve pool. Furthermore, Sort-11 is extremely fast, performing sorts in 1-2 min that would take Datatrieve nearly an hour, if it could do the job at all.

CONCLUSION

The data base definitions discussed in this paper push the capabilities of Datatrieve-11 close to their practical limits of performance. For more complex linkage among records and fields, a more complete data base management system should be considered.

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```

IDU11270.31PRELOGRRM.CMD
DELETE PRELOGRRM;
DEFINE PROCEDURE PRELOGRRM
FINISH
DEFINE FILE FOR PRELOG SUPERSEDE, KEY = CWR-NO;
READY PRELOG WRITE
READY CWR SHARED READ
DECLARE CWR-TMP PIC IS 9(16) USAGE IS INTEGER;
DECLARE REC-TMP PIC IS 9(16) USAGE IS INTEGER;
DECLARE INPUT-TMP USAGE IS DATE;
DECLARE LOG-TMP USAGE IS DATE;
DECLARE SITE-TMP PIC IS X(12);
DECLARE PROJECT-TMP PIC IS X(16);
DECLARE MEMO-TMP USAGE IS DATE;
DECLARE ASSIGNED-TMP USAGE IS DATE;
DECLARE COMP-TMP USAGE IS DATE;
DECLARE BY-TMP PIC IS X(12);
DECLARE CUT-STATUS-TMP PIC IS X(16);
DECLARE FIRST-TMP PIC IS X(16);
DECLARE LAST-TMP PIC IS X(16);
DECLARE ORG-TMP PIC IS X(12);
DECLARE MAIL-TMP PIC IS X(16);
DECLARE PHONE-TMP PIC IS X(16);
DECLARE NEW-PORT-TMP PIC IS X(16);
DECLARE NEW-PART-TMP PIC IS X(16);
DECLARE NEW-STATUS-TMP PIC IS X(16);
DECLARE NEW-SPEED-TMP PIC IS X(16);
DECLARE NEW-TYPE-TMP PIC IS X(16);
DECLARE TO-DEV-TYP-TMP PIC IS X(12);
DECLARE TO-DEV-NO-TMP PIC IS X(16);
DECLARE OLD-PORT-TMP PIC IS X(16);
DECLARE OLD-PART-TMP PIC IS X(16);
DECLARE OLD-STATUS-TMP PIC IS X(16);
DECLARE OLD-SPEED-TMP PIC IS X(16);
DECLARE OLD-TYPE-TMP PIC IS X(16);
DECLARE FROM-DEV-TYP-TMP PIC IS X(12);
DECLARE FROM-DEV-NO-TMP PIC IS X(16);
DECLARE TO-TA-TMP PIC IS X(16);
DECLARE TO-BLDG-TMP PIC IS X(16);
DECLARE TO-ROOM-TMP PIC IS X(16);
DECLARE TO-BOX-TMP PIC IS X(16);
DECLARE FROM-TA-TMP PIC IS X(16);
DECLARE FROM-BLDG-TMP PIC IS X(16);
DECLARE FROM-ROOM-TMP PIC IS X(16);
DECLARE FROM-BOX-TMP PIC IS X(16);
FOR AA IN CWR WITH REC-TYP=1 AND SITE NE 'EP' BEGIN
  CWR-TMP = AA.CWR-NO
  REC-TMP = AA.REC-TYP
  INPUT-TMP = AA.INPUT-DATE
  LOG-TMP = AA.LOG-DATE
  SITE-TMP = AA.SITE
  PROJECT-TMP = AA.PROJECT-NO
  MEMO-TMP = AA.MEMO-DATE
  FIRST-TMP = A.FIRST-NAME
  LAST-TMP = A.LAST-NAME
  ORG-TMP = A.DIV1A.GRP
  MAIL-TMP = A.MAIL-STOP
  PHONE-TMP = A.PHONE
END
FOR B IN CWR WITH CWR-NO=CWR-TMP AND REC-TYP=4 BEGIN
  NEW-PORT-TMP = B.PORT-NO
  NEW-PART-TMP = B.PRT-NO
  NEW-STATUS-TMP = B.PRT-STA

```

Figure 12. Procedure to transfer data from CWR to flat work file

```

NEW-SPEED-TMP = B.PPT-SPD
NEW-TYPE-TMP = B.PRT-TYP
TO-DEV-TYP-TMP = B.DEV-TYP
TO-DEV-NO-TMP = B.DEV-NO
END
FOR C IN CWR WITH CWR-NO=CWR-TMP AND REC-TYP=5 BEGIN
  OLD-PORT-TMP = C.PRT-NO
  OLD-PART-TMP = C.PRT-STA
  OLD-STATUS-TMP = C.PRT-STA
  OLD-SPEED-TMP = C.PPT-SPD
  OLD-TYPE-TMP = C.PRT-TYP
  FROM-DEV-TYP-TMP = C.DEV-TYP
  FROM-DEV-NO-TMP = C.DEV-NO
END
FOR D IN CWR WITH CWR-NO=CWR-TMP AND REC-TYP=6 BEGIN
  FROM-TA-TMP = D.TA
  FROM-BLDG-TMP = D.BLDG
  FROM-ROOM-TMP = D.ROOM
  FROM-BOX-TMP = D.TRM-RCF
END
FOR E IN CWR WITH CWR-NO=CWR-TMP AND REC-TYP=7 BEGIN
  TO-TA-TMP = E.TA
  TO-BLDG-TMP = E.BLDG
  TO-ROOM-TMP = E.ROOM
  TO-BOX-TMP = E.TRM-RCF
END
FOR F IN CWR WITH CWR-NO=CWR-TMP AND REC-TYP=8 BEGIN
  ASSIGNED-TMP = F.ASSIGNED-DATE
  COMP-TMP = F.COMPLETION-DATE
  BY-TMP = F.CUT-BY
  CUT-STATUS-TMP = F.CUT-STATUS
END
FOR G IN CWR WITH CWR-NO=CWR-TMP AND REC-TYP=11 BEGIN
  ASSIGNED-TMP = G.ASSIGNED-DATE
  CUT-TMP = G.COMPLETION-DATE
  BY-TMP = G.CUT-BY
  CUT-STATUS-TMP = G.CUT-STATUS
END
STORE Z IN PRELOG USING BEGIN
  Z.CWR-NO = CWR-TMP
  Z.REC-TYP = REC-TMP
  Z.INPUT-DATE = INPUT-TMP
  Z.LOG-DATE = LOG-TMP
  Z.SITE = SITE-TMP
  Z.PROJECT-NO = PROJECT-TMP
  Z.MEMO-DATE = MEMO-TMP
  Z.FIRST = FIRST-TMP
  Z.LAST = LAST-TMP
  Z.ORG = ORG-TMP
  Z.MAIL-STOP = MAIL-TMP
  Z.PHONE = PHONE-TMP
  Z.NEW-PORT-NO = NEW-PORT-TMP
  Z.NEW-PART = NEW-PART-TMP
  Z.NEW-STATUS = NEW-STATUS-TMP
  Z.NEW-SPEED = NEW-SPEED-TMP
  Z.NEW-TYPE = NEW-TYPE-TMP
  Z.TO-DEV-TYPE = TO-DEV-TYP-TMP
  Z.TO-DEV-NO = TO-DEV-NO-TMP
  Z.OLD-PORT-NO = OLD-PORT-TMP
  Z.OLD-PART = OLD-PART-TMP
  Z.OLD-STATUS = OLD-STATUS-TMP
  Z.OLD-SPEED = OLD-SPEED-TMP
  Z.OLD-TYPE = OLD-TYPE-TMP
  Z.FROM-DEV-TYPE = FROM-DEV-TYP-TMP
  Z.FROM-DEV-NO = FROM-DEV-NO-TMP
  Z.TO-TA = TO-TA-TMP
  Z.TO-BLDG = TO-BLDG-TMP
  Z.TO-ROOM = TO-ROOM-TMP
  Z.TO-BOX-NO = TO-BOX-TMP
  Z.FROM-TA = FROM-TA-TMP
  Z.FROM-BLDG = FROM-BLDG-TMP
  Z.FROM-ROOM = FROM-ROOM-TMP
  Z.ASSIGNED-DATE = ASSIGNED-TMP
  Z.COMPLETION-DATE = COMP-TMP
  Z.BY = BY-TMP
  Z.STATUS = CUT-STATUS-TMP
END

```

Figure 12. cont.