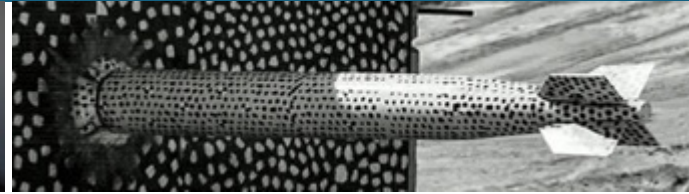
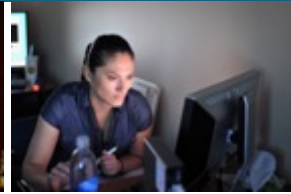




Parallel Memory-Efficient Computation of Symmetric Higher-Order Joint Moment Tensors



Zitong Li

liz20@wfu.edu

(Wake Forest University)

Hemanth Kolla

hnkolla@sandia.gov

(Sandia National Laboratories)

Eric Phipps

etphipp@sandia.gov



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Outline



1. Higher-order joint moments and cumulants.
2. Computational aspects and scientific data.
3. Our algorithmic contributions.
4. HPC implementations.
5. Results.

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Higher-order joint statistical moments



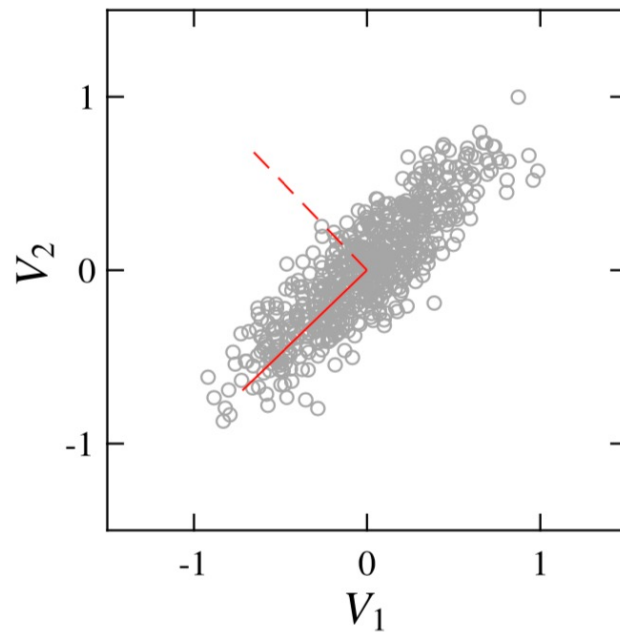
- Scientific analyses rarely look at marginal/joint moments higher than 2nd (variance/covariance).
 - Financial modelling has used coskewness and cokurtosis.
- For multi-variate non-Gaussian statistics important information is present in higher joint moments.
- Definition: For a vector of random variables, $[\tilde{x}_1, \tilde{x}_1, \dots, \tilde{x}_c]$, centered around mean i.e. $\mathbb{E}(\tilde{x}_1) = 0$:
 - Moments: $\tilde{m}_{i,j} = \mathbb{E}(\tilde{x}_i \tilde{x}_j)$, $\tilde{m}_{i,j,k} = \mathbb{E}(\tilde{x}_i \tilde{x}_j \tilde{x}_k)$, $\tilde{m}_{i,j,k,l} = \mathbb{E}(\tilde{x}_i \tilde{x}_j \tilde{x}_k \tilde{x}_l)$ where $i, j, k, l \in \{1, \dots, c\}$
 - Cumulants: $q_{i,j} = \tilde{m}_{i,j}$, $q_{i,j,k} = \tilde{m}_{i,j,k}$, $q_{i,j,k,l} = \tilde{m}_{i,j,k,l} - \tilde{m}_{i,j}\tilde{m}_{k,l} - \tilde{m}_{i,k}\tilde{m}_{j,l} - \tilde{m}_{i,l}\tilde{m}_{j,k}$
- If the random variables are joint-Gaussian, all cumulants of order > 2 are zero.
- A d^{th} -order moment/cumulant is a 'supersymmetric' tensor of order- d :
 - e.g. for a 3rd-order moment, $\tilde{m}_{i,j,k} = \tilde{m}_{i,k,j} = \tilde{m}_{j,i,k} = \tilde{m}_{j,k,i} = \tilde{m}_{k,i,j} = \tilde{m}_{k,j,i}$ for any i, j, k .

Applications of higher-order joint moments



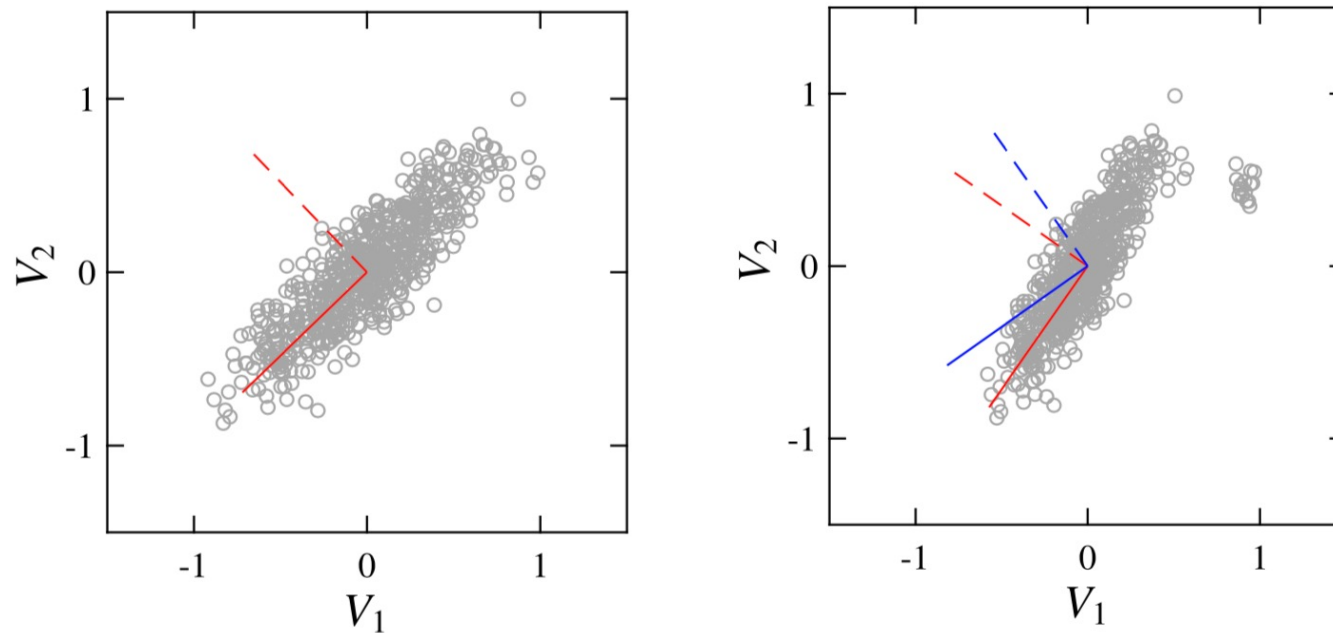
- Hitherto mostly used in financial modelling; portfolio risk assessment and asset pricing.
- Independent Component Analysis (ICA) algorithms based on eigen decomposition of 4th-order cumulants (Cardoso 1989, Comon & Cardoso 1990).
- ICA has been used for assessment of climate models (Fodor & Kamath 2003), source identification in stream water temperatures (Middleton *et al.*, 2015).
- Hyperspectral imaging: band selection and small target detection (Geng *et al.*, 2015, Głomb *et al.*, 2018), ICA-based dimensionality reduction (Wang & Chang 2006).
- Medical electrodiagnostics: artifact detection in EEG (Delorme *et al.*, 2007), feature identification in EMG (Domino *et al.*, 2019), feature extraction and classification in ECG (Yu & Chou 2008, Kutlu & Kuntalp 2012).
- Anomaly detection in multi-variate data (Peña & Prieto 2001, Konduri *et al.*, 2019).

For non-Gaussian multi-variate statistical processes higher-order joint moments are informative
(co-skewness is 3rd-order tensor, co-kurtosis is 4th-order tensor)



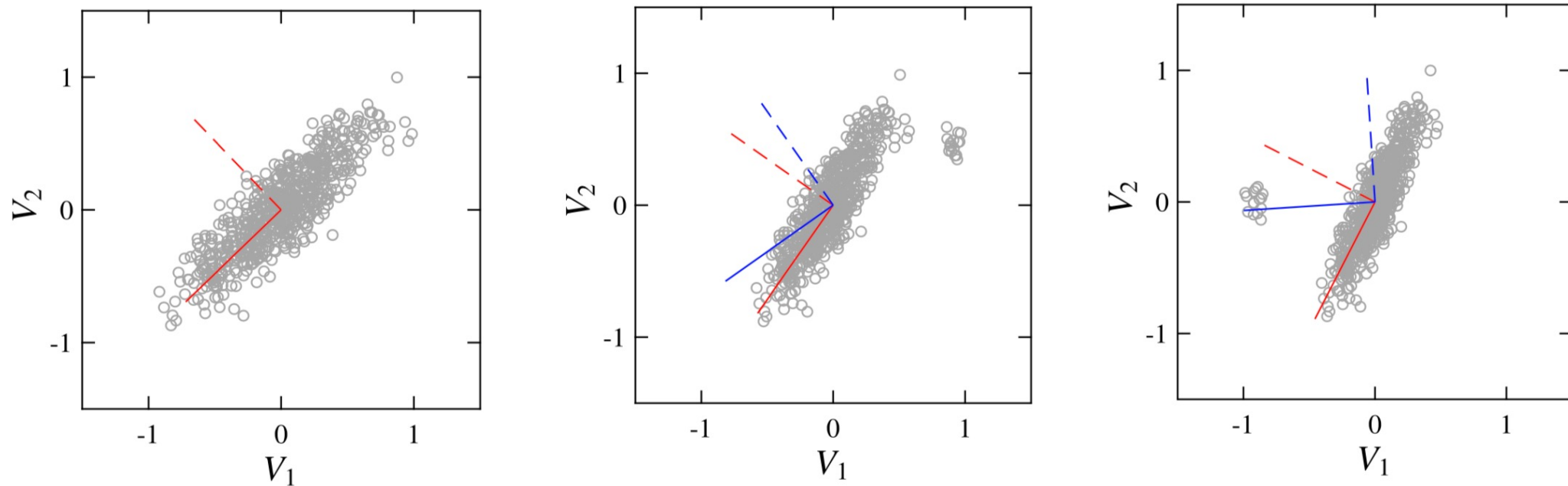
Red: Eigenvectors of Covariance (Principal Component Analysis). Denote directions of maximal variance

For non-Gaussian multi-variate statistical processes higher-order joint moments are informative
(co-skewness is 3rd-order tensor, co-kurtosis is 4th-order tensor)



Red: Eigenvectors of Covariance (Principal Component Analysis). Denote directions of maximal variance
Blue: 'Principal Kurtosis Vectors'. Obtained through HOSVD of co-kurtosis tensor.

PCA vectors not sensitive to outliers, Principal Kurtosis Vectors are.



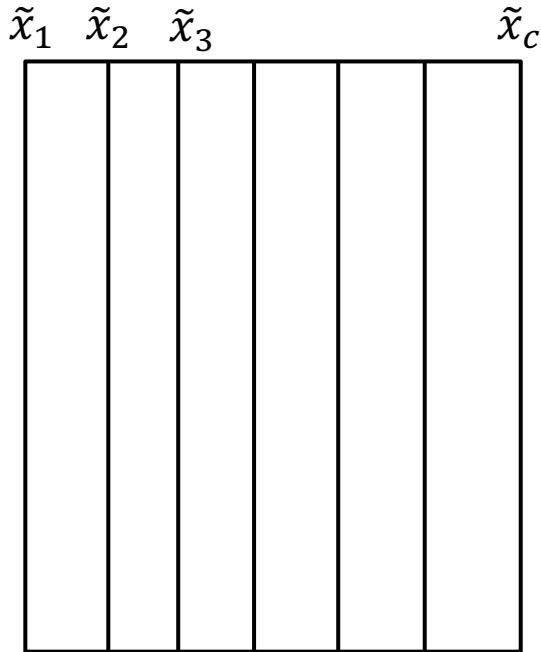
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Naïve computation of moment tensor



$X \in \mathbb{R}^{r \times c}$



```
for  $i_1 = 1:c$ 
  for  $i_2 = 1:c$ 
    .....
    for  $i_d = 1:c$ 
      for row = 1:r
         $m_{i_1, i_2, \dots, i_d} += X(\text{row}, i_1) * \dots * X(\text{row}, i_d)$ 
```

- Input matrix, r – grid points, time steps
- Typically $r \gg c$
- Naïve computation of d^{th} -order moment .
- Computational complexity $\sim \mathcal{O}(rdc^d)$.

Leveraging symmetry



- Symmetry: Full moment tensor has c^d elements, but many are duplicated
 - Number of unique elements: $\binom{c+d-1}{d}$
- Blocked Compact Symmetric Storage (BCSS) (Schatz et al., 2014).
 - Number of blocks to compute: $\binom{(c/s)+d-1}{d}$
 - Number of elements per block: s^d
 - Number of elements to be computed: $s^d \binom{(c/s)+d-1}{d}$
 - Potential savings $\sim \mathcal{O}(d!)$

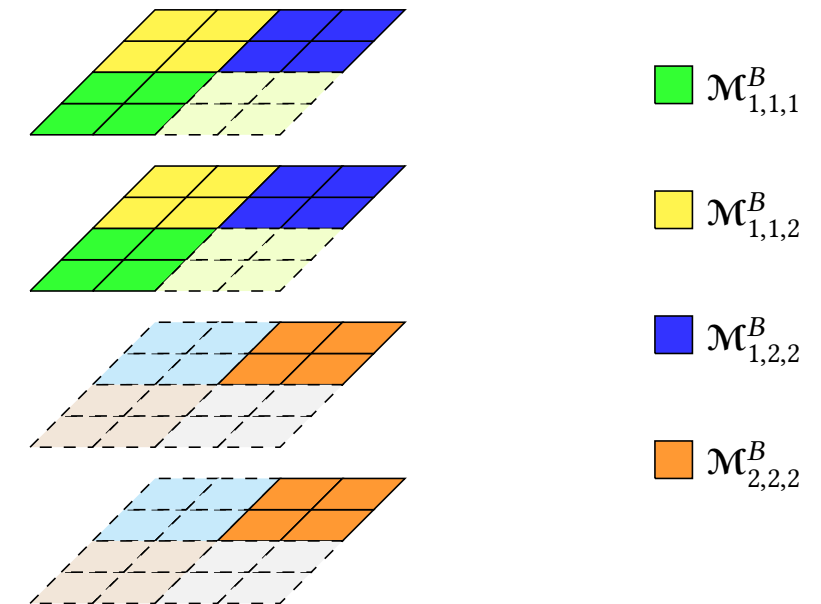
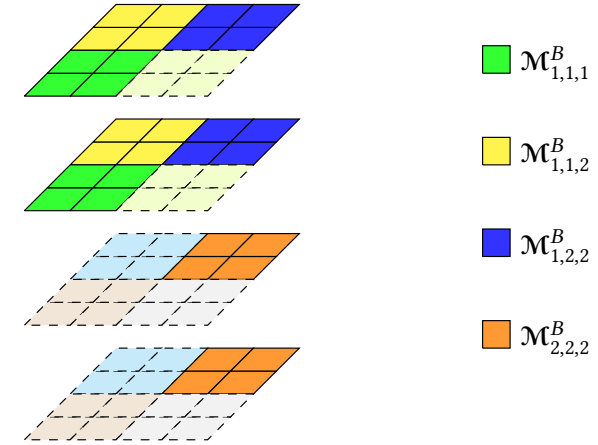


Fig. A symmetric 4x4x4 tensor divided into 8 blocks of 2x2x2 as per the BCSS format



- Domino *et al.*, (2018) leverage symmetry and BCSS to compute only unique subset of blocks.
- Focus was on computation of cumulant tensors:
 - Presented a formula for $\mathcal{C}_d = f(\mathcal{M}_d, \mathcal{C}_2, \dots, \mathcal{C}_{d-2})$
 - Involves sum of outer-products of $\mathcal{C}_2, \dots, \mathcal{C}_{d-2}$
- Compute moment tensor (\mathcal{M}_d) subblocks using nested loops
- Parallelize along the row dimension
- Speedup of $\mathcal{M}_4 \sim 24$ (relative to naïve full tensor computation)
- Speedup of $\mathcal{C}_4 \sim 100$



```

for  $i_1 = 1:s$ 
  for  $i_2 = 1:s$ 
    .....
    for  $i_d = 1:s$ 
      for row = 1:r
         $m_{i_1,i_2,\dots,i_d} += X(\text{row}, i_1) * \dots * X(\text{row}, i_d)$ 
    
```



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Refactoring of moment tensor block computation



- We propose refactoring using following ingredients:
 - $x_j^{\circ d} = \underbrace{x_j \circ x_j \circ \dots \circ x_j}_{d\text{-way outer product}}$
 - Matricisation of a d-way symmetric tensor: $Mat_p : \mathbb{R}^{(c \times c \times \dots \times c)} \rightarrow \mathbb{R}^{c^p \times c^{d-p}}$
 - Matrix Khatri-Rao product \odot - Kronecker product of individual columns of two matrices
- Motivation: \mathcal{M}_d can be expressed as sum of outer products of row vectors of X (Sherman & Kolda 2020):
 - $\mathcal{M}_d = \frac{1}{r} \sum_{j=1}^r x_j^{\circ d}$, where $x_j = X[j, :]$,
 - Can be easily pictured for covariance: $\mathcal{M}_2 = \frac{1}{r} (X^T X)$
- Final result (proof & details in paper): $Mat_p(\mathcal{M}) = \frac{1}{r} \left(\overset{p}{\bigodot} X^T \right) \left(\overset{d-p}{\bigodot} X^T \right)^T$
- Recommendation: $p = d/2$

Refactoring of moment tensor block computation



- Refactored expression applies identically to sub moment tensor sub blocks: $Mat_p(\mathcal{M}) = \frac{1}{r} \left(\bigodot^p X^T \right) \left(\bigodot^{d-p} X^T \right)^T$
- Input would be column slices of X instead of the entire matrix

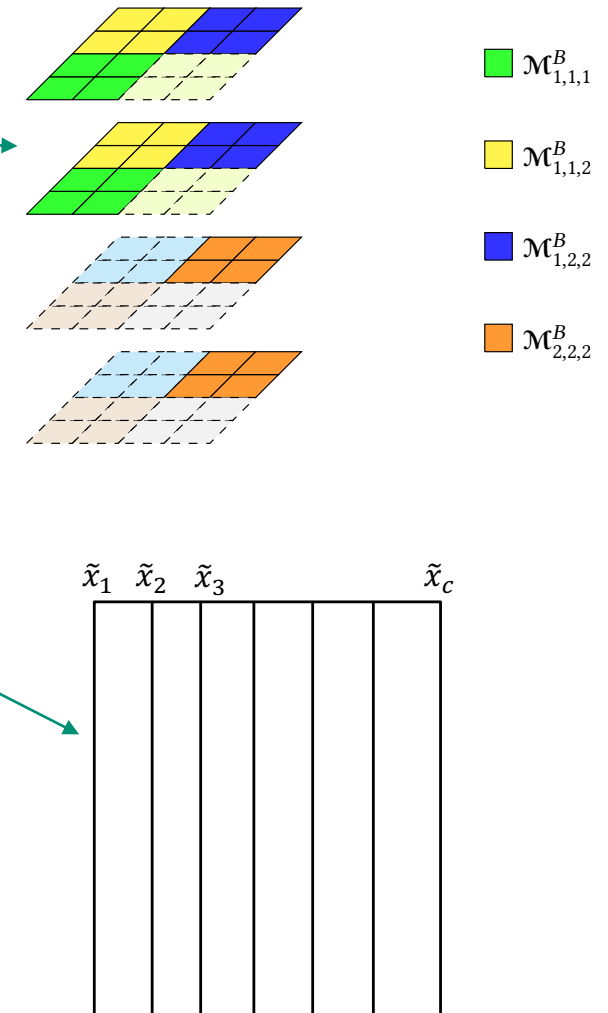
- Computational complexity:
$$\overbrace{(2rs^{d/2} + (2r-1)s^d)}^{\text{Khatri-Rao products}} \overbrace{\binom{(c/s)+d-1}{d}}^{\text{\# of blocks}}.$$

- Less expensive than naïve (nested for-loops) by factor $\approx d/2$.
- Cache complexity: Matrix Khatri-Rao products are very cache-friendly, leading to fewer cache misses:
 - We recommend X to be row-major order. Matrix transpose (if need be) is $\sim \mathcal{O}(\frac{rc}{L})$ cache misses.
 - Refactoring incurs fewer cache misses by at least $\sim \mathcal{O}(s^{d/2})$ (various scenarios detailed in paper).

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Hierarchical parallelism

- Each subblock of \mathcal{M}_d can be computed independently:
 - For distributed-memory this could be along nodes/processes/ranks.
 - For shared-memory this can be along thread parallel units (e.g. GPU warps).
- Traversal along row-dimension of X can be parallel-reduced:
 - For distributed-memory this is aligned with domain decomposition.
 - For shared-memory this can be tiling (to reduce memory requirements).
- The matrix Khatri-Rao product (due to refactoring) exposes another level of parallelism.



Performance portable implementation with Kokkos



- Kokkos – a C++ library for performance portable implementations.
- Detailed semantics to express data and compute parallelism. Hardware details of memory layout and parallelism units are abstracted out.
- We provide Kokkos implementations with three levels of parallelism:
 - “Thread Team” parallelism (e.g. GPU warp)
 - Tile parallelism (row dimension of X)
 - Thread-level parallelism

Algorithm 3 Parallel Algorithm for computing the 4th moment tensor

```

1: function  $\mathcal{M}^{(4)} = \text{4THMOMENTTENSOR}(X, n, t, s)$ 
2:    $\triangleright s = \text{block size}, t = \text{tile size}, n = \# \text{ of teams}, X \in \mathbb{R}^{r \times c}$ 
3:    $\bar{r} = r/t$   $\triangleright \# \text{ of row tiles}$ 
4:    $nbm = \text{ceil}(c/s)$   $\triangleright \# \text{ of blocks on each mode}$ 
5:    $nb = \binom{nbm+3}{4}$   $\triangleright \text{total } \# \text{ of unique blocks in } \mathcal{M}^{(4)}$ 
6:    $\bar{nb} = nb/n$   $\triangleright \# \text{ of blocks each team computes}$ 
7:   teams_parallel_for( $i = 1, \dots, n$ )
8:     [ $b_s = (i-1)\bar{nb} + 1, b_e = i\bar{nb}$ ]  $\triangleright \text{blocks of this team}$ 
9:     for  $j = b_s, \dots, b_e$  do
10:       $\mathbb{T}(j; nbm) \mapsto (i_1, i_2, i_3, i_4)$   $\triangleright \text{multi-index this block}$ 
11:       $X_{i_l} = X_{:, (i_l-1)s+1:i_ls}, i_l = (i_1, \dots, i_4)$   $\triangleright \text{column slices}$ 
12:      for  $k = 1, \dots, \bar{r}$  do
13:         $X_{k, i_l} = X_{(k-1)t+1:kt, i_l}$   $\triangleright k^{th} \text{ row tile}$ 
14:         $Y = \text{KHATRIRAOPRODUCT}(X_{k, i_2}^T, X_{k, i_i}^T)$ 
15:         $Z = \text{KHATRIRAOPRODUCT}(X_{k, i_4}^T, X_{k, i_3}^T)$ 
16:         $\mathcal{M}_{(i_1, i_2, i_3, i_4)}^B += \frac{1}{\bar{r}} YZ^T$ 

```

Algorithm 2 Parallel Algorithm for the Khatri-Rao Product

```

1: function  $C = \text{KHATRIRAOPRODUCT}(A, B)$ 
2:    $\triangleright A \in \mathbb{R}^{a \times c}, B \in \mathbb{R}^{b \times c}$ 
3:   threads_parallel_for( $i = 1, \dots, ab$ )
4:     vector_parallel_for( $j = 1, \dots, c$ )
5:        $(\alpha, \beta) = \mathbb{T}(i; a, b)$ 
6:        $R_{i,j} = X_{\alpha,j} X_{\beta,j}$ 

```

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Serial performance



- All experiments for 4th-order moment tensor, \mathcal{M}_4 .
- Serial Julia implementation compared against reference implementation (Domino *et al.*, 2018), AMD EPYC 7302.
- Parameters: r – rows of X , c – cols of X , s – size of each subblock

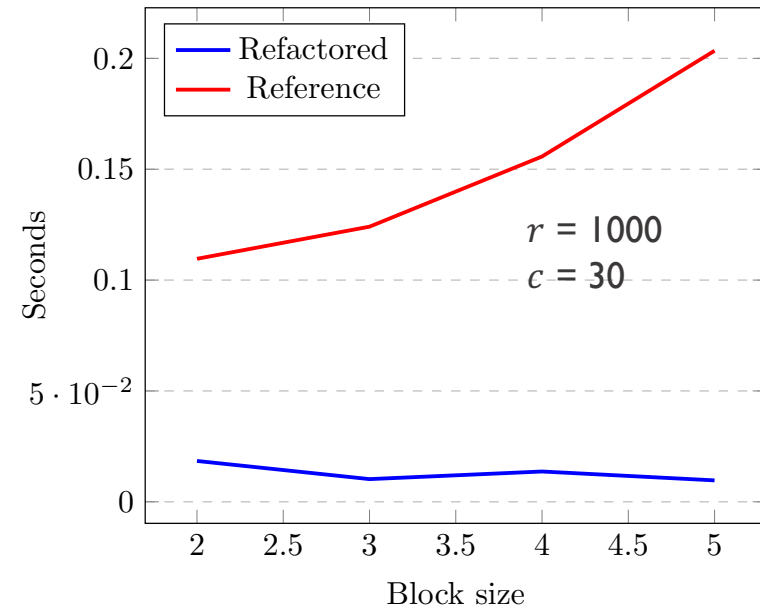


Fig. Trend with varying block size

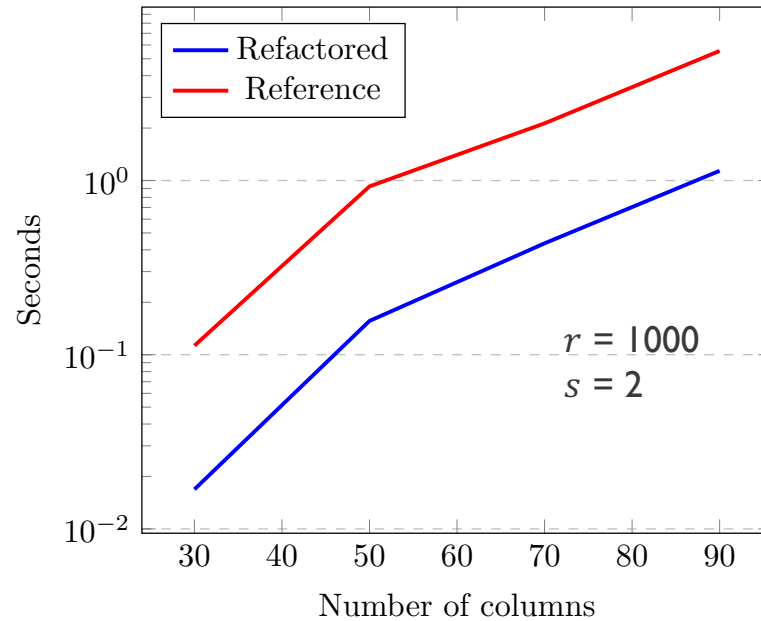


Fig. Trend with varying #columns

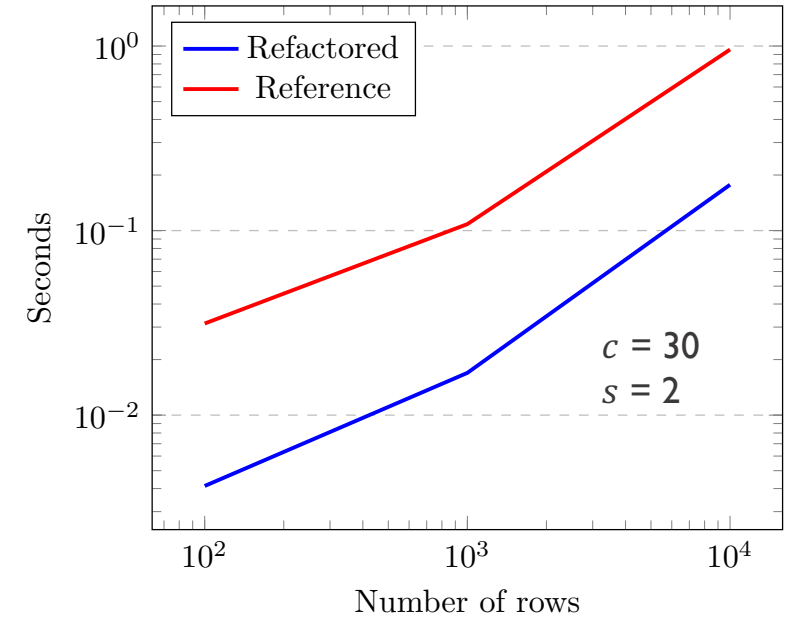


Fig. Trend with varying #rows



- All experiments for 4th-order moment tensor, \mathcal{M}_4 .
- Kokkos implementation compared against reference implementation (Domino *et al.*, 2018), NVIDIA Tesla V100.
- Parameters: r – rows of X , c – cols of X , s – size of each subblock

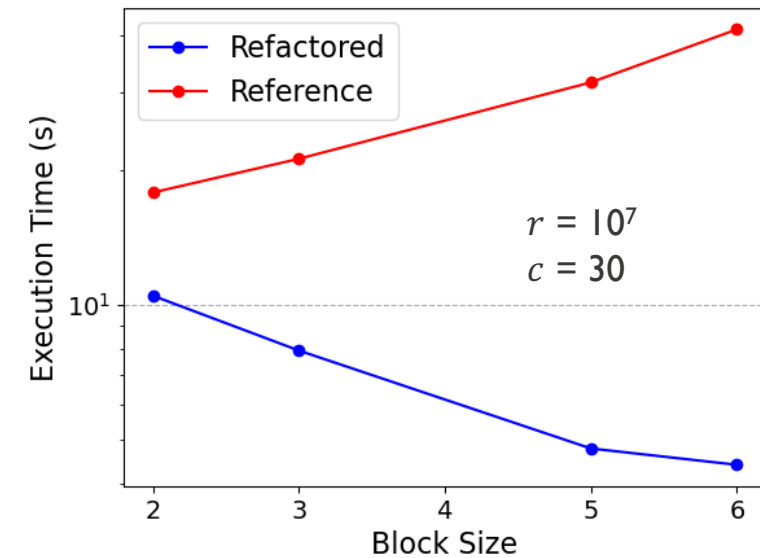


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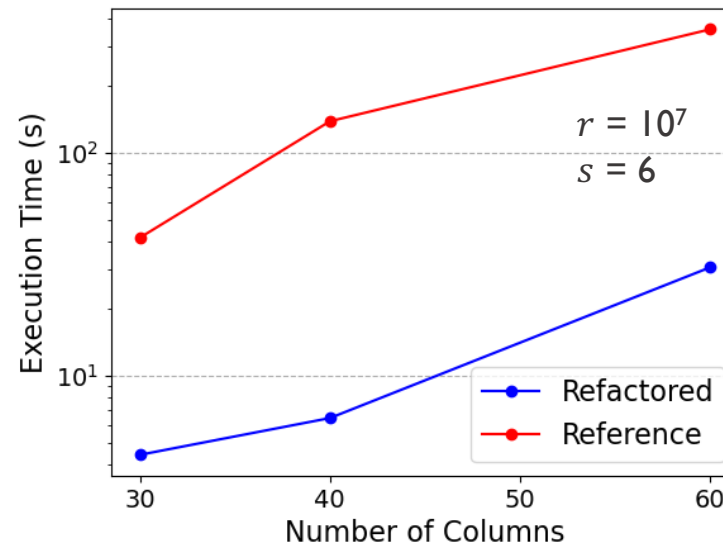


Fig. Trend with varying #columns

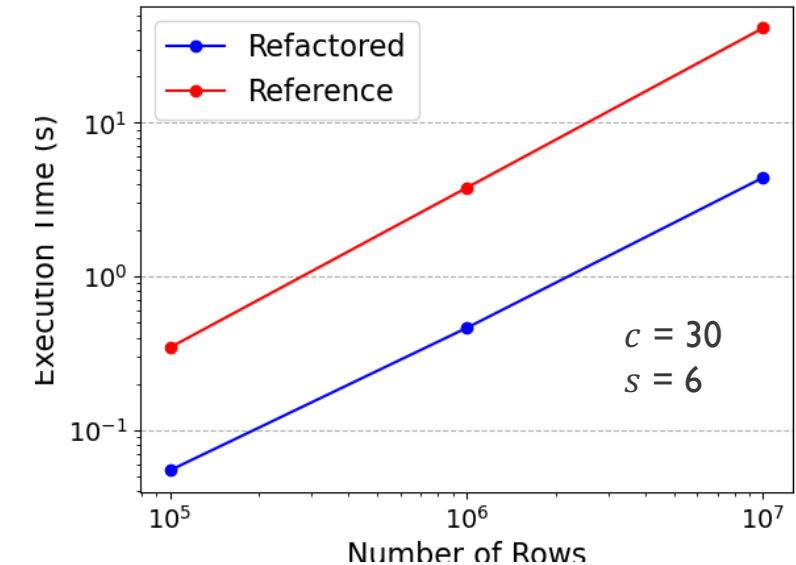


Fig. Trend with varying #rows

Cache performance



- All experiments for 4th-order moment tensor, \mathcal{M}_4 .
- Kokkos implementation compared against reference implementation (Domino *et al.*, 2018), NVIDIA Tesla V100.
- Flops and MemOps (L1-cache) measured with nvprof.

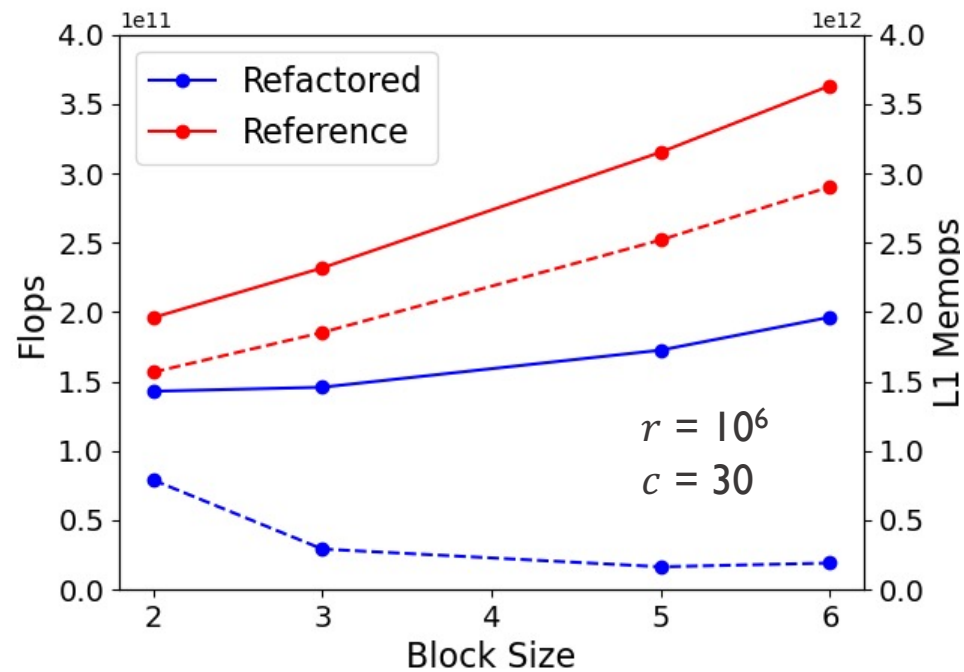


Fig. Trend with varying block size

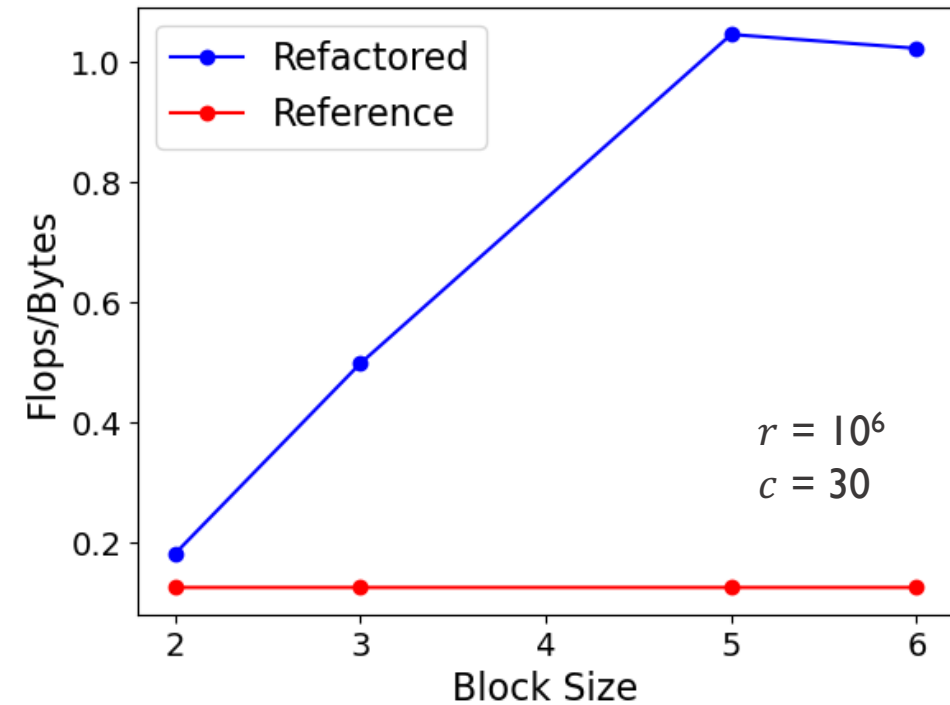


Fig. Trend with varying block size



Thank You



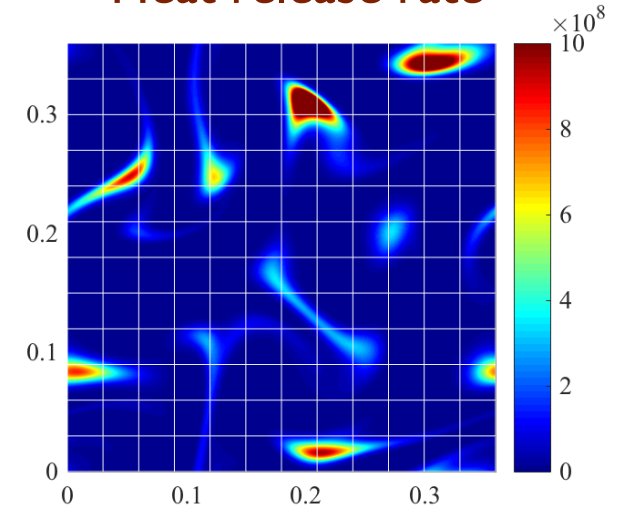
Backup

Formalizing Distributed Rare Event Detection

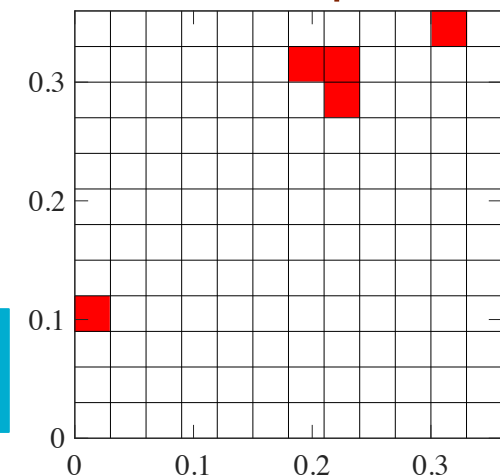


- Compute Principal Kurtosis Vectors on each data partition (e.g. processor).
- Compare the vectors amongst partitions in space and/or time:
 - Proposed Feature moment metrics (fraction of the kurtosis attributable to each variable) to quantify orientation of Kurtosis vectors.
 - FMMs sum to unity, akin to discrete distribution.
 - Divergence metric (Hellinger distance) to compare across partitions.
- Most computation (cokurtosis tensor and principal vectors) is local.
- Communication only of a small vector of numbers (FMMs).

Heat release rate



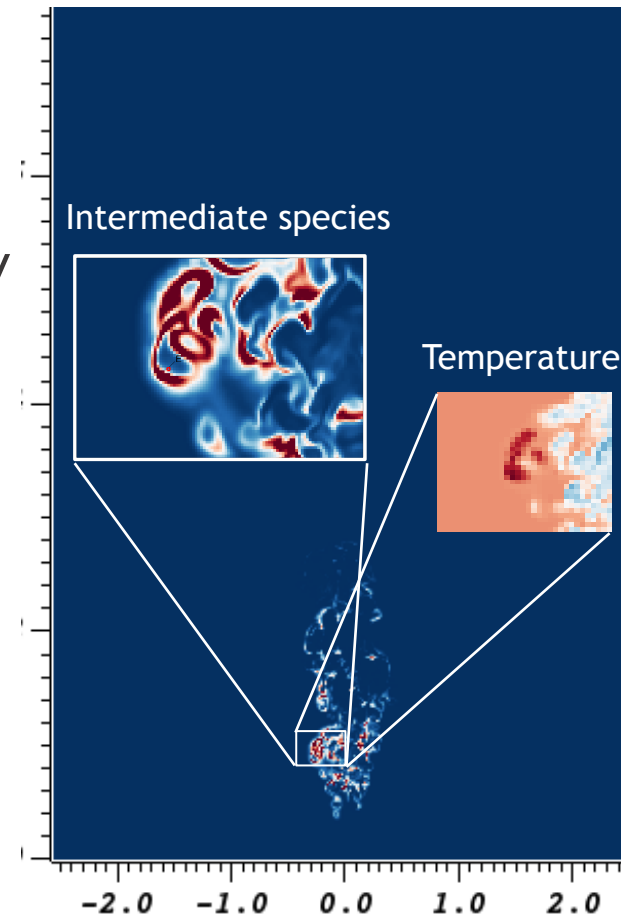
Anomalous partitions



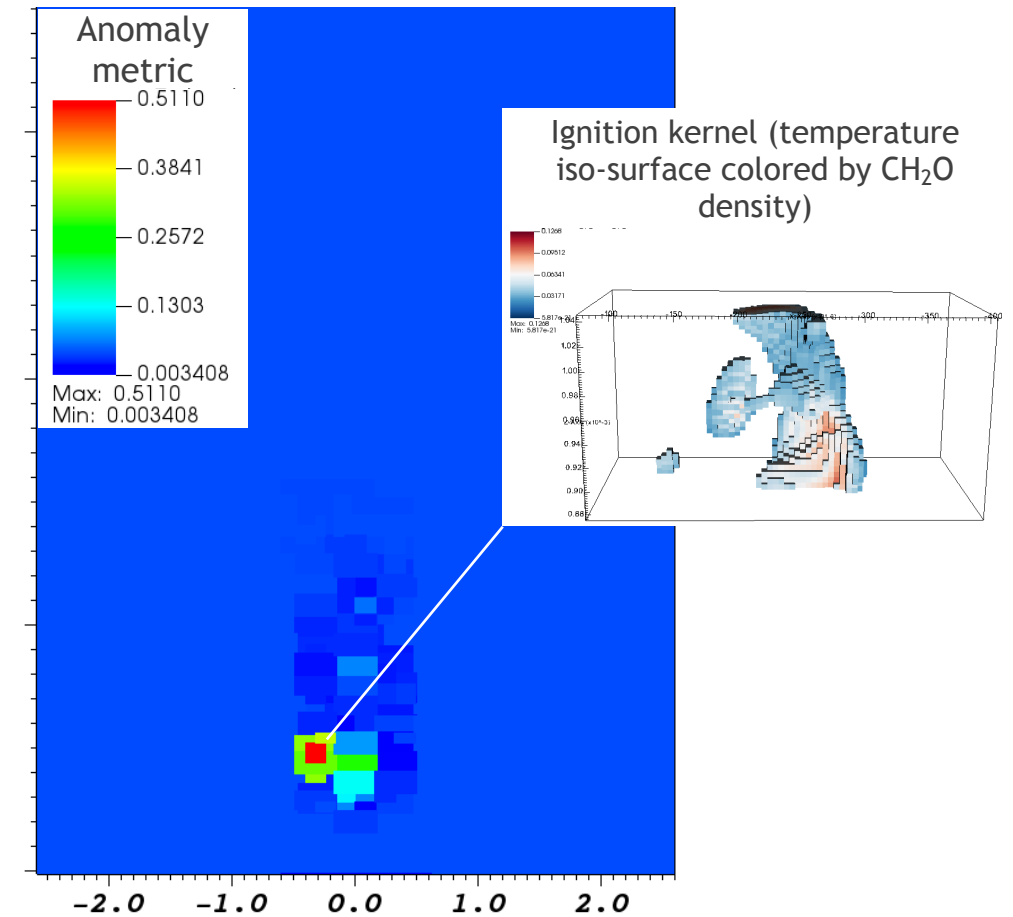
“Anomaly detection in scientific data using joint statistical moments.” K. Aditya, H. Kolla, W.P. Kegelmeyer, T.M. Shedd, J. Ling, W.L. Davis IV, Journal of Computational Physics, 2019.

Pele: PeleLM, adaptive-mesh low Mach number hydrodynamics code for reacting flows

Identification currently based on ad-hoc thresholds



Validation: co-kurtosis tensor-based unsupervised anomaly detection



Contributors: Martin Rieth, Jackie Chen, Marco Arienti, Janine Bennett (Sandia Natl. Labs), Matt Larsen (formerly at LLNL)