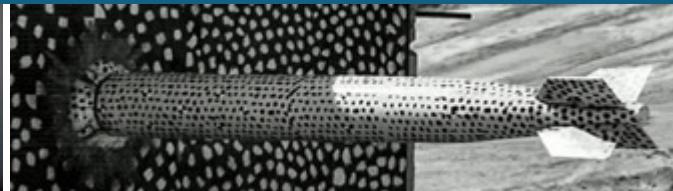
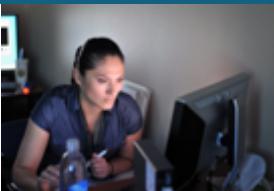




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# Scalable verification of quantum algorithm circuits



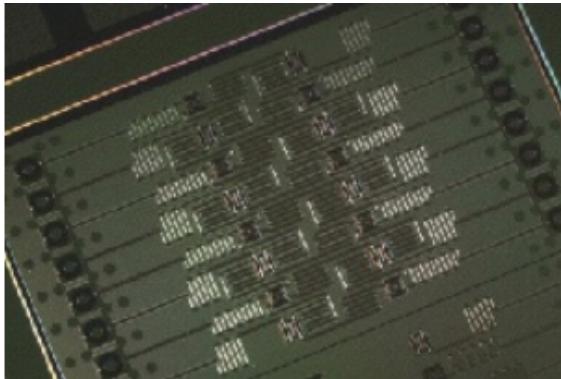
## Mohan Sarovar

Stefan Seritan, Erik Nielsen, Kenneth Rudinger, Kevin Young, Robin Blume-Kohout, **Tim Proctor**

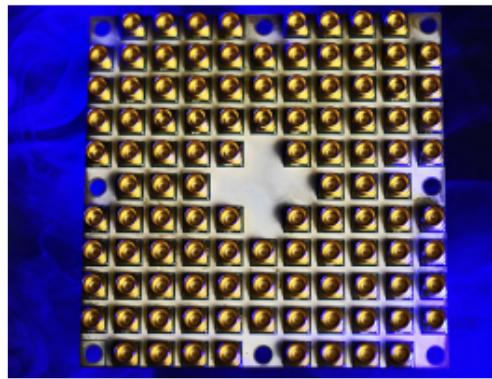
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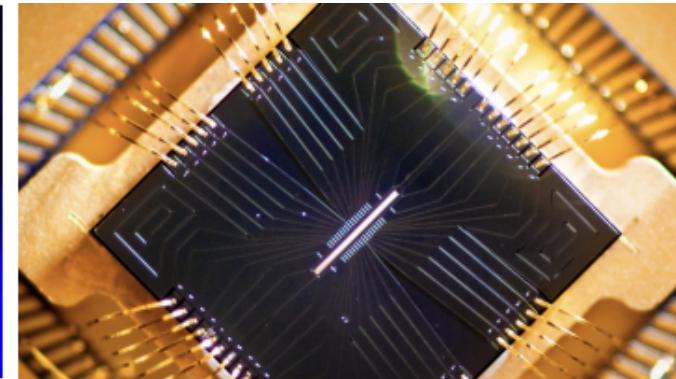
# Verification of quantum computer output



IBM



Intel



Sandia

In the next decade: 100s – 1000s of qubits + poly depth circuits

- How do we verify the correctness of the output of such systems?
  - Many interesting applications will have output that is not efficiently verifiable by classical computers
  - Microscopic modeling will be impossible

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**Honeywell says quantum computers will outpace standard verification in '18 to 24 months'**

Michael Vizard @mrvizard March 26, 2021 6:50 AM

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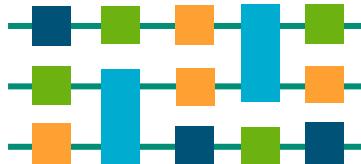
# Scalable verification



**Verification goal:** estimate fidelity of actual circuit with ideal circuit

We are given a circuit of interest  $C$ , which must contain only:

- Arbitrary single-qubit gates
- Two-qubit gates that are self-inverse and Clifford (not an essential assumption).



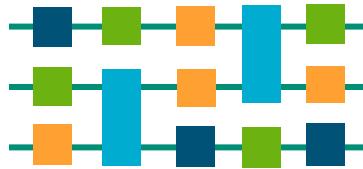
Estimate fidelity of an implementation of this circuits using data from executing circuits sampled from **three different mirror circuit [1] ensembles**

- Each ensemble is a slight variation of the circuit followed by its inverse

# The three ensembles

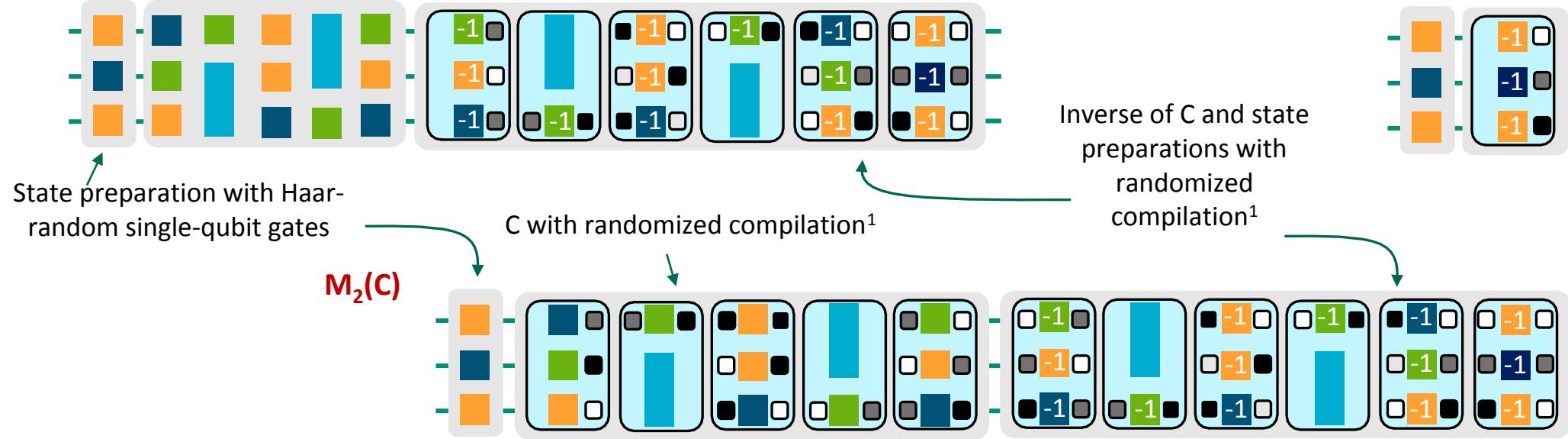


Original circuit

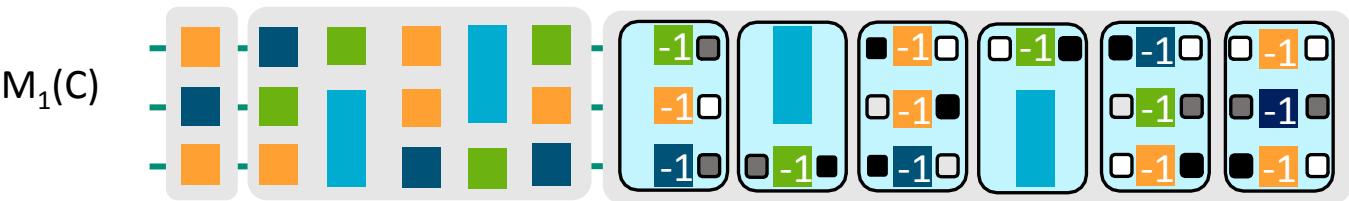


**$M_1(C)$**

The circuit  $C$

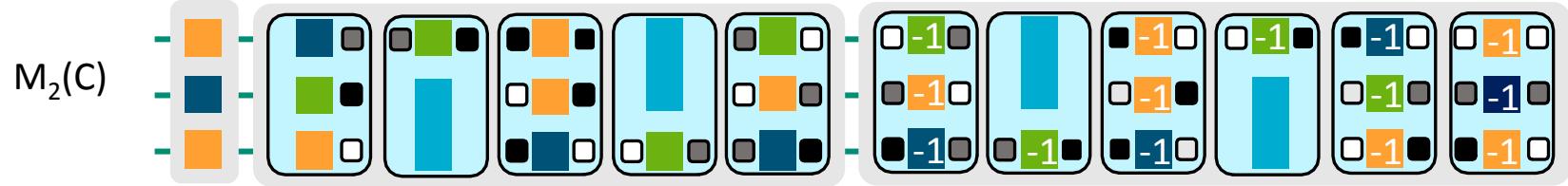


# The three ensembles

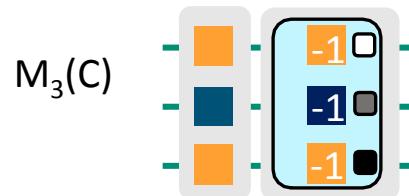


$F(C)$  = The process fidelity with which  $C$  is implemented

$$S_1 = \text{Success Prob}(M_1(C)) \approx F(\text{SPAM}) * F(C) * F(\text{RC}[C^{-1}])$$



$$S_2 = \text{Success Prob}(M_2(C)) \approx F(\text{SPAM}) * F(\text{RC}[C]) * F(\text{RC}[C^{-1}]) \approx F(\text{SPAM}) * F(\text{RC}[C^{-1}])^2$$

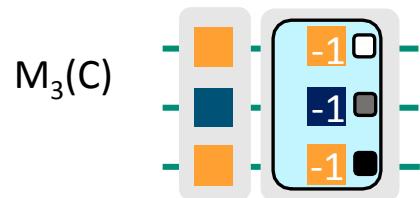
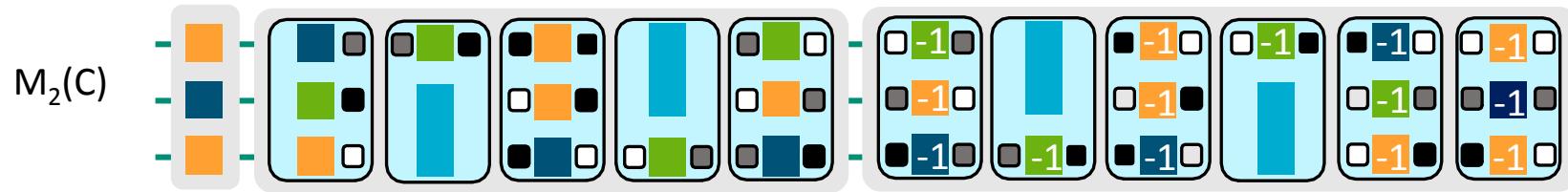
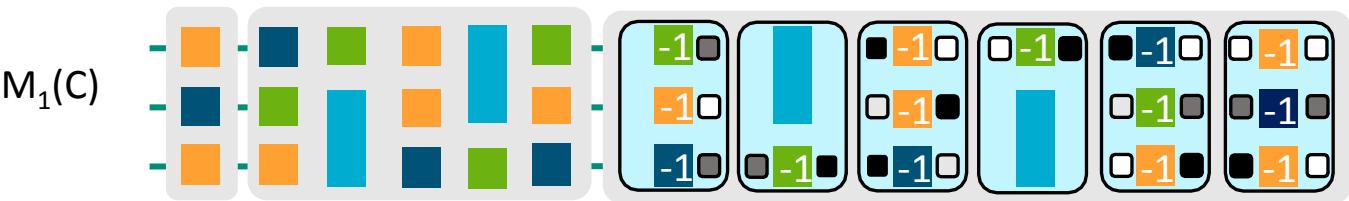


$$S_3 = \text{Success Prob}(M_3(C)) \approx F(\text{SPAM})$$

$$F(C) \approx S_1 / (S_2 \times S_3)^{1/2}$$

The actual formula we use is slightly more complex, because (process) fidelities don't multiply...

# Circuit fidelity estimate



Our estimate for  $F(C)$

$$\chi_F(c) = 1 - \frac{4^n - 1}{4^n} \left( 1 - \frac{\mathbb{E}[\gamma(M_1(c))]}{\sqrt{\mathbb{E}[\gamma(M_2(c))]\mathbb{E}[\gamma(M_3(c))]}} \right)$$

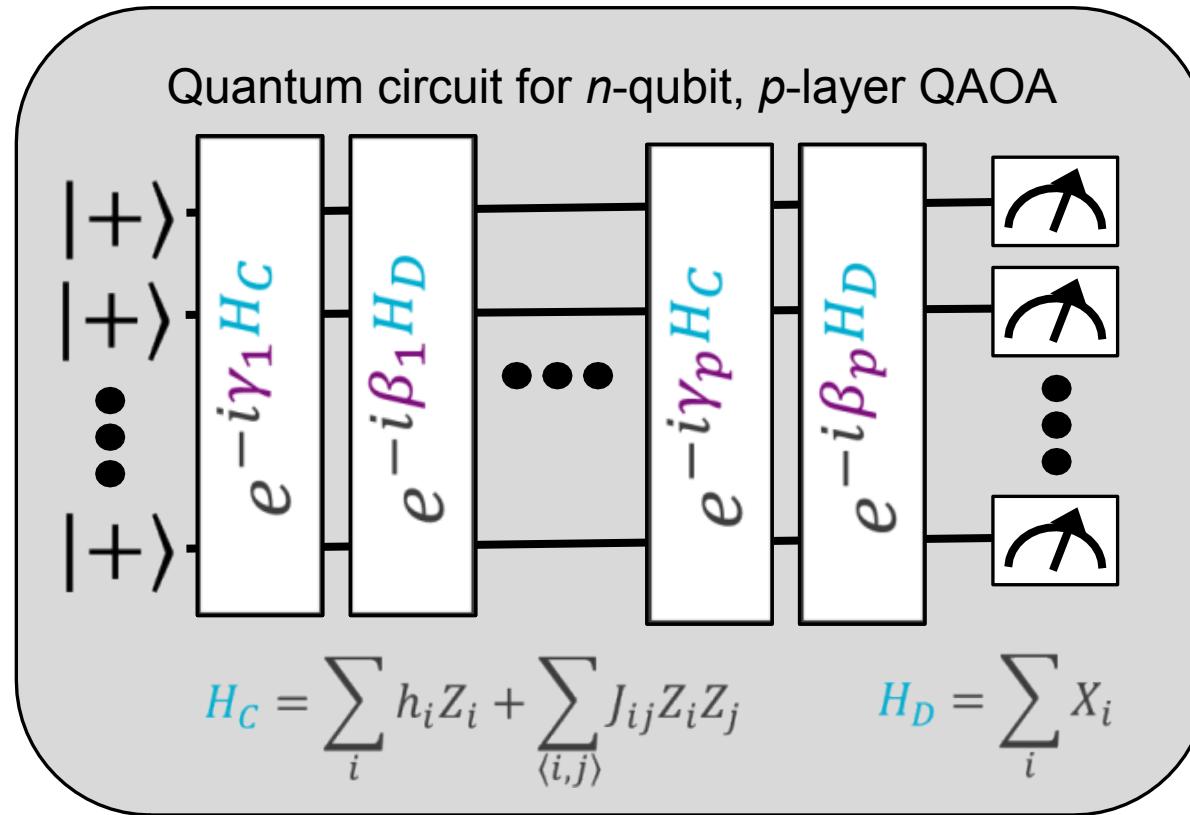
This is efficient to estimate, for any number of qubits, except when  $F(C)$  is very small!

$$\gamma = \frac{4^n}{4^n - 1} \left[ \sum_{k=0}^n \left( -\frac{1}{2} \right)^k h_k \right] - \frac{1}{4^n - 1},$$

# Illustrations



Quantum approximate algorithm (QAOA) circuits

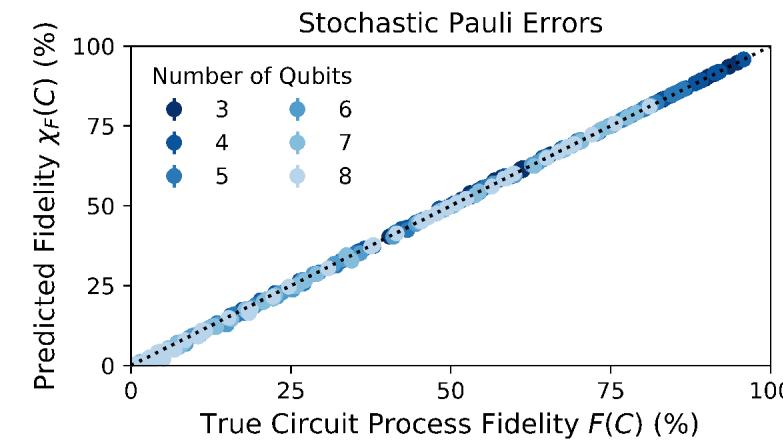


- Variable number of qubits, QAOA layers ranging from  $p=1$  to  $p=10$

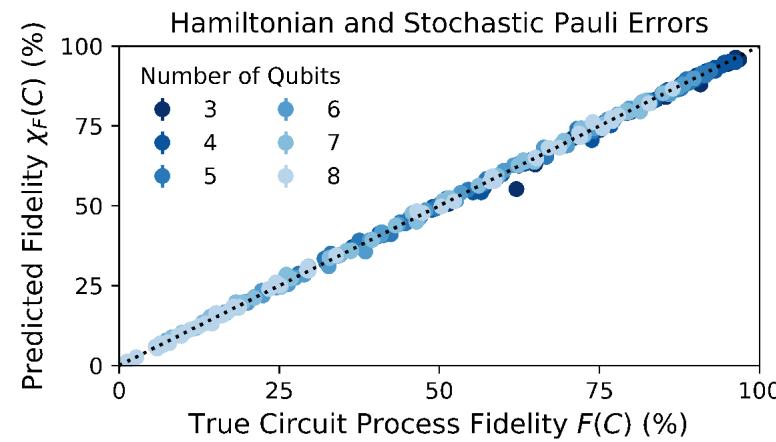
# Illustrations



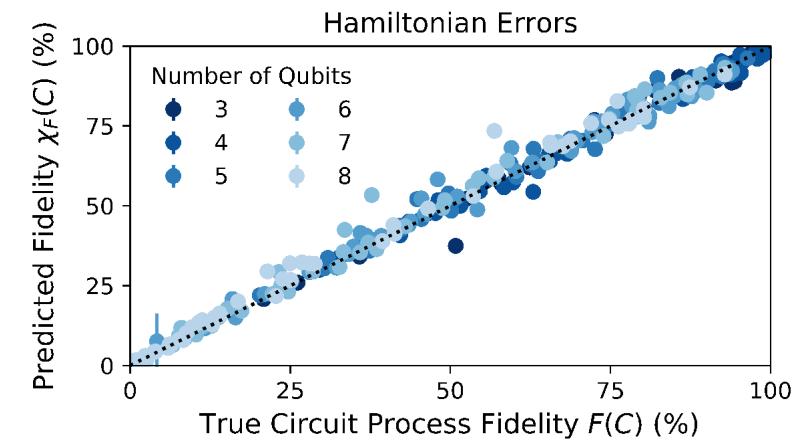
## Comparison of estimated process fidelity to actual process fidelity



Perfect correlation under stochastic noise model



Excellent correlation under noise model with equal stochastic and Hamiltonian contributions



Good correlation, with some outliers, under noise model with *only* Hamiltonian/coherent errors

- $n = 3 - 8$  qubits
- $p = 1, 2, 5, 8, 10$  QAOA layers
- For each  $(n, p)$  there are 10 problem instances (300 total circuits).
- 1000 of each type of mirror circuit for each QAOA circuit.
- Each QAOA instance (+ the 3000 corresponding mirror circuits) has its own randomly sampled error model.

# Summary



- We have developed a *scalable* way to verify accuracy of quantum circuit implementations
- Scalability:
  - No classical simulation of quantum circuits necessary
  - Requires execution of an ensemble of auxiliary circuits of roughly twice the depth
  - Analysis and sample complexity have weak dependence on *number of qubits, circuit depth*
- Solves one of the major challenges emerging in the NISQ era of quantum computing