

A Cooperative Compiler and Runtime Checkpoint/Restart Approach for Kokkos

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MOTIVATION

Performance Portability

- ❑ Write one implementation that
 - can run on multiple platforms
 - can enable suitable memory access patterns for different platforms
 - can exploit platform-specific features
- ❑ Examples:
 - Kokkos
 - RAJA
 - DPC++



Kokkos

□ What's Kokkos?

- Kokkos enables single performance portable codes
 - ✓ CPUs, NVIDIA GPUs, ...
- Kokkos provides data abstractions critical for performance portability not available in OpenMP or OpenACC (Kokkos::View)
- Essential parallel constructs
 - ✓ Kokkos::parallel_for();
 - ✓ Kokkos::parallel_reduce();
 - ✓ Kokkos::parallel_scan();



Kokkos Example

```
1  View<int*> x("x", N);           // Construct a View (length: N)
2  // Parallel For
3  parallel_for(RangePolicy<>(0, N), // 1D Range (0 to N)
4              [=] (int i) {        // C++11 lambda
5                  x(i) = i;         // Computational Body
6              });
7
8  int result = 0;
9  // Parallel Reduce
10 parallel_reduce(RangePolicy<>(0, N), // 1D Range (0 to N)
11                [=] (int i, int &update) { // C++11 lambda
12                    update += x(i);         // Computational Body
13                }, result);                // Value to update
14 std::cout << result << std::endl;
```



Enabling Resiliency in Performance Portable Programs

□ Why important?

- Increasing the number of nodes from 10,000 to 22,000 increases the probability of failures by a factor of 20x on the Blue Waters System [1]

□ Typical approach: Manual checkpointing

- e.g., for Kokkos [2]

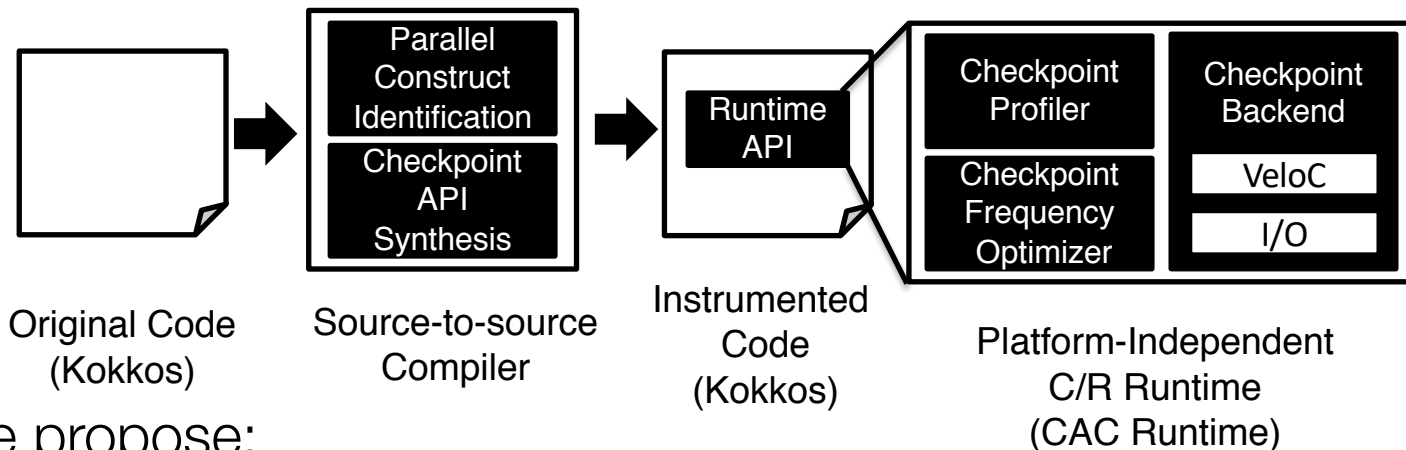
□ Research question: Any way to automate it?

[1] Martino et al. “Measuring and Understanding Extreme-Scale Application Resilience: A Field Study of 5,000,000 HPC Application Runs” (DSN’15)

[2] Morales et al. “Towards High Performance Resilience Using Performance Portable Abstractions” (Euro-Par’21)



Big picture



□ We propose:

- a source-to-source compiler that analyzes Kokkos programs and automatically generates Checkpoint/restart API
- a runtime provides
 - ✓ 1) platform-independent Checkpoint/restart API routines
 - ✓ 2) optimizes the frequency of checkpointing when possible

Cooperate



Related Work

	Language	IR	Note
Elkhouly et al.	Sequential C + directive on data access	LLVM IR	Sequential
Bronevetsky et al.	OpenMP C	ROSE	OpenMP
Lehr et al.	MPI + C	LLVM IR	Type Safe Checking Only
Shahneous et al.	OpenSHMEM	LLVM IR	Suggestion Only
Our Approach	Kokkos (C++)	Clang AST	Full translation

Note: compilers are typically used to perform selective instruction duplication (SID) e.g., Swift



Related work (Cont'd)

❑ clang-tidy for Kokkos (not for resilience)

```
$ clang-tidy kokkos-implicit-this-capture.cpp -checks=-\*,kokkos\* -- -  
I$PWD/Inputs/kokkos
```

...

```
warning: Lambda passed to parallel_for implicitly captures this. [kokkos-  
implicit-this-capture]
```

```
10, KOKKOS_LAMBDA(int x) { my_int = x; });
```

<https://github.com/kokkos/llvm-project/tree/master/>

<https://github.com/kokkos/llvm-project/tree/master/clang-tools-extra/clang-tidy/kokkos>



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DESIGN

High-level Overview (lambda)

Original Code

```
1 Kokkos::View<int*> A("A", N);
2 Kokkos::View<int*> B("B", N);
3 for (int t = 0; t < T; t++) {
4     Kokkos::parallel_for("modA",
5         N, KOKKOS_LAMBDA (int i) {
6         A(i) = B(i);
7     });
8     ...
9 }
```

Translated

```
1 Kokkos::View<int*> A("A", N);
2 Kokkos::View<int*> B("B", N);
3 loop_enter();
4 for (int t = 0; t < T; t++) {
5     loop_start(t);
6     Kokkos::parallel_for("modA",
7         N, KOKKOS_LAMBDA (int i) {
8         A(i) = B(i);
9     });
10    checkpoint(A);
11    ...
12    loop_end(t);
13 }
14 loop_exit();
```

High-level Overview (functor)

Original Code

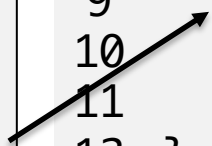
```
1 Kokkos::View<int*> A("A", N);
2 Kokkos::View<int*> B("B", N);
3 for (int t = 0; t < T; t++) {
4     Functor functor(A, B);
5     Kokkos::parallel_for("modA",
6         N, functor);
```

```
7 ...
8 }

struct Functor {
    void operator()(const size_t i)
    const {
        A[idx] = B[idx];
    }
    // synthesized by translator
    void checkpoint() { checkpoint(A); }
};
```

Translated

```
1 Kokkos::View<int*> A("A", N);
2 Kokkos::View<int*> B("B", N);
3 loop_enter();
4 for (int t = 0; t < T; t++) {
5     loop_start(t);
6     Functor functor(A, B);
7     Kokkos::parallel_for("modA",
8         N, functor);
9     functor.checkpoint(A);
10    ...
11    loop_end(t);
12 }
13 loop_exit();
```



Compiler

- ❑ Open question: How to analyze C++ program?
 - LLVM IR-based
 - ✓ Pros: Lots of analyzer and optimizer passes
 - ✓ Cons: It's very hard to locate the start/end of a specific function
 - Clang AST-based (Our approach)
 - ✓ Pros: it's relatively easy to locate the start/end of a specific function
 - ✓ Cons: its analyzers/optimizers would not be as powerful as LLVM's ones

```
// ClangAST
LambdaExpr ... '(lambda at test.cpp:55:40)'
| ...
| - CompoundStmt ...
|   BinaryOperator ... int' lvalue '='
|   | - CXXOperatorCallExpr ... 'int' lvalue '()'
|   | | ...
|   | | - DeclRefExpr ... 'const Kokkos::View<int *>' ... 'A'
```



Compiler (Cont'd)

❑ Checkpoint API insertion

1. Identify Kokkos constructs in user's program
 - ✓ Kokkos::parallel_for, reduce, scan
2. Analyze the body of a lambda expression/the parenthesis operator and identify which View is used
3. Insert checkpoint API immediately after Kokkos::parallel_for
 - ✓ Optimization: do not checkpoint a View that is read-only

❑ Loop API Insertion

- Insert loop_enter(), loop_start(), loop_end(), loop_exit() to all for-loops

Note: the current implementation performs the above actions for each lambda and it does not do any data flow analyses to analyze the lifetime of Views.



How robust is the compiler?

- ❑ Our compiler is able to translate 9 Kokkos applications including 3 real-world applications

Benchmark Suite	Benchmark
N/A	Heat-Dist, NimbleSM
Polybench MPI (stencil)	Jacobi-1D/2D, Heat-3D, FDTD-2D, Seidel-2D
Mantevo	miniAero/miniFE



CAC Runtime

- ❑ Provides a generic checkpoint/restart API independent of datatype and checkpoint runtime
 - Uses Kokkos::ViewHolderBase for de-templating Views but keeping memory information
 - ✓ Kokkos::View<int*>, Kokkos::View<double*>, ...
 - Keeps format of Resilient Kokkos [1] configuration file to enable changing checkpoint backends without application code changes
 - ✓ VeloC, StdFile, ...



CAC Runtime (Cont'd)

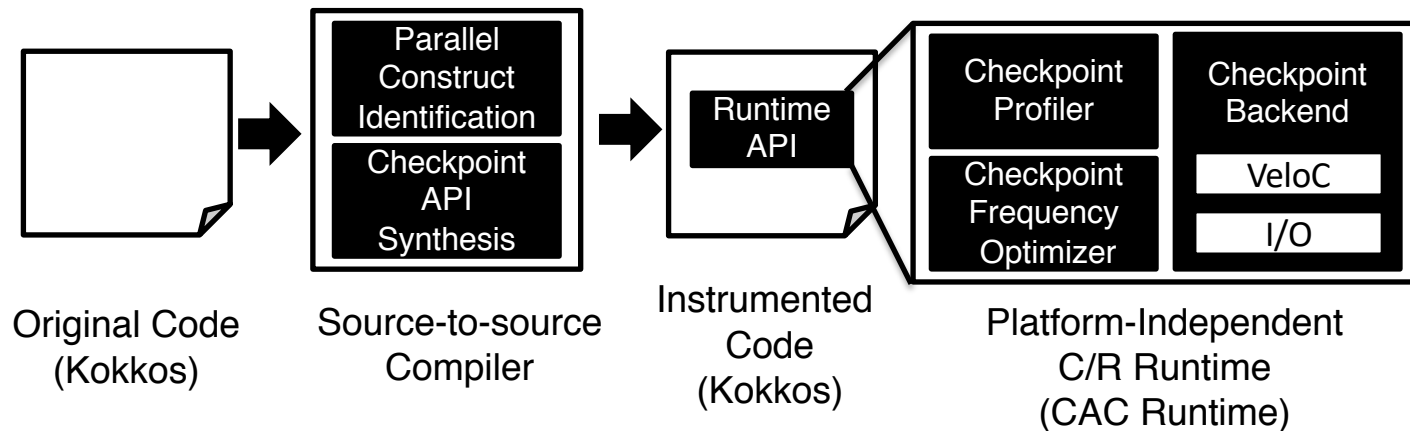
- ❑ Gathers checkpoint regions within loops for write aggregations
- ❑ Automatically updates loop iteration variables after recovery
- ❑ Dynamically optimizes checkpoint frequency
 - User provides estimated MTBF
 - Calculates time-weighted average checkpoint cost



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IMPLEMENTATION

Implementation



- ❑ Compiler: implemented as a Clang's tool (executable)
- ❑ Runtime: a standalone library which depends on
 - Kokkos-Resilience (Proprietary)
 - VeloC: <https://veloc.readthedocs.io/en/latest/>



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PRELIMINARY EVALUATIONS

Applications

❑ Platform:

- Sandia Machine: Kahuna
 - ✓ a dual socket intel E5-2683v3 2.00GHz CPU (28 cores)
 - ✓ High memory data-analytics machine (256GB)
 - ✓ 56GB/s FDR Infiniband network
 - ✓ CephFS distributed filesystem

❑ Application: Heat-distribution (MPI+Kokkos)

- 4-points iterative stencil kernel
- Variants:
 - ✓ Raw VeloC
 - ✓ Auto-translated by the compiler + CACRuntime



Heat-distribution

□ Parameters

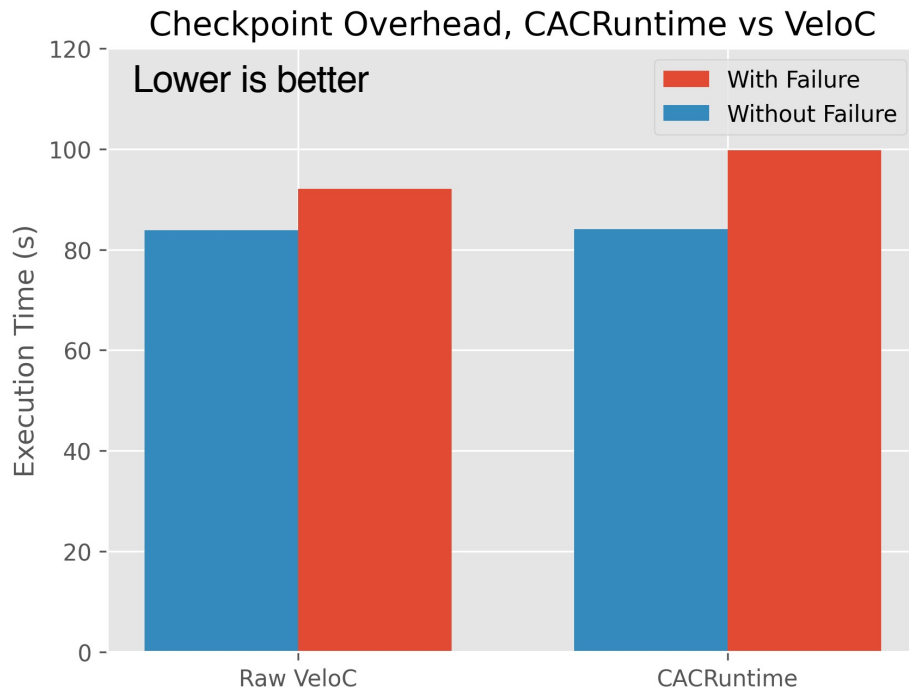
- 8 nodes
 - ✓ 256MB checkpoint per node
- 6 Checkpoints per (non-failing) run

□ Contrived MTTF

- On the order of seconds



Heat-distribution Results



❑ CACRuntime

- Very similar to the baseline without failures
- Some more overhead in recovery
 - ✓ Doesn't know to skip restoring initialization views
 - ✓ Checkpoints an additional time after recovery
- Our version checkpoints as frequently as the raw version after frequency optimization
- Our version does not require any manual checkpointing



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CONCLUSIONS

Conclusions

- ❑ Introduced A Cooperative Compiler and Runtime Checkpoint/Restart Approach for Kokkos
 - Compiler automatically generates checkpoint/restart API
 - ✓ robust enough to translate 9 Kokkos+MPI applications
 - Runtime automatically
 - ✓ performs write aggregation
 - ✓ optimizes the frequency of checkpointing
 - ✓ updates loop iteration variables after recovery



Future Work

- ❑ Evaluate other applications
 - An MPI version of Polybench
 - Mantevo
 - NimbleSM
- ❑ Compare our approach with Kokkos-Resilience [1]
- ❑ Enhance the CACRuntime
 - Asynchronous Checkpointing
 - Runtime Alias Checking
- ❑ Use GPUs

