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The Basics of Orthogonal Polynomials

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The Basics of Orthogonal Polynomials

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Abstract

To understand the mathematics behind Uncertainty Quantification (UQ), one first needs to understand the basics of orthogonal polynomials, which this report covers.

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Definition 1 (Coefficient Sequence). A coefficient sequence \mathbf{c} is simply a finite sequence of complex numbers, such that, either (a) \mathbf{c} is empty (i.e., the sequence has no elements), or (b) \mathbf{c} is non-empty and its last element is non-zero (i.e., $\mathbf{c} = [c_0, c_1, c_2, \dots, c_i]$, where $c_i \neq 0$).

Definition 2 (Polynomial). Given a coefficient sequence \mathbf{c} , we define a function $p_{\mathbf{c}}$, that takes a complex number z as input, and returns another complex number $p_{\mathbf{c}}(z)$ as output. This function, known as the “polynomial” corresponding to the coefficient sequence \mathbf{c} , is given by:

$$p_{\mathbf{c}}(z) = \begin{cases} c_0 + c_1z + c_2z^2 + \dots + c_iz^i = \sum_{j=0}^i c_jz^j, & \text{if } \mathbf{c} = [c_0, c_1, c_2, \dots, c_i] \text{ is non-empty, and} \\ 0, & \text{if } \mathbf{c} \text{ is empty.} \end{cases} \quad (1)$$

Or, we could use “lambda function” notation (e.g., as in Python® [1]), to say:

$$p_{\mathbf{c}} = \begin{cases} \lambda z : c_0 + c_1z + c_2z^2 + \dots + c_iz^i = \lambda z : \sum_{j=0}^i c_jz^j, & \text{if } \mathbf{c} = [c_0, c_1, c_2, \dots, c_i] \text{ is non-empty, and} \\ \lambda z : 0, & \text{if } \mathbf{c} \text{ is empty.} \end{cases} \quad (2)$$

Remark 3. If \mathbf{c} is empty, we say that $p_{\mathbf{c}}$ is “identically zero”.

Observation 4 (Canonicity of Coefficient Sequences and Polynomials). If $p_{\mathbf{c}_1}$ and $p_{\mathbf{c}_2}$ are identical polynomial functions (i.e., $p_{\mathbf{c}_1}(z) = p_{\mathbf{c}_2}(z)$, for every complex number z), then their corresponding coefficient sequences, \mathbf{c}_1 and \mathbf{c}_2 , must be identical as well. Conversely, if two coefficient sequences \mathbf{c}_1 and \mathbf{c}_2 are identical, then their corresponding polynomial functions, $p_{\mathbf{c}_1}$ and $p_{\mathbf{c}_2}$, will also be identical (by Definition 2).

Definition 5 (Degree of a Polynomial). Given a polynomial $p_{\mathbf{c}}$, corresponding to the coefficient sequence \mathbf{c} , we define the degree of $p_{\mathbf{c}}$ as follows:

$$\text{degree}(p_{\mathbf{c}}) = \begin{cases} i, & \text{if } \mathbf{c} = [c_0, c_1, c_2, \dots, c_i] \text{ is non-empty, and} \\ -1, & \text{if } \mathbf{c} \text{ is empty.} \end{cases} \quad (3)$$

Definition 6 (Monic Polynomial). A polynomial $p_{\mathbf{c}}$, corresponding to the coefficient sequence \mathbf{c} , is called “monic” if (a) \mathbf{c} is non-empty, and (b) the highest coefficient $c_i = 1$ (where $\mathbf{c} = [c_0, c_1, c_2, \dots, c_i]$).

Remark 7. We denote the set of complex numbers by \mathbb{C} , the set of real numbers by \mathbb{R} , and the set of integers by \mathbb{Z} .

Definition 8 (Real Polynomial). A polynomial $p_{\mathbf{c}}$, corresponding to the coefficient sequence \mathbf{c} , is called “real” if either, (a) \mathbf{c} is empty, or (b) (if \mathbf{c} is non-empty) the coefficients in \mathbf{c} are all real; that is, $c_j \in \mathbb{R}$ for every integer j such that $0 \leq j \leq i$ (where $\mathbf{c} = [c_0, c_1, c_2, \dots, c_i]$).

Definition 9 (Moment Functional). Given an infinite sequence of complex numbers $\mu = [\mu_0, \mu_1, \mu_2, \dots]$, we define a function \mathcal{L}_{μ} , that takes a polynomial $p_{\mathbf{c}}$ as input, and returns a complex number $\mathcal{L}_{\mu}(p_{\mathbf{c}})$

as output. This function, known as the “moment functional” associated with the “moment sequence” μ , is given by:

$$\mathcal{L}_\mu(p_c) = \begin{cases} \mathcal{L}_\mu\left(\lambda z : \sum_{j=0}^i c_j z^j\right) = \sum_{j=0}^i c_j \mu_j, & \text{if } c = [c_0, c_1, c_2, \dots, c_i] \text{ is non-empty, and} \\ 0, & \text{if } c \text{ is empty.} \end{cases} \quad (4)$$

Observation 10. With the above definition, we have $\mathcal{L}_\mu(\lambda z : z^i) = \mu_i$, for every integer $i \geq 0$.

Observation 11. \mathcal{L}_μ is “additive”. That is,

$$\mathcal{L}_\mu(p_{c_1} + p_{c_2}) = \mathcal{L}_\mu(\lambda z : p_{c_1}(z) + p_{c_2}(z)) = \mathcal{L}_\mu(p_{c_1}) + \mathcal{L}_\mu(p_{c_2}), \text{ for any two polynomials } p_{c_1} \text{ and } p_{c_2}. \quad (5)$$

Observation 12. \mathcal{L}_μ is “homogeneous”. That is,

$$\mathcal{L}_\mu(\alpha p_c) = \mathcal{L}_\mu(\lambda z : \alpha p_c(z)) = \alpha \mathcal{L}_\mu(p_c), \text{ for any polynomial } p_c \text{ and any complex number } \alpha. \quad (6)$$

Observation 13. \mathcal{L}_μ is “linear”, because it satisfies the additivity and homogeneity conditions above.

Example 14 (Weighted Integral Moment Functional). This moment functional will be of particular interest to us. Given a non-negative “weight function” w , with a positive measure, that is defined on the real interval $[a, b]$, we define the corresponding moment sequence $\mu = \{\mu_i \mid i \in \mathbb{Z}, i \geq 0\}$ to be:

$$\mu_i = \int_{x=a}^b w(x) x^i dx \quad (7)$$

In particular, we will be interested in the case where the weight function w above happens to be the probability density function of a random variable.

Definition 15 (Orthogonal Polynomial Sequence or OPS). Given a moment functional \mathcal{L}_μ , we define an infinite sequence of polynomials $p_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ (corresponding to the coefficient sequences $[c_0, c_1, c_2, \dots]$ respectively) to be an orthogonal polynomial sequence (or OPS) with respect to \mathcal{L}_μ , if the following three conditions are satisfied:

1. For every integer $i \geq 0$, p_{c_i} is a polynomial of degree i ,
2. For every two integers i and j such that $i \geq 0, j \geq 0$, and $i \neq j$, we have:

$$\mathcal{L}_\mu(p_{c_i} \times p_{c_j}) = \mathcal{L}_\mu(\lambda z : p_{c_i}(z) \times p_{c_j}(z)) = 0, \text{ and} \quad (8)$$

3. For every integer $i \geq 0$, we have:

$$\mathcal{L}_\mu(p_{c_i}^2) = \mathcal{L}_\mu(\lambda z : p_{c_i}(z) \times p_{c_i}(z)) \neq 0. \quad (9)$$

Observation 16. Given an arbitrary moment sequence μ , there may or may not exist an OPS associated with \mathcal{L}_μ .

Observation 17. Let $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ be an OPS, with respect to the moment functional \mathcal{L}_μ . Then, $\mathbf{q}_{\mathcal{L}_\mu} = [\beta_0 p_{c_0}, \beta_1 p_{c_1}, \beta_2 p_{c_2}, \dots]$ is also an OPS with respect to \mathcal{L}_μ , for any sequence of non-zero complex numbers $[\beta_0, \beta_1, \beta_2, \dots]$. Thus, if an OPS with respect to \mathcal{L}_μ exists, then there exist infinitely many OPSes with respect to \mathcal{L}_μ .

Definition 18 (Monic OPS). An OPS $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ is called “monic” if every polynomial $(p_{c_0}, p_{c_1}, p_{c_2}, \dots)$ in the OPS is monic (as in [Definition 6](#)).

Definition 19 (K-OPS). Given (a) a moment functional \mathcal{L}_μ , and (b) an infinite sequence of non-zero complex numbers, $\mathbf{K} = [K_0, K_1, K_2, \dots]$, we define $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ to be a **K-OPS** with respect to \mathcal{L}_μ , if it is an OPS with respect to \mathcal{L}_μ that, in addition to satisfying the regular OPS conditions of [Definition 15](#), also satisfies the following:

$$\mathcal{L}_\mu(z^i p_{c_i}) = \mathcal{L}_\mu(\lambda z : z^i \times p_{c_i}(z)) = K_i, \text{ for every integer } i \geq 0. \quad (10)$$

Definition 20 (Normalised OPS, or Orthonormal Polynomial Sequence). Given a moment functional \mathcal{L}_μ , we define $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ to be a “normalised” OPS with respect to \mathcal{L}_μ , or an **orthonormal** polynomial sequence with respect to \mathcal{L}_μ , if it is an OPS with respect to \mathcal{L}_μ that, in addition to satisfying the regular OPS conditions of [Definition 15](#), also satisfies the following:

$$\mathcal{L}_\mu(p_{c_i}^2) = 1, \text{ for every integer } i \geq 0. \quad (11)$$

Definition 21 (Positively Oriented Complex Number). We say that a complex number $z = \operatorname{Re}(z) + \sqrt{-1} \cdot \operatorname{Im}(z)$ is “positively oriented” if it lies either in the open right half of the complex plane, or on the positive imaginary axis. That is, we have:

$$(\operatorname{Re}(z) > 0) \text{ or } (\operatorname{Re}(z) = 0 \text{ and } \operatorname{Im}(z) > 0) \quad (12)$$

Definition 22 (Positively Oriented OPS). Given a moment functional \mathcal{L}_μ , we define $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ to be a “positively oriented” OPS with respect to \mathcal{L}_μ if it is an OPS with respect to \mathcal{L}_μ that, in addition to satisfying the regular OPS conditions of [Definition 15](#), is also such that, for each $i \geq 0$, the highest coefficient in p_{c_i} (i.e., the coefficient of z^i in the expansion of $p_{c_i}(z)$) is positively oriented.

Definition 23 (Negatively Oriented OPS). Given a moment functional \mathcal{L}_μ , we define $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ to be a “negatively oriented” OPS with respect to \mathcal{L}_μ if it is an OPS with respect to \mathcal{L}_μ that, in addition to satisfying the regular OPS conditions of [Definition 15](#), is also such that, for each $i \geq 0$, the highest coefficient in p_{c_i} (i.e., the coefficient of z^i in the expansion of $p_{c_i}(z)$) lies in the region of the complex plane given by $\{z \mid \operatorname{Re}(z) < 0 \text{ or } (\operatorname{Re}(z) = 0 \text{ and } \operatorname{Im}(z) < 0)\}$ (i.e., the open left half of the complex plane, plus the negative imaginary axis).

Observation 24. An OPS may be neither positively nor negatively oriented.

Theorem 25. Let $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ be an OPS, with respect to the moment functional \mathcal{L}_μ . Then, $p_{c_0} = \lambda z : c$, where c is a non-zero complex number.

Proof. From the first condition of [Definition 15](#), p_{c_0} is a polynomial of degree 0. That is, we have:

$$\text{degree}(p_{c_0}) = 0. \quad (13)$$

From [Definition 5](#), this is possible only when the coefficient sequence c_0 has exactly one element. This element, from [Definition 1](#), must be a complex number. If we call this complex number c , we have:

$$c_0 = [c]. \quad (14)$$

Applying [Definition 2](#) to the above, we have:

$$p_{c_0} = \lambda z : c. \quad (15)$$

Now, all that's left to prove is that $c \neq 0$. This we can see from [Definition 1](#): c , being the last element of a non-empty coefficient sequence c_0 , has to be non-zero. \square

Theorem 26. *Let $p_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ be an OPS, with respect to the moment functional \mathcal{L}_μ . Then, for every $i \geq 1$, $\mathcal{L}_\mu(p_{c_i}) = 0$.*

Proof. From [Theorem 25](#), let $p_{c_0} = \lambda z : c$, where c is a non-zero complex number. Then, for every $i \geq 1$, we have:

$$\begin{aligned} \mathcal{L}_\mu(p_{c_i}) &= \frac{c}{c} \mathcal{L}_\mu(p_{c_i}) \\ &= \frac{1}{c} \mathcal{L}_\mu(c p_{c_i}) && \text{(from } \text{Observation 12)} \\ &= \frac{1}{c} \mathcal{L}_\mu(p_{c_0} \times p_{c_i}) \\ &= \frac{1}{c} 0 && \text{(from the second condition of } \text{Definition 15)} \\ &= 0. \end{aligned} \quad (16)$$

\square

Definition 27 (Vector Spaces of Polynomials). For every $i \geq 0$, we define \mathbb{P}_i to be the vector space of polynomials of degree at most i , over the field of complex numbers \mathbb{C} . Operations such as field addition, field multiplication, vector addition, and multiplication of a vector by a scalar, are assumed to be done “the natural way” and are not defined precisely here.

Observation 28. For every $i \geq 0$, the set of polynomials given by $\{\lambda z : z^j \mid 0 \leq j \leq i\}$ forms a basis for the vector space \mathbb{P}_i . Thus, the dimension of \mathbb{P}_i is $i + 1$ [2].

Theorem 29 (Basis Sets of Orthogonal Polynomials). *Let $p_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ be an OPS, with respect to the moment functional \mathcal{L}_μ . Then, for every $i \geq 0$, the set of polynomials $\{p_{c_0}, p_{c_1}, p_{c_2}, \dots, p_{c_i}\}$ forms a basis for the vector space \mathbb{P}_i .*

Proof. For every $i \geq 1$, because p_{c_i} is of degree i , it cannot be written as a linear combination of the polynomials $\{p_{c_0}, p_{c_1}, p_{c_2}, \dots, p_{c_{i-1}}\}$, which all have smaller degrees. Therefore, the polynomials $\{p_{c_0}, p_{c_1}, p_{c_2}, \dots, p_{c_i}\}$ are linearly independent. But these are $i+1$ linearly independent “vectors” from the vector space \mathbb{P}_i , which, by [Observation 28](#), is of dimension $i+1$. Therefore, by the dimension theorem for vector spaces (see, for example, [2]), the polynomials $\{p_{c_0}, p_{c_1}, p_{c_2}, \dots, p_{c_i}\}$ must form a basis for \mathbb{P}_i . \square

Theorem 30. Let $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ be an OPS, with respect to the moment functional \mathcal{L}_μ . Then, for each $i \geq 0$, $\mathcal{L}_\mu(q \times p_{c_i}) = 0$ for every polynomial q of degree smaller than i , and $\mathcal{L}_\mu(q \times p_{c_i}) \neq 0$ for every polynomial q of degree equal to i .

Proof. Let q be a polynomial of degree j , which is at most i .

Now, if q is identically zero (i.e., $j = -1$), then $q \times p_{c_i}$ will also be identically zero for every $i \geq 0$, and thus $\mathcal{L}_\mu(q \times p_{c_i})$ will trivially be zero (from [Definition 9](#)) for every $i \geq 0$. So let us just consider situations where $j \geq 0$. That is,

$$\text{degree}(q) = j, \text{ where } 0 \leq j \leq i. \quad (17)$$

Then, by the “basis sets of orthogonal polynomials” theorem ([Theorem 29](#)), we can write q as a linear combination of the polynomials in $\{p_{c_0}, p_{c_1}, p_{c_2}, \dots, p_{c_j}\}$. That is, we have:

$$q = \sum_{k=0}^j \alpha_k p_{c_k} \quad (18)$$

Also, in the expansion above, the last coefficient α_j will be non-zero. This is because the left hand side (q) of the equation above is of degree j , and if α_j were zero, the linear combination on the right hand side would have a degree strictly smaller than j , which would be a contradiction. Thus, we have:

$$\alpha_j \neq 0. \quad (19)$$

Multiplying [Eq. \(18\)](#) by p_{c_i} and applying \mathcal{L}_μ to both sides, we get:

$$\begin{aligned} \mathcal{L}_\mu(q \times p_{c_i}) &= \mathcal{L}_\mu\left(\left(\sum_{k=0}^j \alpha_k p_{c_k}\right) \times p_{c_i}\right) \\ &= \mathcal{L}_\mu\left(\sum_{k=0}^j \alpha_k (p_{c_k} \times p_{c_i})\right) \\ &= \sum_{k=0}^j \alpha_k \mathcal{L}_\mu(p_{c_k} \times p_{c_i}) \quad (\text{by the linearity of } \mathcal{L}_\mu, \text{ i.e., } \text{Observation 13}) \end{aligned} \quad (20)$$

Now, from [Eq. \(17\)](#), $j \leq i$. If $j < i$, each term in the summation above vanishes due to the second condition from [Definition 15](#). But if $j = i$, the last term in the summation (corresponding to

$k = j = i$) alone survives. Thus, we have:

$$\mathcal{L}_\mu(q \times p_{c_i}) = \begin{cases} 0, & \text{if } j < i, \text{ and} \\ \alpha_i \mathcal{L}_\mu(p_{c_i}^2), & \text{if } j = i. \end{cases} \quad (21)$$

But if $j = i$, we have $\alpha_i \neq 0$ from Eq. (19), and we also have $\mathcal{L}_\mu(p_{c_i}^2) \neq 0$ from the third condition of Definition 15. Therefore, Eq. (21) allows us to conclude that, for every $i \geq 0$:

$$\mathcal{L}_\mu(q \times p_{c_i}) = \begin{cases} 0, & \text{if } \text{degree}(q) < i, \text{ and} \\ \neq 0, & \text{if } \text{degree}(q) = i, \end{cases} \quad (22)$$

where we have replaced j by $\text{degree}(q)$ using Eq. (17). \square

Theorem 31 (Converse of Theorem 30). *Let $\mathbf{p} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ be an infinite sequence of polynomials (with coefficient sequences $[c_0, c_1, c_2, \dots]$ respectively), where p_{c_i} has degree i , for every $i \geq 0$. Also, let \mathcal{L}_μ be a moment functional, with respect to the moment sequence $[\mu_0, \mu_1, \mu_2, \dots]$. Suppose that, for every $i \geq 0$, we have $\mathcal{L}_\mu(q \times p_{c_i}) = 0$ for every polynomial q of degree smaller than i , and that $\mathcal{L}_\mu(q \times p_{c_i}) \neq 0$ for every polynomial q of degree equal to i . Then, \mathbf{p} is an OPS with respect to \mathcal{L}_μ .*

Proof. Since, for every $i \geq 0$, p_{c_i} is given to be a polynomial of degree i , the sequence \mathbf{p} clearly satisfies the first condition of Definition 15. Now, choose any $i > 0$ and $j > 0$. There are only three possibilities:

1. $i < j$. In this case, letting “ q ” = p_{c_i} and “ i ” = j in “ $\mathcal{L}_\mu(q \times p_{c_i}) = 0$ for every polynomial q of degree smaller than i ”, we get $\mathcal{L}_\mu(p_{c_i} \times p_{c_j}) = 0$.
2. $i > j$. In this case, letting “ q ” = p_{c_j} and “ i ” = i in “ $\mathcal{L}_\mu(q \times p_{c_i}) = 0$ for every polynomial q of degree smaller than i ”, we get $\mathcal{L}_\mu(p_{c_j} \times p_{c_i}) = \mathcal{L}_\mu(p_{c_i} \times p_{c_j}) = 0$.
3. $i = j$. In this case, letting “ q ” = p_{c_i} and “ i ” = j ($= i$) in “ $\mathcal{L}_\mu(q \times p_{c_i}) \neq 0$ for every polynomial q of degree equal to i ”, we get $\mathcal{L}_\mu(p_{c_i} \times p_{c_j}) = \mathcal{L}_\mu(p_{c_i}^2) \neq 0$.

Thus, we have:

$$\mathcal{L}_\mu(p_{c_i} \times p_{c_j}) = \begin{cases} 0, & \text{if } i \neq j, \text{ and} \\ \mathcal{L}_\mu(p_{c_i}^2) \neq 0, & \text{if } i = j, \end{cases} \quad (23)$$

which shows that the sequence \mathbf{p} satisfies the second and third conditions of Definition 15. Therefore, \mathbf{p} is an OPS with respect to \mathcal{L}_μ . \square

Theorem 32. *Let $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ be an OPS, with respect to the moment functional \mathcal{L}_μ . Then, for every $i \geq 0$, given any polynomial q of degree i , we can write q as:*

$$\begin{aligned} q &= \sum_{j=0}^i \alpha_j p_{c_j}, \\ \text{where } \alpha_j &= \frac{\mathcal{L}_\mu(q \times p_{c_j})}{\mathcal{L}_\mu(p_{c_j}^2)}, \text{ for every } 0 \leq j \leq i. \end{aligned} \quad (24)$$

Proof. From [Theorem 29](#), we already know that the set of polynomials $\{p_{c_0}, p_{c_1}, p_{c_2}, \dots, p_{c_i}\}$ forms a basis for the vector space \mathbb{P}_i , for every $i \geq 0$. Also, q , being a polynomial of degree i , is a member of the vector space \mathbb{P}_i . Therefore, q can be written as a linear combination of the polynomials $\{p_{c_0}, p_{c_1}, p_{c_2}, \dots, p_{c_i}\}$. That is, we have:

$$q = \sum_{j=0}^i \alpha_j p_{c_j} \quad (25)$$

Multiplying both sides of the above equation by p_{c_k} (where $0 \leq k \leq i$), and applying \mathcal{L}_μ to both sides, we get:

$$\begin{aligned} \mathcal{L}_\mu(q \times p_{c_k}) &= \mathcal{L}_\mu\left(\left(\sum_{j=0}^i \alpha_j p_{c_j}\right) \times p_{c_k}\right) \\ &= \mathcal{L}_\mu\left(\sum_{j=0}^i \alpha_j (p_{c_j} \times p_{c_k})\right) \\ &= \sum_{j=0}^i \alpha_j \mathcal{L}_\mu(p_{c_j} \times p_{c_k}) \quad (\text{by the linearity of } \mathcal{L}_\mu, \text{ i.e., } \text{Observation 13}) \end{aligned} \quad (26)$$

From the second condition in [Definition 15](#), we see that only the term corresponding to $j = k$ survives in the summation above; the rest of the terms vanish. Therefore, we have:

$$\mathcal{L}_\mu(q \times p_{c_k}) = \alpha_k \mathcal{L}_\mu(p_{c_k} \times p_{c_k}) = \alpha_k \mathcal{L}_\mu(p_{c_k}^2), \text{ for every } k \in \{0, 1, 2, \dots, i\} \quad (27)$$

Dividing both sides by $\mathcal{L}_\mu(p_{c_k}^2)$ (which we know to be non-zero from the third condition of [Definition 15](#)) we get:

$$\alpha_k = \frac{\mathcal{L}_\mu(q \times p_{c_k})}{\mathcal{L}_\mu(p_{c_k}^2)}, \text{ for every } k \in \{0, 1, 2, \dots, i\} \quad (28)$$

Changing the dummy index variable from k to j in [Eq. \(28\)](#), we get:

$$\alpha_j = \frac{\mathcal{L}_\mu(q \times p_{c_j})}{\mathcal{L}_\mu(p_{c_j}^2)}, \text{ for every } j \in \{0, 1, 2, \dots, i\} \quad (29)$$

Combining [Eq. \(25\)](#) and [Eq. \(29\)](#), we get:

$$\begin{aligned} q &= \sum_{j=0}^i \alpha_j p_{c_j} \\ \text{where } \alpha_j &= \frac{\mathcal{L}_\mu(q \times p_{c_j})}{\mathcal{L}_\mu(p_{c_j}^2)}, \text{ for every } 0 \leq j \leq i. \end{aligned} \quad (30)$$

□

Theorem 33 (Uniqueness of OPS upto constant factors). *Given a moment functional \mathcal{L}_μ . Let $\mathbf{p} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ and $\mathbf{q} = [q_{d_0}, q_{d_1}, q_{d_2}, \dots]$ be two orthogonal polynomial sequences associated with \mathcal{L}_μ , with coefficient sequences $[c_0, c_1, c_2, \dots]$ and $[d_0, d_1, d_2, \dots]$ respectively. Then, the polynomials in \mathbf{p} and those in \mathbf{q} will only differ by constant factors; that is, for every $i \geq 0$, $q_{d_i} = \beta_i p_{c_i}$, where $[\beta_0, \beta_1, \beta_2, \dots]$ is a sequence of non-zero complex numbers.*

Proof. Applying [Theorem 32](#), we can write the polynomial q_{d_i} (which is of degree i) as a linear combination of the polynomials $\{p_{c_0}, p_{c_1}, p_{c_2}, \dots, p_{c_i}\}$, as follows:

$$q_{d_i} = \sum_{j=0}^i \alpha_j p_{c_j},$$

where $\alpha_j = \frac{\mathcal{L}_\mu(q_{d_i} \times p_{c_j})}{\mathcal{L}_\mu(p_{c_j}^2)}.$ (31)

But applying [Theorem 30](#), with q_d as the sequence of orthogonal polynomials, q_{d_i} as the i^{th} polynomial (with degree i) in the sequence, and p_{c_j} (for $0 \leq j \leq i$) as the “arbitrary” polynomial of degree at most i , we get:

$$\mathcal{L}_\mu(p_{c_j} \times q_{d_i}) = \mathcal{L}_\mu(q_{d_i} \times p_{c_j}) \begin{cases} = 0, & \text{if } i < j, \text{ and} \\ \neq 0, & \text{if } i = j, \end{cases} \quad (32)$$

for every j from 0 to i .

From [Eq. \(31\)](#) and [Eq. \(32\)](#), we see that only the last term (corresponding to $j = i$) in the summation on the right hand side of [Eq. \(31\)](#) survives; the rest of the terms vanish because $\alpha_j = 0$ for every j from 0 to $i - 1$. Therefore, we have:

$$q_{d_i} = \underbrace{\frac{\mathcal{L}_\mu(q_{d_i} \times p_{c_i})}{\mathcal{L}_\mu(p_{c_i}^2)}}_{\beta_i} p_{c_i}, \text{ for every } i \geq 0. \quad (33)$$

The complex numbers β_i in the equation above are all non-zero. We know this because each β_i is a fraction, where we know the numerator to be non-zero from the “ $i = j$ ” case of [Eq. \(32\)](#), and the denominator to be non-zero from the third condition of [Definition 15](#). □

Remark 34. [Theorem 33](#) tells us that, given a moment functional \mathcal{L}_μ , orthogonal polynomial sequences with respect to \mathcal{L}_μ are essentially unique — except for the non-zero constant factors β_i above.

But often, for concreteness, we will need to single out a particular OPS for discussion, without having to worry about these constant factors. In such cases, we will typically follow one of the three approaches below:

1. We will insist that the OPS be monic (as in [Definition 18](#)), or

2. We will specify an infinite sequence of non-zero complex numbers, $\mathbf{K} = [K_0, K_1, K_2, \dots]$, and we will insist that the OPS be a \mathbf{K} -OPS with respect to \mathcal{L}_μ (as in [Definition 19](#)), or
3. We will insist (a) that the OPS be normalised (as in [Definition 20](#)), and (b) that the OPS be either positively (as in [Definition 22](#)) or negatively (as in [Definition 23](#)) oriented.

Each of the above is enough to guarantee uniqueness. For example, given a moment functional \mathcal{L}_μ , if there exists an OPS for it, there will also exist a monic OPS for it, and this monic OPS will be unique (see [Theorem 35](#)). Similarly, given a moment functional \mathcal{L}_μ , if there exists an OPS for it, there will also exist a unique \mathbf{K} -OPS for it, with respect to any sequence of non-zero complex numbers \mathbf{K} (see [Theorem 36](#)). And finally, given a moment functional \mathcal{L}_μ , if there exists an OPS for it, there will also exist a unique positively/negatively oriented, normalised OPS for it (see [Theorem 37](#)).

Moreover, given a moment functional \mathcal{L}_μ , and any OPS for it, it is easy to get from it the corresponding monic OPS, or the corresponding \mathbf{K} -OPS with respect to any sequence of non-zero complex numbers \mathbf{K} , or the corresponding positively/negatively oriented, normalised OPS.

We state the following theorems without proof.

Theorem 35. *Given a moment functional \mathcal{L}_μ , such that an OPS with respect to \mathcal{L}_μ exists. Then, there exists a unique monic OPS with respect to \mathcal{L}_μ .*

Theorem 36. *Given (a) a moment functional \mathcal{L}_μ , such that an OPS with respect to \mathcal{L}_μ exists, and (b) an infinite sequence of non-zero complex numbers, $\mathbf{K} = [K_0, K_1, K_2, \dots]$. Then, there exists a unique \mathbf{K} -OPS with respect to \mathcal{L}_μ .*

Theorem 37. *Given a moment functional \mathcal{L}_μ , such that an OPS with respect to \mathcal{L}_μ exists. Then, there exists a unique positively oriented, normalised OPS with respect to \mathcal{L}_μ , and a unique negatively oriented, normalised OPS with respect to \mathcal{L}_μ .*

Definition 38 (Moment Matrices). Given a moment sequence $\mu = [\mu_0, \mu_1, \mu_2, \dots]$, we define an infinite sequence of square matrices, $\mathbf{M}_\mu = [\mathbf{M}_{(0, \mu)}, \mathbf{M}_{(1, \mu)}, \mathbf{M}_{(2, \mu)}, \dots]$, known as “moment matrices”, as follows:

$$\mathbf{M}_{(i, \mu)} = \begin{bmatrix} \mu_0 & \mu_1 & \mu_2 & \cdots & \mu_i \\ \mu_1 & \mu_2 & \mu_3 & \cdots & \mu_{i+1} \\ \mu_2 & \mu_3 & \mu_4 & \cdots & \mu_{i+2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \mu_i & \mu_{i+1} & \mu_{i+2} & \cdots & \mu_{2i} \end{bmatrix}, \text{ for every } i \geq 0. \quad (34)$$

Definition 39 (Moment Determinants). Given a moment sequence $\mu = [\mu_0, \mu_1, \mu_2, \dots]$, we define an infinite sequence of complex numbers, $\mathbf{d}_\mu = [d_{(0, \mu)}, d_{(1, \mu)}, d_{(2, \mu)}, \dots]$, known as “moment determinants”, as follows:

$$d_{(i, \mu)} = \det(\mathbf{M}_{(i, \mu)}) = \begin{vmatrix} \mu_0 & \mu_1 & \mu_2 & \cdots & \mu_i \\ \mu_1 & \mu_2 & \mu_3 & \cdots & \mu_{i+1} \\ \mu_2 & \mu_3 & \mu_4 & \cdots & \mu_{i+2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \mu_i & \mu_{i+1} & \mu_{i+2} & \cdots & \mu_{2i} \end{vmatrix}, \text{ for every } i \geq 0. \quad (35)$$

Theorem 40 (Existence of OPS). *Let \mathcal{L}_μ be a moment functional, corresponding to the moment sequence $\mu = [\mu_0, \mu_1, \mu_2, \dots]$. Then, a necessary and sufficient condition for the existence of an OPS with respect to \mathcal{L}_μ is that the moment matrices $[\mathbf{M}_{(0, \mu)}, \mathbf{M}_{(1, \mu)}, \mathbf{M}_{(2, \mu)}, \dots]$ (as defined in Definition 38) are all non-singular, i.e., the moment determinants $[d_{(0, \mu)}, d_{(1, \mu)}, d_{(2, \mu)}, \dots]$ (as defined in Definition 39) are all non-zero.*

Theorem 41. *Given (a) a moment functional \mathcal{L}_μ , and (b) a \mathbf{K} -OPS $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ with respect to \mathcal{L}_μ , corresponding to the sequence $\mathbf{K} = [K_0, K_1, K_2, \dots]$ (as in Definition 19). Then, for every $i \geq 0$, the highest coefficient $c_{(i, i)}$ of p_{c_i} , i.e., the last element of \mathbf{c}_i , or the coefficient multiplying z^i in the expansion of $p_{c_i}(z)$, is given by:*

$$c_{(i, i)} = \frac{K_i d_{(i-1, \mu)}}{d_{(i, \mu)}}, \quad (36)$$

where $d_{(-1, \mu)}$ is defined to be 1.

Definition 42 (S-positive definiteness). Given an infinite set of real numbers S , a moment functional \mathcal{L}_μ is called “positive definite with respect to S ”, or “S-positive definite”, if, for every polynomial p that is (a) non-negative on S (i.e., $p(x) \geq 0$ for every $x \in S$), and (b) not identically zero, we have $\mathcal{L}_\mu(p) > 0$.

Definition 43 (Positive definiteness). A moment functional \mathcal{L}_μ is called “positive definite” if it is S-positive definite, when $S = (-\infty, \infty)$.

Theorem 44 (S-positive definiteness implies positive definiteness). *Let \mathcal{L}_μ be an S-positive definite moment functional, with respect to the infinite set S . Then, \mathcal{L}_μ is also positive definite.*

Theorem 45. *Let \mathcal{L}_μ be a positive definite moment functional, corresponding to the moment sequence $\mu = [\mu_0, \mu_1, \mu_2, \dots]$. Then, for every $i \geq 0$, μ_i is real and $\mu_{2i} > 0$.*

Theorem 46. *Let \mathcal{L}_μ be a positive definite moment functional. Then, an OPS with respect to \mathcal{L}_μ exists.*

Theorem 47. *Let p be a polynomial that is non-negative on $(-\infty, \infty)$, i.e., $p(x) \geq 0$ for all real x . Then, p can always be written as the sum of squares of two real polynomials. That is, we have:*

$$p = q^2 + r^2, \quad (37)$$

where q and r are real polynomials (as in Definition 8).

Theorem 48. *Let \mathcal{L}_μ be a moment functional, corresponding to the moment sequence $\mu = [\mu_0, \mu_1, \mu_2, \dots]$. Then, \mathcal{L}_μ is positive definite if and only if (a) μ_i is real for every $i \geq 0$, and (b) $d_{(i, \mu)} > 0$ for every $i \geq 0$.*

Theorem 49. *Let $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ be an OPS, with respect to the moment functional \mathcal{L}_μ . Then, if (a) p_{c_i} is real for every $i \geq 0$, and (b) $\mathcal{L}_\mu(p_{c_i}^2) > 0$ for every $i \geq 0$, then \mathcal{L}_μ is positive definite.*

Theorem 50 (Three-term Recurrence Formula). *Let $\mathbf{p}_{\mathcal{L}_\mu} = [p_{c_0}, p_{c_1}, p_{c_2}, \dots]$ be an OPS, with respect to the moment functional \mathcal{L}_μ . Then, there exist three sequences of complex numbers, $\gamma = [\gamma_1, \gamma_2, \gamma_3, \dots]$, $\delta = [\delta_1, \delta_2, \delta_3, \dots]$, and $\theta = [\theta_2, \theta_3, \theta_4, \dots]$, such that, for every $z \in \mathbb{C}$, we have:*

$$p_{c_1}(z) = \left(\frac{z - \delta_1}{\gamma_1} \right) p_{c_0}(z), \text{ and}$$

$$p_{\mathbf{c}_i}(z) = \left(\frac{z - \delta_i}{\gamma_i} \right) p_{\mathbf{c}_{i-1}}(z) - \left(\frac{\theta_i}{\gamma_i} \right) p_{\mathbf{c}_{i-2}}(z), \text{ for every } i \geq 2. \quad (38)$$

In particular, we have:

$$\begin{aligned} \gamma_i &= \frac{\mathcal{L}_\mu(z p_{\mathbf{c}_{i-1}} p_{\mathbf{c}_i})}{\mathcal{L}_\mu(p_{\mathbf{c}_i}^2)}, \text{ for every } i \geq 1, \\ \delta_i &= \frac{\mathcal{L}_\mu(z p_{\mathbf{c}_{i-1}}^2)}{\mathcal{L}_\mu(p_{\mathbf{c}_{i-1}}^2)}, \text{ for every } i \geq 1, \text{ and} \\ \theta_i &= \frac{\mathcal{L}_\mu(z p_{\mathbf{c}_{i-2}} p_{\mathbf{c}_{i-1}})}{\mathcal{L}_\mu(p_{\mathbf{c}_{i-2}}^2)}, \text{ for every } i \geq 2. \end{aligned} \quad (39)$$

References

- [1] <https://docs.python.org/3/tutorial/controlflow.html#lambda-expressions>.
- [2] https://en.wikipedia.org/wiki/Dimension_theorem_for_vector_spaces.

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