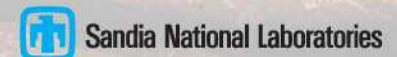


Virtual Machine Monitors as Secure Service Platforms

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Welcome

- Categories of hypervisors
- Type I versus Type II hypervisors for VM IaaS
- Challenges to VM IaaS
- Potential VM services
- Virtualization as a secure platform
- Requirements for building a secure virtual platform
- Interesting academic architectures





VM Infrastructure as a Service

“Adding services via a virtual machine is analogous to adding network services via a firewall. Both virtual machines and firewalls intercept actions at a universal, low-level interface, and both must overcome performance and semantic-gap problems. Just as network firewalls have proven useful for adding network services, we believe virtual machines will prove useful for adding services for the entire computer.”

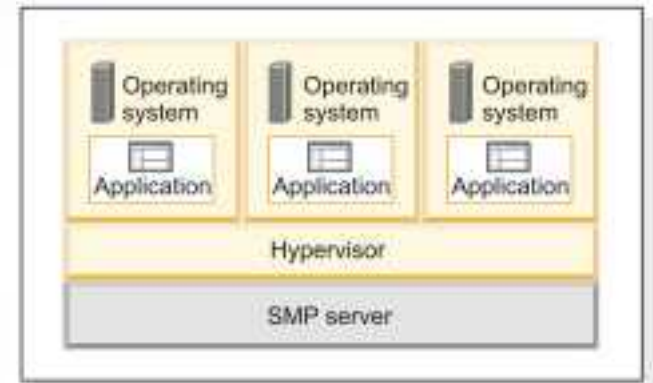
- Chen, P.M., Noble, B.D., “When Virtual Is Better Than Real,” *In Proceedings of the 2001 Workshop on Hot Topics in Operating Systems (HotOS)*, pages 133– 138, May 2001.



Two Categories of Hypervisors

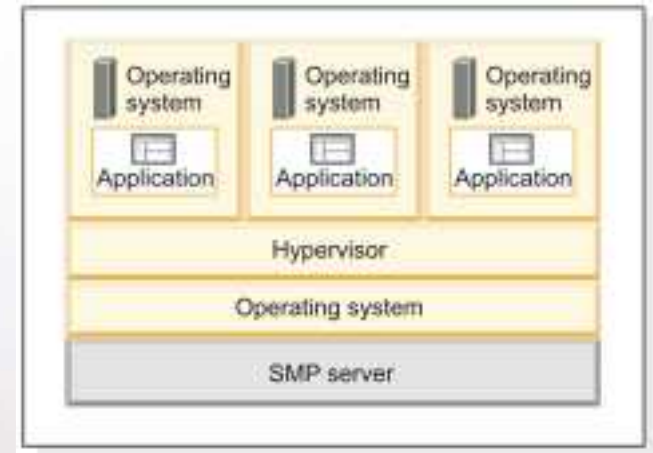
■ Type I Hypervisors

- Bare Metal
- Examples:
 - ◆ VMware ESXi
 - ◆ Hyper-V
 - ◆ “Thin” hypervisors



■ Type II Hypervisors

- Hosted
- Examples:
 - ◆ VMware Workstation
 - ◆ VirtualPC
 - ◆ Parallels

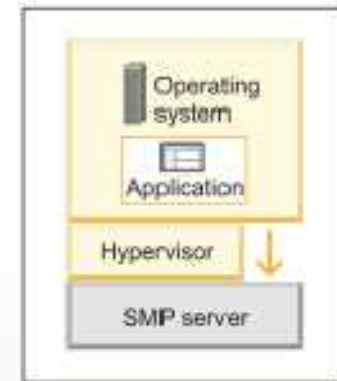


- IBM, “IBM Systems Virtualization,” 2005.



“Thin” Hypervisors

- Subclass of Type I hypervisor
- Single guest
- Guest allowed to interact with some hardware directly
- Hypervisor virtualizes as little as possible to maintain itself, its security model, its features, and its running guest
- Examples: Blue Pill, Vitriol, BitVisor, MAVMM





When One Is Better Than Two

- Most security tools and research today are built with hosted hypervisors
- These require that your VM is an image and runs in a hosted OS. This is good for large servers and honey nets, etc...
- However, for end-point workstations this is less desirable
- Building on top of bare metal hypervisors allows an pre-existing OS to run inside a VM-based security tool
- This likely means only one OS running at a time
- Most principles still apply when using a type two hypervisor





Challenges to VM IaaS

■ Challenges

- ✓ Performance [1]
- Semantic gap [1]
- Introspection subversion [2]
- Hypervisor trust

1. Chen, P.M., Noble, B.D., “When Virtual Is Better Than Real,” *In Proceedings of the 2001 Workshop on Hot Topics in Operating Systems (HotOS)*, pages 133– 138, May 2001.
2. Bahram, S., Jiang, X., Wang, Z., Grace, M., Li, J., Xu, D., “DKSM: Subverting Virtual Machine Introspection for Fun and Profit,” NC State University, TR-2010-4, February 2010.





Potential VM Services

- Secure logging/reporting
- Intrusion prevention and detection
- Monitoring guest components
 - SSDT/IDT hook detection
 - Kernel integrity enforcement
- Debug registers –
 - Strict monitoring, control, and usage of
- ...





Virtualization as a Platform

■ Static “thin” security hypervisors

- BitVisor
- rkAnalyzer (based on BitVisor)
- MAVMM

■ Extensible “thin” security hypervisors

- ?

■ ...DANGER...

- Run-time loadable “ring -1” code
- Easier exploitation by malware
- Larger attack surface
- Sounds a lot like attacking a kernel





VMM Assurance

- ***“... a VMM is a simple-enough mechanism that we can reasonably hope to implement it correctly.” [1]***
- ***“We also assume that SecVisor does not have any vulnerabilities.” [2]***
- **Et cetera...**

1. Garfinkel, T., Rosenblum, M., “A Virtual Machine Introspection Based Architecture for Intrusion Detection,” In *Proceedings of the Networked and Distributed Systems Security Symposium*, February 2003.
2. Seshadri, A., Luk, M., Qu, N., Perrig, A., “SecVisor: A Tiny Hypervisor to Guarantee Lifetime Kernel Code Integrity for Commodity Oses,” In *Proceedings of the ACM Symposium on Operating Systems Principles (SOSP 2007)*, October 2007.





Virtualization as a Secure Platform

- **Can't assume the hypervisor is secure**
 - Do everything we can to secure it
- **Detect something wrong?**
 - Bail out (ala PatchGuard)
- **Attestation for known, good initial state**
- **How do we trust loadable modules**
 - Only load trusted, signed modules (e.g. Driver Signing)
- **Sounds a lot like securing a kernel**





Designing a Secure Virtual Platform

■ Employ hardware protection

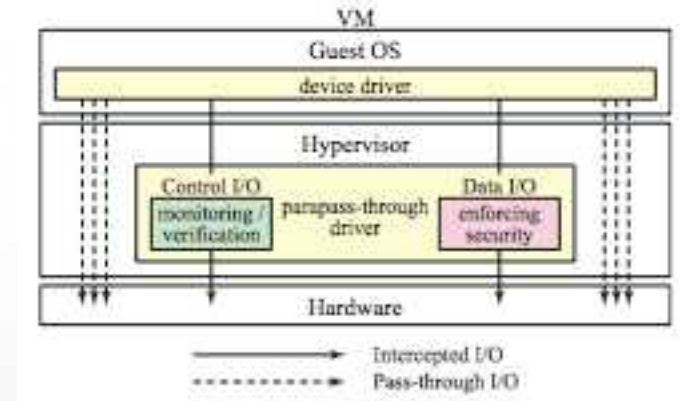
- Paging –
 - ◆ Extended Page Tables (EPT) on Intel
 - ◆ Rapid Virtualization Index (RVI) (aka Nested Page Tables) on AMD
 - ◆ Shadow paging
- DMA –
 - ◆ Virtualization Technology for Direct I/O (Vt-d) on Intel
 - ◆ AMD I/O Virtualization (AMD-Vi aka IOMMU) on AMD
- TPM –
 - ◆ “Thin hypervisor” starts before OS
 - ◆ Load-time hypervisor integrity
 - ◆ Measure similar to Secure Boot



Designing a Secure Virtual Platform

■ Software protection

- Best practices, code auditing, etc...
- Control particular driver interactions
- Hypervisor run-time integrity
- HyperSafe approach



[1]

1. Takahiro Shinagawa et al., "BitVisor: A Thin Hypervisor for Enforcing I/O Device Security," *Proceedings of the 2009 ACM SIGPLAN/SIGOPS Virtual Execution Environments (VEE '09)*, March 2009.





HyperSafe

- **Non-Bypassable Memory Lockdown**
 - Paging-based protections
- **Ensure lifetime $W \oplus X$ for all hypervisor code pages**
 - i.e. After initial setup no page that has ever been writable will be allowed to become executable
- **Ensure no pages are double-mapped to violate this property**
- **Write-protect page tables**
- **Hypervisor uses page-faults in hypervisor to enforce**





HyperSafe continued...

- **Restricted Pointer Indexing**
 - Used for control-flow integrity (CFI)
 - Modification to the compilation
- **Pre-compute indirect control-flow and use function entrance and exit wrappers to enforce this control flow**
- **Mark data pages which contain the control flow tables as read-only**
- Wang, Z., Jiang, X., “HyperSafe: A Lightweight Approach to Provide Lifetime Hypervisor Control-Flow Integrity,” In *Proceedings of the Thirty First IEEE Symposium on Security & Privacy*, Oakland 2010.





Building an Extensible Hypervisor

■ Requirements:

- Kernel best practices
 - ◆ NX, measured launch, integrity checking, etc...
- Utilize hardware
 - ◆ EPT/RVI
 - ◆ Vt-d/AMD-Vi
 - ◆ TPM
- Driver control or device virtualization
 - ◆ Required for interposition of devices
 - ◆ Must be able to know the state of hardware
 - ◆ Doesn't necessarily need full driver





Building an Extensible Hypervisor

■ Requirements:

- Module loader
 - ◆ Signing enforced by loader
 - ◆ Modules not allowed to allocate executable memory
- Symbol resolution
 - ◆ Imported internal symbols resolved for module
 - ◆ Exported symbols from module maintained
- Inspection/interposition API
 - ◆ Inspect guest memory
 - ◆ Assist bridging semantic gap
 - ◆ Interpose on data transfer (i.e. to/from hard disk, network, etc...)





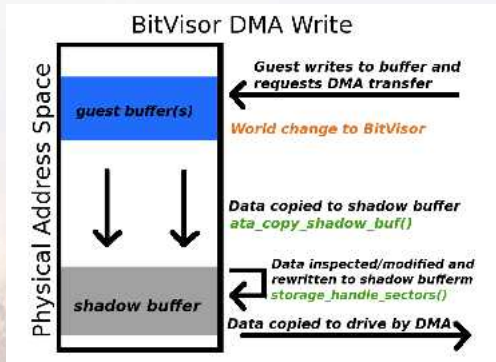
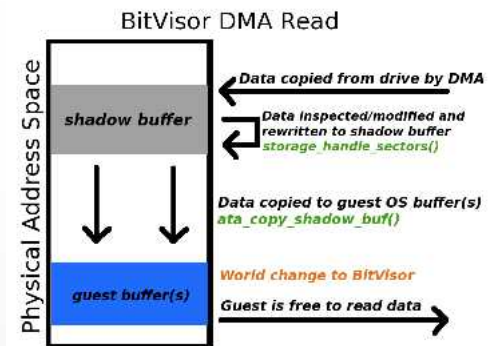
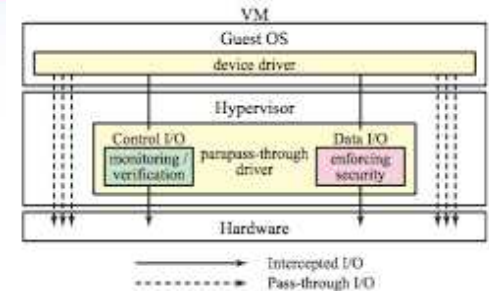
FLOW: A VM Service Case Study

- **FLOW is a disk transfer inspector built on top of the BitVisor hypervisor**
- **Interposes on all disk reads and writes**
- **May sanitize DMA contents before data is written to disk or used by the OS**
- **Does NOT currently reconstruct files or higher-level semantics (i.e. it operates only on DMA transaction buffers)**



FLOW: Implementation

- BitVisor uses the built-in ability to capture all disk reads and writes to implement encryption
- Single choke point, `storage_handle_sectors()`, for hard disk data interposition
- Modified to implement a simple signature scanner/sanitizer
- ❖ This is where an interface could be added to allow additional computation without modifying original code



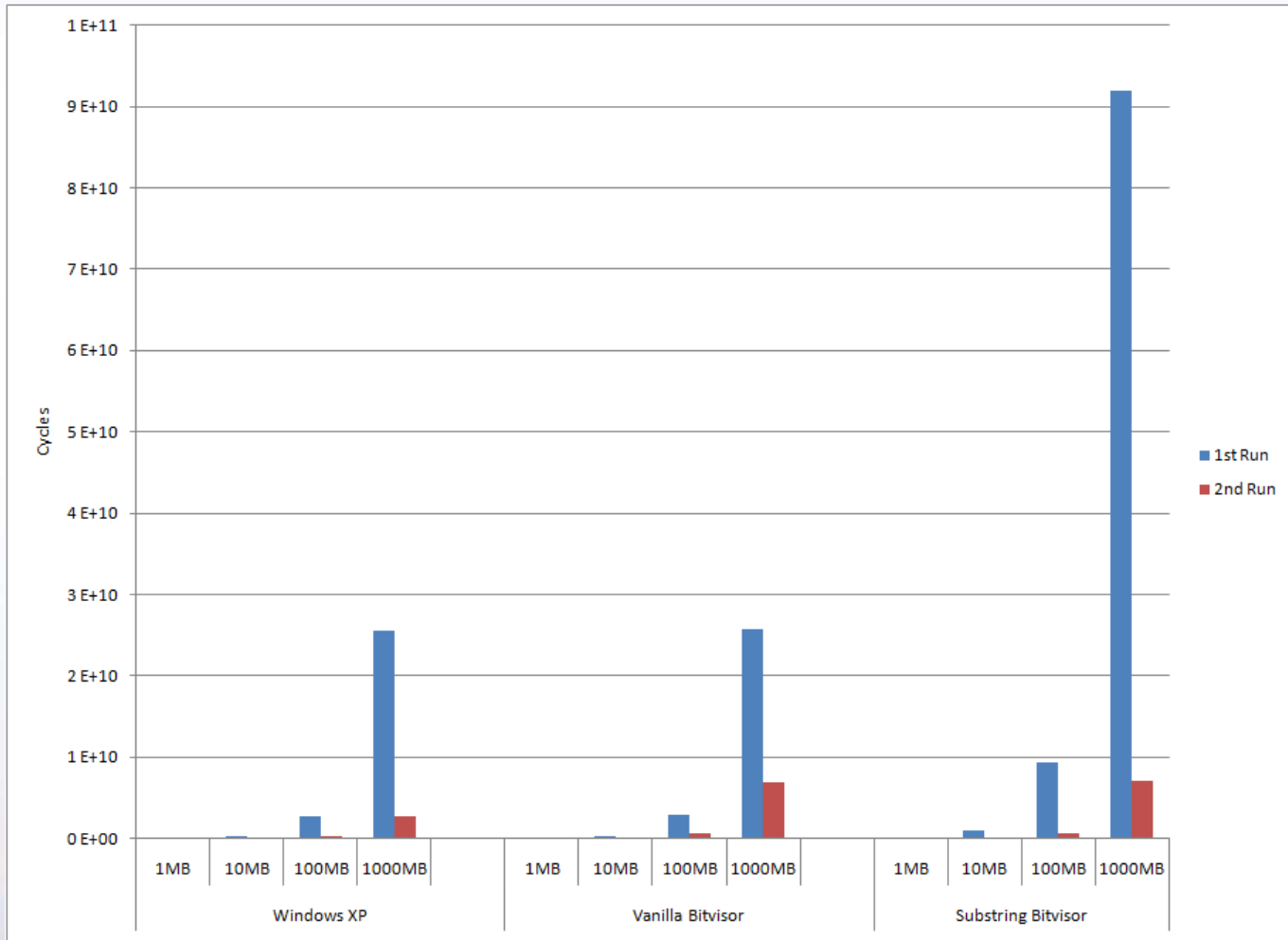


FLOW: Timing Analysis

- Created files of random data
- Timed transfers of size: 1 MB, 10MB, 100MB, and 1GB
- Read file twice
 1. Initial read (no cache)
 2. Second read (with cache from first read)
- Tested three environments:
 - Plain Windows XP
 - Windows XP inside vanilla BitVisor 1.1
 - Windows XP inside BitVisor with substring signatures



FLOW: Performance





FLOW: Outcome

- Inspecting data does add overhead, in addition to the hypervisor
- However, this overhead is insignificant relative to the cost of performing the disk I/O
- In the past, similar results have been shown with network traffic
- ❖ I/O inspection adds minor overhead, but provides very powerful positioning to respond to malicious/undesirable data
- ❖ The hypervisor design should allow the analysis performed on data to be extensible





Parting thoughts

- **Can't implement service on top of hypervisors that are assumed to be secure**
- **Must work toward protected environment**
- **There is still a lot of research to be done in hypervisor security**
- **Hypervisor enables a more closed eco-system that may benefit from restrictions that are less practical in the kernel**
- **Developing extensible hypervisors provides a quick means to add services**



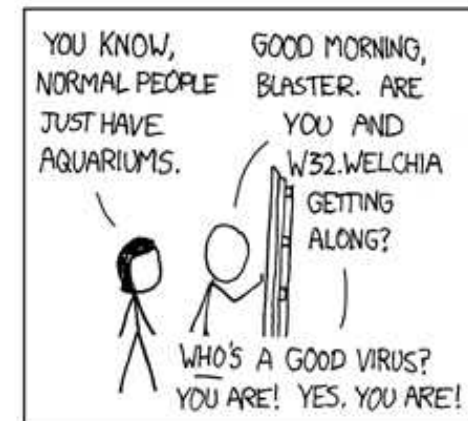
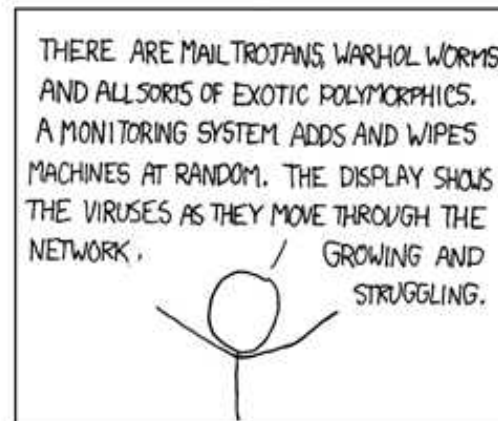
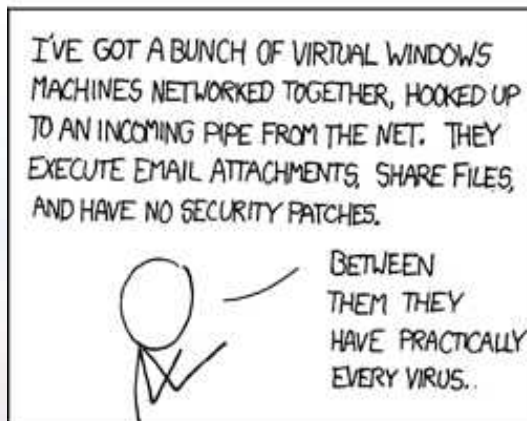
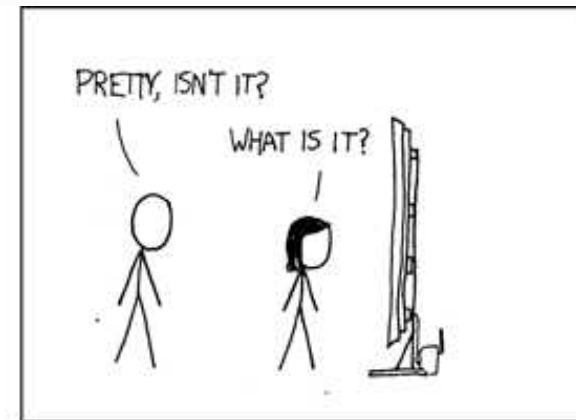
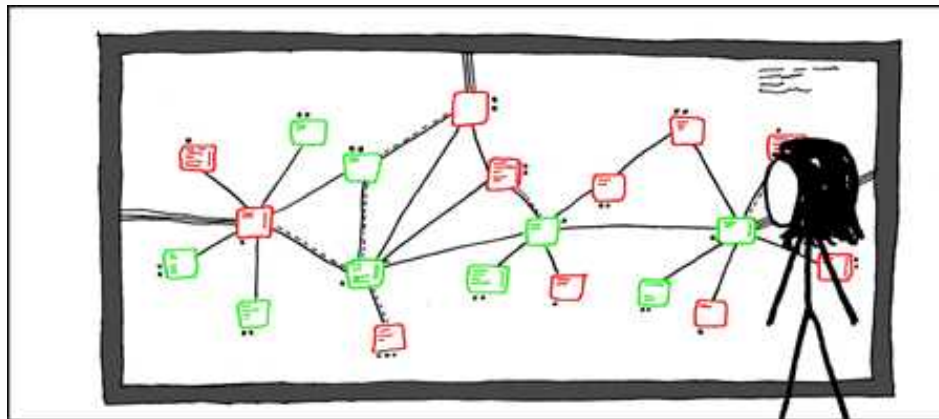


Interesting VM Architectures

- Garfinkel, T., Pfaff, B., Chow, J., Rosenblum, M., Boneh, D., “Terra: A Virtual Machine-Based Platform for Trusted Computing,” in *Proceedings of the 9th ACM Symposium on Operating Systems Principles*, pp. 193–206, 2003.
- Sailer, R., Valdez, E., Jaegar, T., Perez, R., Doorn, L., Griffin, J., Berger, S., “sHype: Secure Hypervisor Approach to Trusted Virtualized Systems,” *IBM Research Report, RC23511, February 2005*.
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- Rutkowska, J., Wojtczuk, R., “Qubes OS Architecture,” Invisible Things Lab, January 2010.
- Douglas, H., “Thin Hypervisor-Based Security Architectures for Embedded Platforms,” Royal Institute of Technology, Sweden, February, 2010.



Questions



- Munroe, R., "Network," <http://xkcd.com/350/>, November 2007.

