



A Scalable Virtualization Environment for HPC

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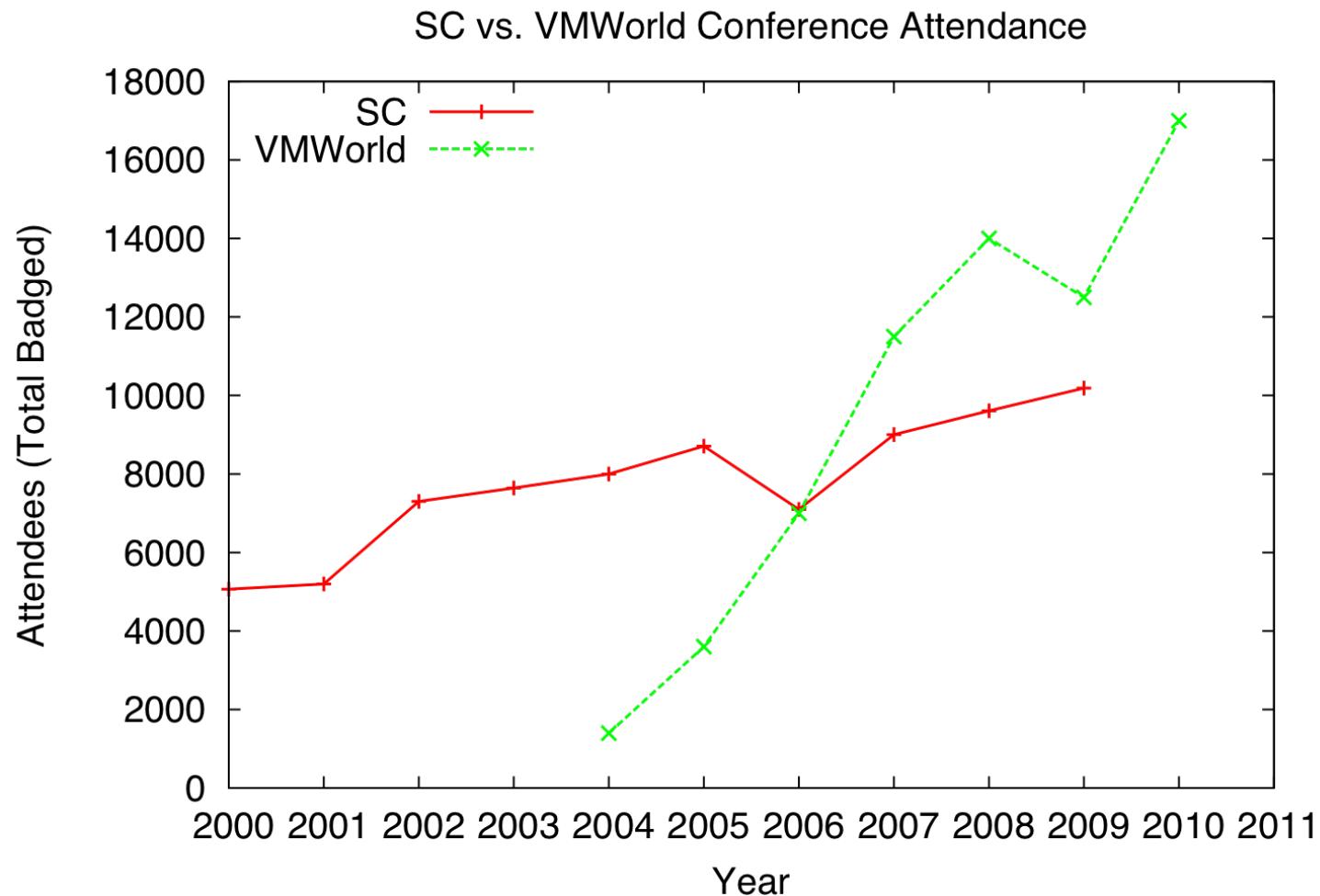


Background – OS Virtualization

- **Treat OS as an application**
- **Major trend in enterprise data center / IT industry over last several years**
- **Motivations**
 - **Server consolidation**
 - **Dynamic workload balancing**
 - **Enhanced security isolation**
 - **On-demand compute capacity, Amazon EC2 “elastic cloud”**
- **Powerful tool for developers, desktop power users**
 - **Run Windows on Linux, run Cplant on laptop, etc.**



Virtualization Seeing Explosive Growth in General Computing Market



Sources: SC web sites, news articles, and blogs



HW-accelerated Virtualization Will Be Baked In

- Any commercially viable platform will have a virtualization story; increasingly sophisticated support
 - x86, AMD, Intel, ...
 - ARM
 - PowerPC
 - Self-virtualizing devices (NICs, GPUs, ...)
- Public clouds beginning to target low/mid HPC
 - Amazon's EC2 Cluster Compute Instances

Can high-end HPC also leverage virtualization?
Does it enable new capabilities?



Key Questions

- **What are the use cases for high-end HPC?**
- **What are the virtualization overheads?**
 - Compute
 - Virtual Memory
 - I/O
- **What can be done to mitigate the overheads?**

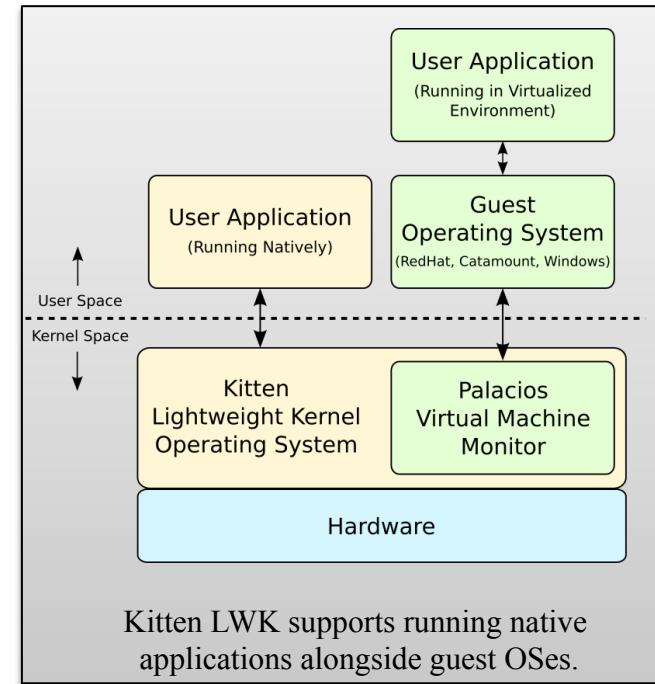


Virtualization Use Cases



Use Case 1: Augment lightweight kernel with VMM to increase flexibility

- Original motivation
- LWK provides high perf. native environment
- VMM allows full-featured guest OS (e.g., Red Hat Linux) to be loaded on-demand
 - Perl, python, matlab, ...
 - COTS databases, simulators, ...
 - You name it
- Approach applies to lightweight Linux distributions like CLE as well





Use Case 2: Tool allowing researchers to test at scale on production machines

- Currently have to request dedicated system time to test prototype system software at scale
 - Long process, difficult to navigate
 - Limited ability to iterate
- Incorporating virtualization into production software stack would allow on-demand loading of custom system software stack(s)
 - Expose effects that only occur at scale
 - VMM can provide enhanced debugging capability compared to native
 - VMM can simulate prototype hardware
 - Issue: performance may be different than native



Use Case 3: Enable New Capabilities

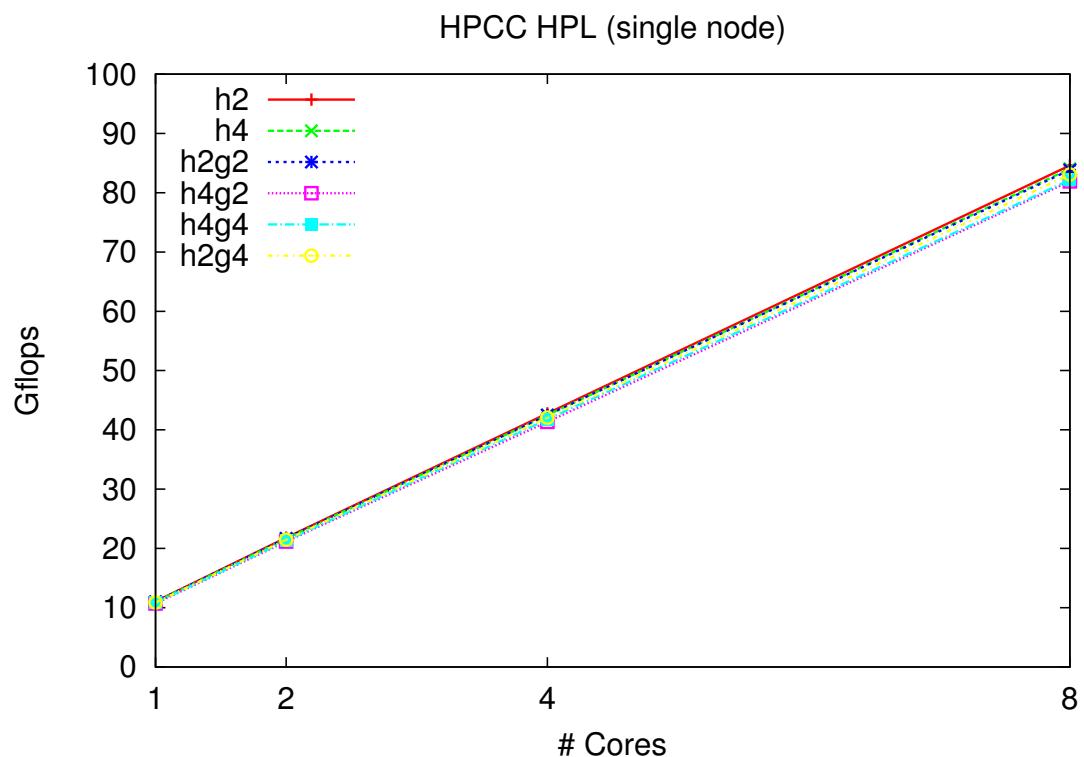
- Perform cybersecurity experiments on capability resources
 - Run commodity OSes + software
 - Multiple virtual nodes per physical node
 - Simulate Internet-scale behavior
- Dynamically replace runtime with one more suitable for the user's workload (e.g., a massive number of small jobs)
- System administrators test new vendor software without taking machine out of production
- Provide backwards capability on future platforms



Virtualization Overheads



Compute Virtualization Essentially Zero



Naming:

h2 = native 2MB paging

n4 = native 4 KB paging

h2g2 = guest memory mapped with 2MB pages, hpcc running in guest using 2 MB pages

h4g2 = guest memory mapped with 4KB pages, hpcc running in guest using 2 MB pages

And so on

Node Configuration:

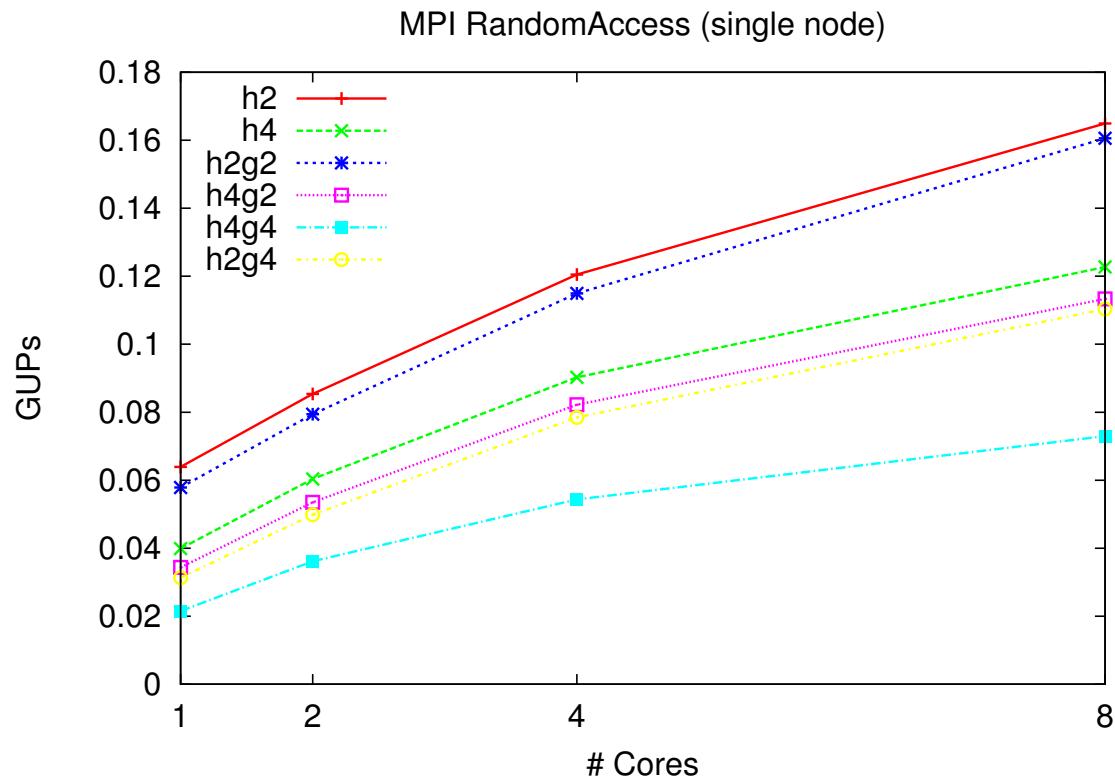
Intel X5570 2.93 GHz (2 sockets, 8 cores)
24 GB RAM (3x 4GB DDR-1333 per socket)
Hyperthreading disabled
Turbo boost disabled

Test Configuration:

Linux 2.6.35, KVM Hypervisor
VCPU to host CPU pinning
Expose NUMA topology to guest
VM uses EPT (aka nested paging)



Memory Virtualization Has Overhead, Using Large Pages Provides Mitigation



Naming:

h2 = native 2MB paging
n4 = native 4 KB paging
h2g2 = guest memory mapped with 2MB pages, hpcc running in guest using 2 MB pages

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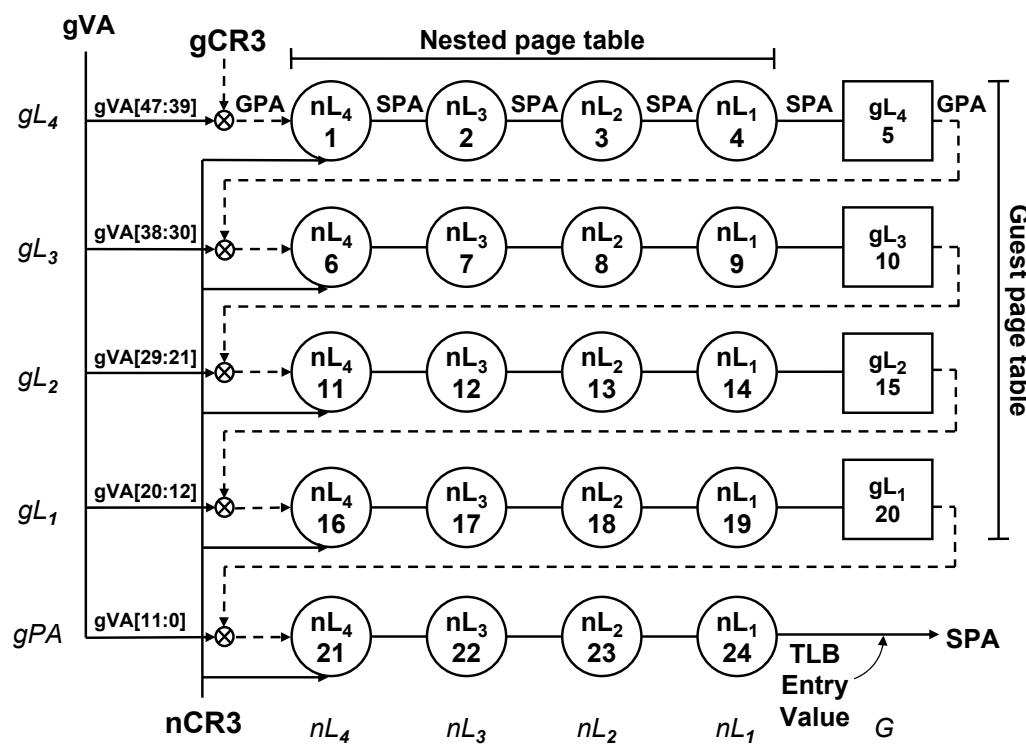
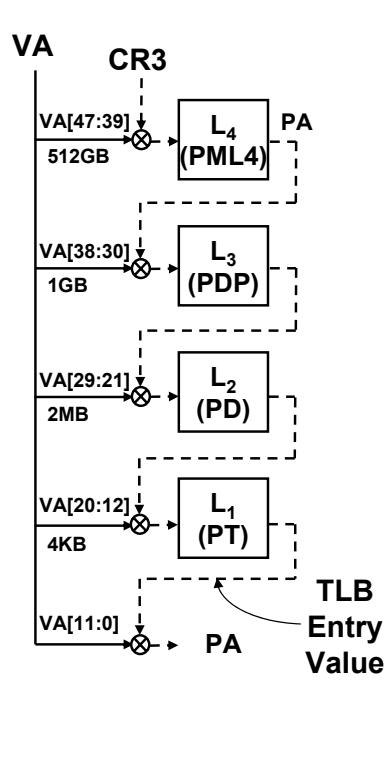
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Nested Paging Memory Virtualization



Normal – 4 levels

Nested – up to 24 memory accesses

Figure from: Ravi Bhargava, Ben Serebrin, Francesco Spanini, and Srilatha Manne.
 Accelerating two-dimensional page walks for virtualized systems.
 In Proceedings ASPLOS'08, March 2008.



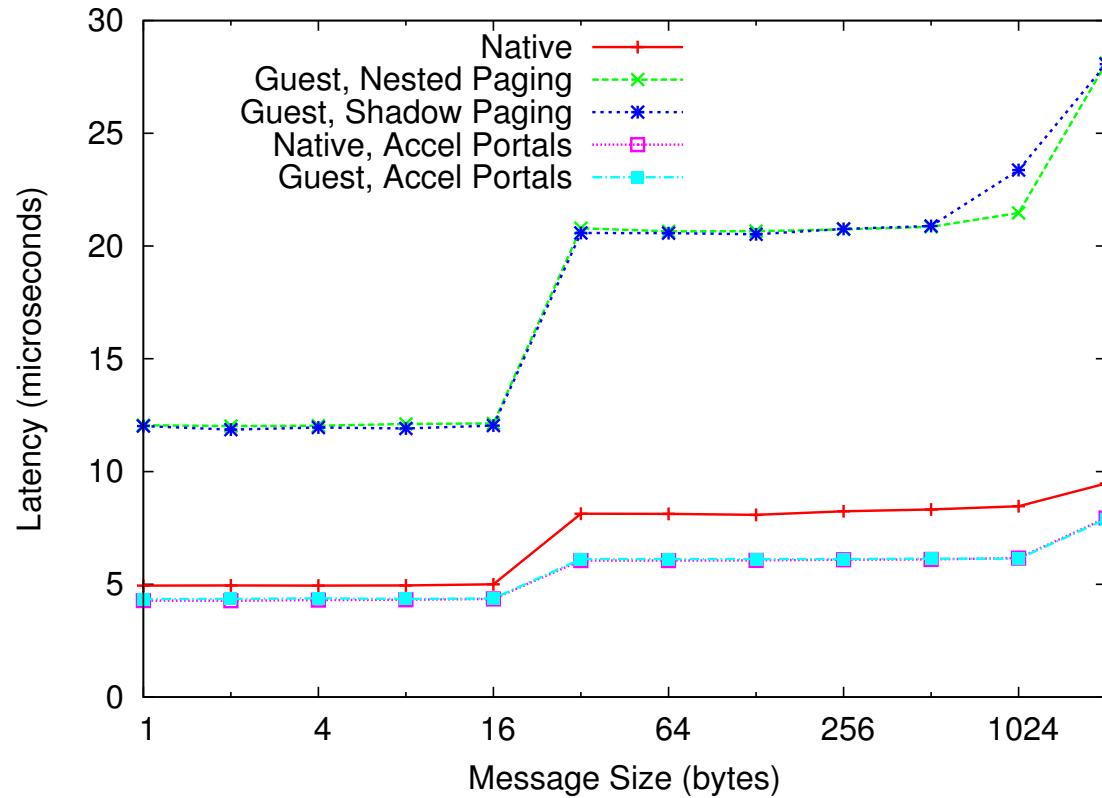
Red Storm Virtualization Experiments

- Testing performed on up to 6240 quad-core Red Storm nodes, also on 48-node test system
- Compared native to guest performance
 - Native = Catamount running on bare metal
 - Guest = Kitten+Palacios running on bare metal, Catamount running as guest OS
- Seastar mapped directly through to guest, interrupts managed by Kitten+Palacios, forwarded to guest
 - Also tested “accelerated portals”, no interrupts
- Compared two guest OS memory management strategies: shadow paging and nested paging



Red Storm PingPong Latency

(Inter-node, SeaStar Passed Through to Guest)

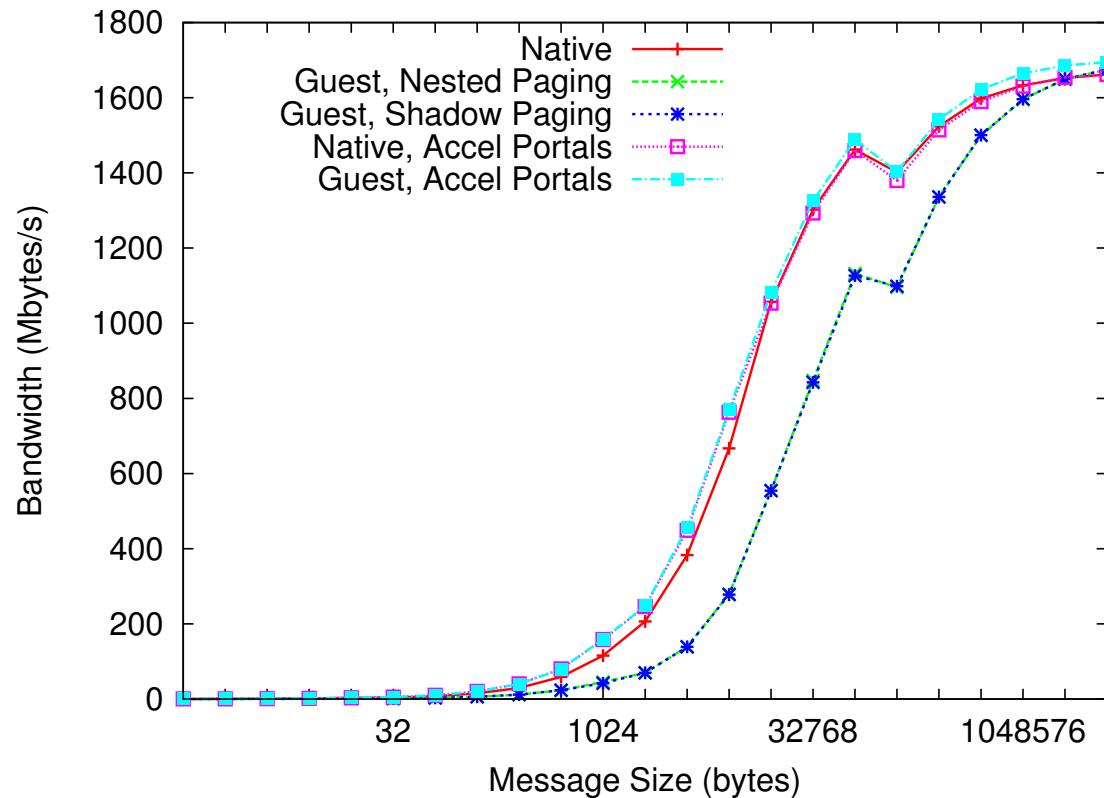


- Interrupt virtualization adds 7 to 14 us overhead for small messages
- Accelerated portals is polling base, so no interrupts.
 - Performance matches native



Red Storm PingPong Bandwidth

(Inter-node, SeaStar Passed Through to Guest)

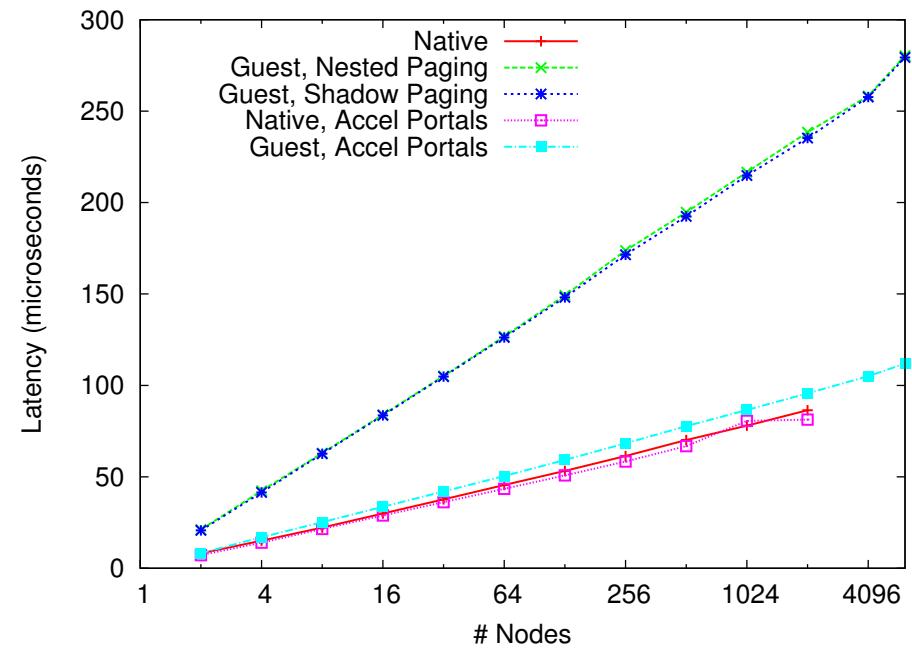
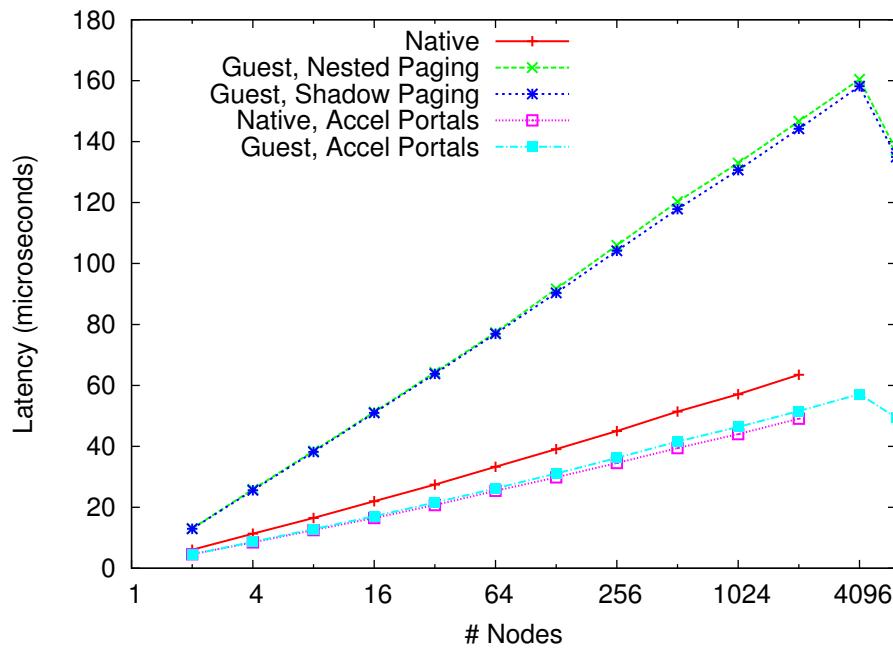


All cases reach same asymptotic bandwidth



Red Storm Reduce and AllReduce Latency

(SeaStar Passed Through to Guest)

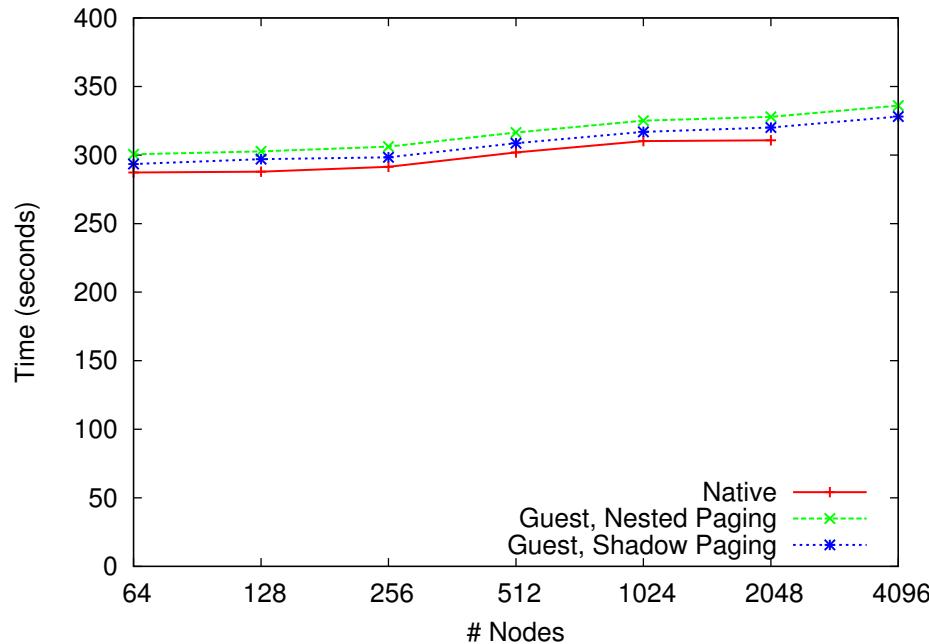


Accelerated Portals Matches Native;
Generic Portals suffers from Interrupt Virtualization
Overhead

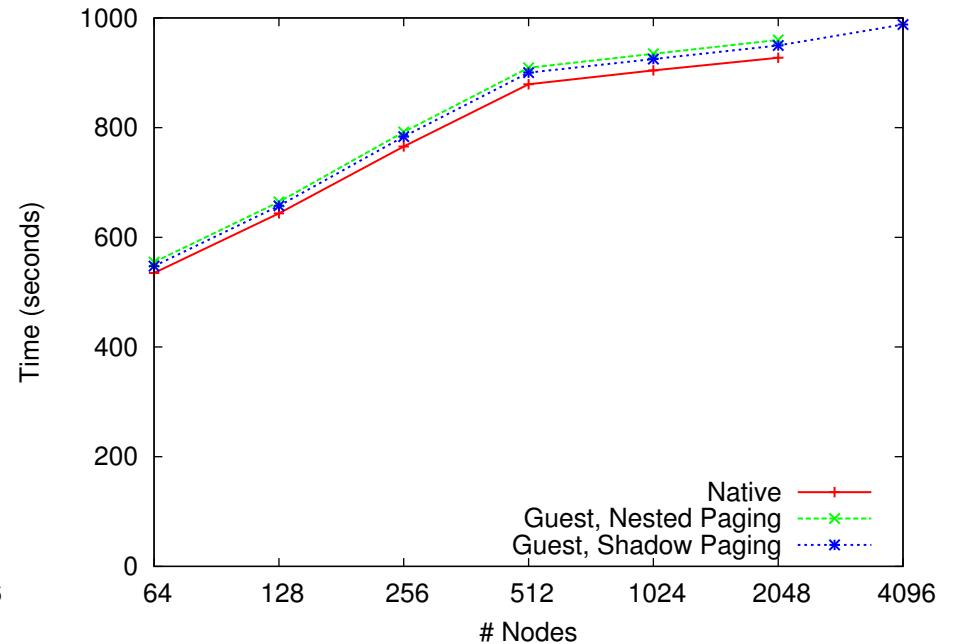


Application Results from Red Storm Virtualization Experiments

CTH Hydrocode (SNL App)



Sage Hydrocode (LANL App)



**Measured < 5% virtualization
overhead for both applications**



Current Project: DOE/ASCR X-Stack

- Objective: Enable X-Stack research and HW/SW co-design for exascale systems by leveraging the virtualization capabilities in modern processors
- Desired Capabilities
 - Enable X-Stack researchers to run new OS stacks at scale on production ASCR systems
 - Test potential architectural innovations at scale as extensions to the virtual machine
 - Measure system performance across multiple hardware/software boundaries
- Example Research
 - Scalable virtualization, VM management tools on modern HPC systems
 - Integration with cycle-accurate simulation/large-scale emulation techniques
 - Explore novel techniques in the VMM, both proposed and potentially in collaboration with other X-Stack or Critical Tech. researchers
- Consortium of researchers from Univ. New Mexico, Northwestern University, Oak Ridge, and Sandia



Conclusion

- **Applying virtualization technology to HPC**
 - Compelling use cases, enable new capabilities
 - Manageable overheads even at scale
- **Next steps:**
 - Test more applications, better characterize overheads for different workload classes
 - Push vendors to incorporate virtualization support in production software stacks
 - Leverage virtualization in exascale research



Backup Slides

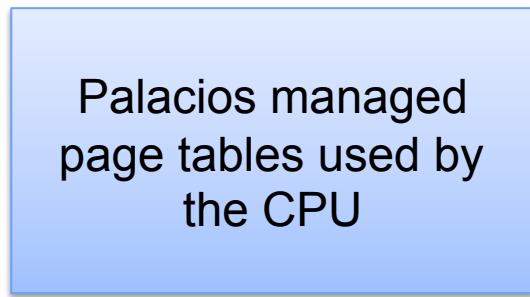


Shadow vs. Nested Paging

No Clear Winner

Shadow Paging

$O(N)$ memory accesses
per TLB miss

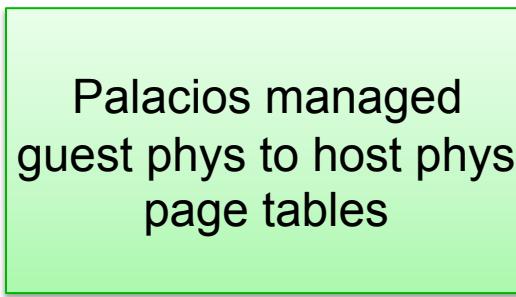


Page Faults

Page tables the guest OS thinks it is using

Nested Paging

$O(N^2)$ memory accesses
per TLB miss



CPU MMU

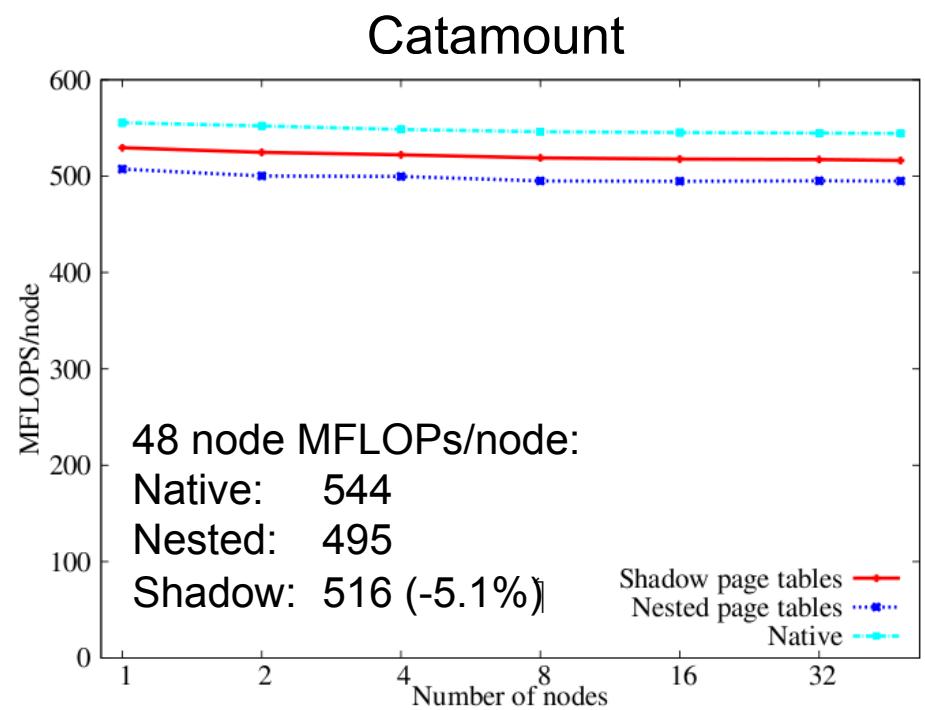
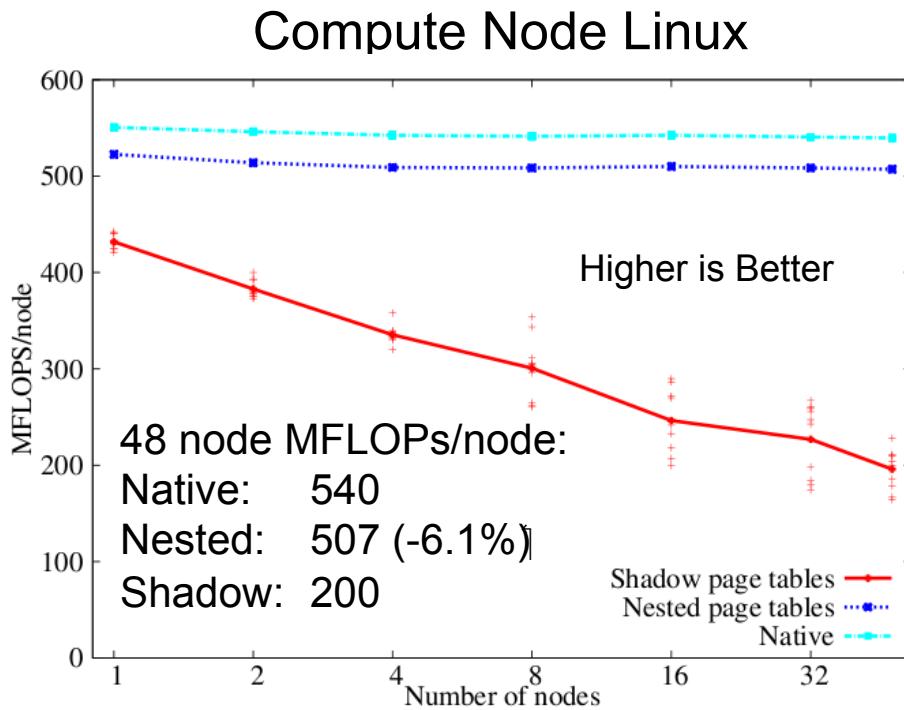
Guest OS managed guest virt to guest phys page tables





Memory Management Depends on Guest

HPCCG CG “Mini-application”



- Poor performance of shadow paging on CNL due to context switching. Could be partially avoided by adding page table caching to Palacios.
- Catamount is essentially doing no context switching, benefiting shadow paging (2n vs. n^2 page table depth issue)