

Exceptional service in the national interest



Photos placed in horizontal position
with even amount of white space
between photos and header

Sirocco Status

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Introduction

- L2 Milestone
 - POSIX compliance
 - Write performance
 - Read capability
- Components
 - Lock server
 - Data location service
 - Scatter/gather API

Lock Server

- POSIX requires atomicity
 - Namespace manipulation
 - File write access
- Related: TAMU work leveraging ASG semantics
- Short term: Lock service
 - Clients can obtain global leased lock
 - N-Party Locks
 - Attached metadata
 - Lock refreshing
 - Lock revocation
 - Knowledge of interspersed locking by others

- Sirocco servers may move data at will
 - Resilience, Capacity, Spite
- Related: Haiying's work, Zhiwei's work
- Short term: Unsophisticated data location
 - Data location depends on netgraph functionality
 - Netgraph does bootstrapping/overlay networking
 - A server broadcasts received requests, short-circuited by:
 - A request cache (keyed by <originating address | operation ID>)
 - TTL
 - Location information *only* is returned to client
 - <Record offset, number of records, location address, version> x n
 - Finite receive buffer yields need for extra information, i.e. “last known-good offset”
 - Catastrophic cancellation of location information

Scatter-Gather API

- Defined wrapper API over ASG API to perform scatter/gather reads
 - Similar to PVFS noncontiguous access, or xread/xwrite in libsysio
- Internally, uses the batching/transactional functionality that we defined for Sirocco
 - ```
ss_tx *tx = start_tx(...);
asg_op(..., tx);
tx_submit(tx);
tx_wait(tx);
```
- The batching made the implementation of this API quick and easy, and API makes certain operations within client quick and easy.
- Is it time to consider adding batching to ASG interface?

# Thanks

- Questions?

# Location: Catastrophic Cancellation

| Batch 1     | Batch 2     | Batch 3       | Batch 4       |
|-------------|-------------|---------------|---------------|
| [0, 1, 2]   | [1, 1, 2]   | [0, 128, 1]   | [0, 128, 3]   |
| [2, 1, 2]   | [3, 1, 2]   | [128, 128, 3] | [128, 128, 1] |
| ...         |             | ...           |               |
| [126, 1, 2] | [127, 1, 2] |               |               |

- [x, y, z] is an extent with offset=x, nrecords=y, and updateID=z
- Merging from left to right, with only 128 slots available
- After merging batches 1 and 2, all slots are filled
- After merging batches <1...2> and 3, all slots are still filled
- After merging batches <1...3> and 4, two slots are filled
  - Correct data from batch 3 ([128, 128, 3]) wasn't "found" because of buffer space constraints
  - Incorrect data from batch for ([128, 128, 1]) was included because capacity was freed during merge
- Solution: Keep offset of known-correct data, keep minimum of all reported offsets through reduction