

Security Games with Limited Surveillance

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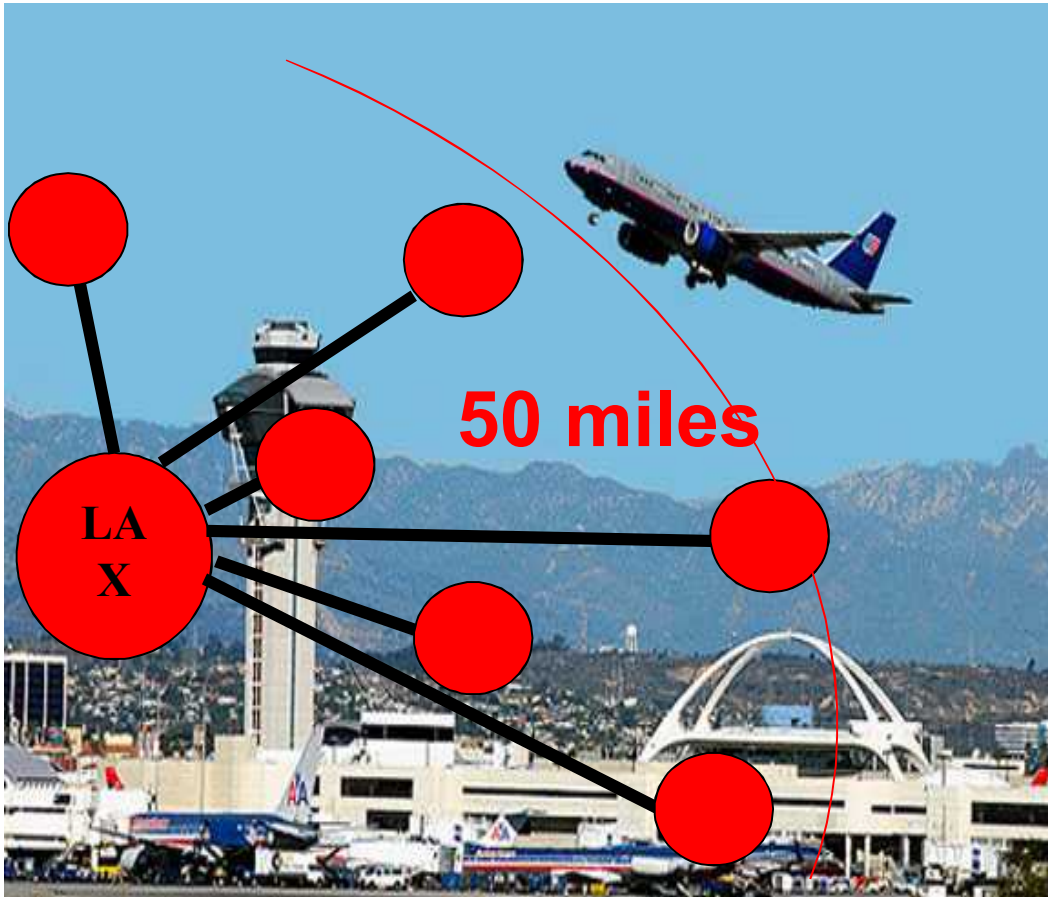
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Motivation: Game Theory for Security

- Limited security resources
- Adversary monitors defenses, exploits patterns



Game Theory: Bayesian Stackelberg Games

- Security allocation: (i) Target weights; (ii) Opponent reaction
- *Stackelberg*: Security forces commit first
- *Bayesian*: Uncertain adversary types
- *Optimal security allocation*: Weighted random
- **Strong Stackelberg Equilibrium (Bayesian)**
 - NP-hard



Adversary ↓



	Target #1	Target #2
Target #1	5, -3	-1, 1
Target #2	-5, 5	2, -1



Problem Statement

**How does the attacker
conduct surveillance and build belief ?**

- **Previous work:**

- *Perfect surveillance*

- *Imperfect surveillance*

- COBRA [Pita et al. 2010], RECON [Yin et al. 2011]

- Rely on **hand-tuned** parameters, cannot reason about surveillance **cost**



Our Contributions

- **Security game with limited observation model**
- **Analysis**
- **Algorithms**
 - **Compute optimal defender strategy**
 - **Estimate attacker observation length**
- **Experiments**

Security Games with Limited Surveillance



will make τ observations



decide mixed strategy X



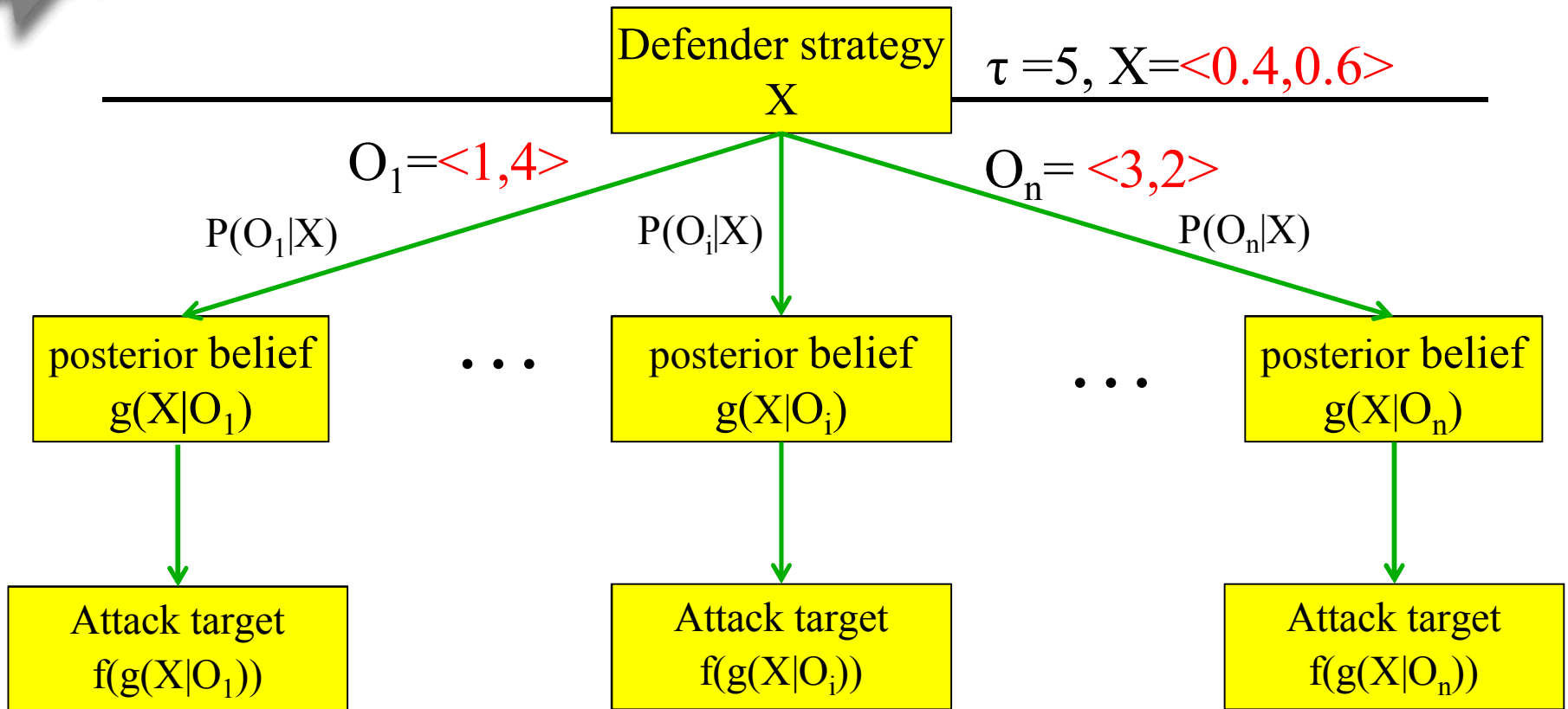
observe τ pure strategies O , update belief, attack the best target



Illustrating Example

- **The scenario:**
 - 3 targets t_1, t_2, t_3 : 3 attacker pure strategies
 - 1 resource: 3 defender pure strategies A_1, A_2, A_3
- **Defender strategy $X = \langle x_1, x_2, x_3 \rangle$**
 - stationary
- **Attacker has a prior belief**
 - e.g., uniform distribution
- **Attacker makes τ (e.g., 5) observations**
 - e.g., $O = A_1, A_2, A_1, A_2, A_3$
 - compact representation: $O = \langle o_1, o_2, o_3 \rangle = \langle 2, 2, 1 \rangle$
- **Best response based on posterior belief**

Defender's Optimization Problem



$$\max_X \sum_O P(O|X) U(X, f(g(X|O)))$$

Optimal Defender Strategy

$$\max \sum_{\mathbf{o} \in \mathcal{O}_\tau} \frac{\tau!}{\prod_{A \in \mathcal{A}} o_A!} \prod_{A \in \mathcal{A}} (x_A)^{o_A} d^{\mathbf{o}}$$

$$\text{s.t.} \quad x_A \in [0, 1] \quad \forall A \in \mathcal{A}$$

$$\sum_{A \in \mathcal{A}} x_A = 1$$

$$c_i = \sum_{A \in \mathcal{A}} x_A A_i \quad \forall i \in T$$

$$d^{\mathbf{o}} = c_{\psi(\mathbf{o})} (R_{\psi(\mathbf{o})}^d - P_{\psi(\mathbf{o})}^d) + P_{\psi(\mathbf{o})}^d \quad \forall \mathbf{o} \in \mathcal{O}_\tau$$



Analysis

Defender utility is always monotone non-increasing with larger τ



More targets protected with larger τ



Safe targets should always not be protected



Unsafe targets should always be protected



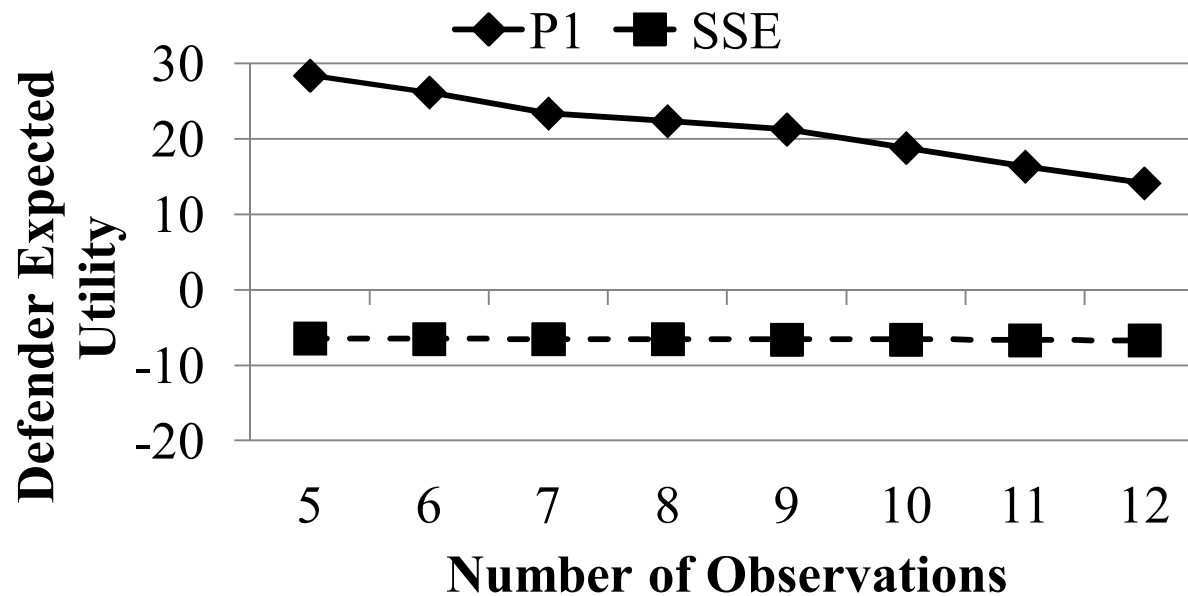
Estimate Attacker Observation Length

$$\operatorname{argmax}_{\tau} \left(\sum_{\mathbf{o} \in \mathcal{O}_{\tau}} \frac{\tau!}{\prod_{A \in \mathcal{A}} o_A!} \prod_{A \in \mathcal{A}} x_A^*(\tau)^{o_A} k^{\mathbf{o}} - \lambda \cdot \tau \right)$$

$\mathbf{x}^*(\tau)$: defender's optimal strategy

$k^{\mathbf{o}}$: attacker's utility when he observes \mathbf{o} and the defender plays $\mathbf{x}^*(\tau)$ strategy

Experimental Results





Conclusions

- **Contributions**

- Security game with limited observation model
- Analysis
- Algorithm for computing optimal defender strategy
- Algorithm for estimating observation length

- **Future work**

- Stopping problem
- Human-subject experiments
- Scalability