



SAND2019-2037C

# Performance Portable SIMD Approach Implementing Block Line Solver For Coupled PDEs

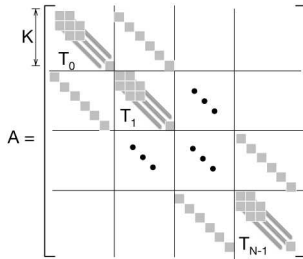
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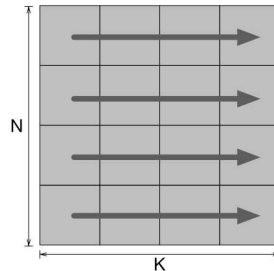
SIAM CSE, Feb 25 - Mar 1, 2019

# Line Implicit Solver

- Consider a block sparse system of equations  $Ax = b$  arising from coupled PDEs.
- Lines are usually formed in the boundary layer to resolve shocks.
- Small blocks represent interactions among a group of variables e.g., velocity, pressure, density and chemical species.
- Default solver for SPARC (Sandia production-level CFD code)



Linear system of equations



Problem domain and extracted lines

# Line Implicit Solver

- By splitting  $A = M - S$ , we obtain

$$(M - S)x = b$$

$$Mx = b + Sx$$

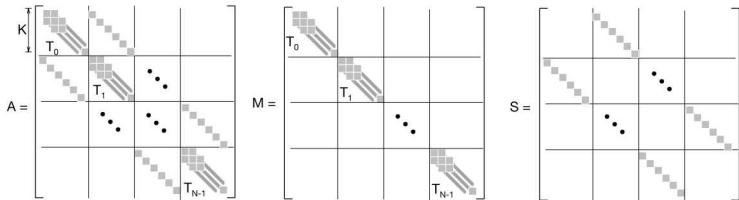
$$x = M^{-1}(b + Sx) = x + M^{-1}(b - Ax),$$

where  $M$  is a set of **block tridiagonal matrices** corresponding to the extracted lines of elements.

- Solve this iteratively

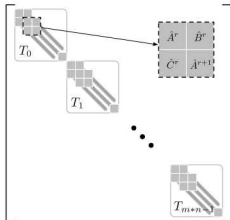
$$x^{k+1} = M^{-1}(b + Sx^k).$$

- This procedure requires solutions of **many block tridiagonal systems** and **efficient block sparse matrix vector multiplication**.



# Solving Many Tridiagonal Systems

- A fixed number of iterations are applied  $x^{k+1} = M^{-1}(b + Sx^k)$ .
- $M$  is factorized once per solution (or each non-linear iteration) and applied multiple times.
- Typical blocksizes  $b$  are 3, 5, 9 and 15, which are related to scientific applications e.g., elasticity, ideal gas and reactive fluid problems.
- Limit memory usage up to 16 GB i.e., MCDRAM on KNL and GPU device memory.
- A typical local problem size in a 3D domain ( $m \times n \times k$ ) is selected as  $128 \times 128 \times 128$  for  $b = 3, 5$  and  $64 \times 64 \times 128$  for  $b = 10, 15$ .
- Batch parallelism ( $m \times n$ ) is considered for solving block tridiagonal systems.



```

1 for  $T$  in  $\{T_0, T_1, \dots, T_{m \times n - 1}\}$  do in parallel
2   for  $r \leftarrow 0$  to  $k - 2$  do
3      $\hat{A}^r := LU(\hat{A}^r)$ ;
4      $\hat{B}^r := L^{-1} \hat{B}^r$ ;
5      $\hat{C}^r := \hat{C}^r U^{-1}$ ;
6      $\hat{A}^{r+1} := \hat{A}^{r+1} - \hat{C}^r \hat{B}^r$ ;
7      $\hat{A}^{k-1} := LU(\hat{A}^{k-1})$ ;

```

*Initialization of the line smoother*

# Efficient Implementation using Compact Data Layout (SIMD)

## Problems using vendor libraries

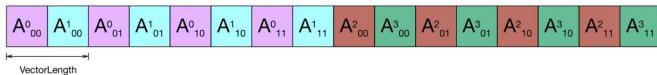
- Standard BLAS/LAPACK are not optimized for solving small problems.
- Using batched BLAS/LAPACK APIs, no data locality is exploited in a sequence of batch operations.

## Our approach: Compact Data Layout

- A new interleaved data layout for batched BLAS/LAPACK for efficient vectorization on small problems.
- Solve a small number of problems in parallel using SIMD instructions.
- Array of blocks is represented by a multi-dimensional array.

$A^0_{00}$	$A^0_{01}$	$A^1_{00}$	$A^1_{01}$	$A^2_{00}$	$A^2_{01}$	$A^3_{00}$	$A^3_{01}$
$A^0_{10}$	$A^0_{11}$	$A^1_{10}$	$A^1_{11}$	$A^2_{10}$	$A^2_{11}$	$A^3_{10}$	$A^3_{11}$

- Compact layout packs data across batch instances with respect to SIMD length.



# SIMD Approach To Solve Small Dense Problems

- To vectorize across batch instances, we use packed vectors as computing unit instead scalar.
- By overloading arithmetic operators, the scalar BLAS are reused and naturally vectorized.

```

template<typename T, int vl>
struct Vector<T,vl> {
    T _data[vl];
};

Vector<T,vl>
operator+(Vector<T,vl> a, Vector<T,vl> b) {
    Vector<T,vl> r_val;
    for (int i=0;i<vl;++i)
        r_val._data[i] = a._data[i] + b._data[i];
    return r_val;
}

```

- Multidimensional array:

```

typedef double scalar_type;
typedef Vector<scalar_type,VectorLength> vector_type;

// LayoutRight: most right consecutive index is stored in the memory contiguously
Kokkos::View<vector_type***,LayoutRight> Av('A', N, Blk, Blk);
// Reinterpretation of rank 3 vector_type view with scalar rank 4 view
Kokkos::View<scalar_type****,LayoutRight> As(Av.data(), N, Blk, Blk, VectorLength);

```

- What about GPUs ?

# Implicit and Explicit (Warp) Vectorizations

## SIMD instructions on CPUs

- A single thread exploits data-level parallelism.
- Vector length is given from hardware specific vector instruction set e.g., AVX512.
- A compiler is responsible for vectorization.

## Warp-based (explicit) vectorization on GPUs

- Warp: a set of concurrent threads executing the same instruction.
- Each thread in a warp processes scalar instructions.
- Vector instructions are dynamically formed by a warp.
- Coalesced memory access by threads is essential for efficient memory operations.
- GPUs have 128 bit memory instruction set (double2 – internal vector type)
  - A benefit using double2 is to reduce the number of memory transactions.

```
using namespace Kokkos;
// Suppose that N = 64, b = 4

// Default values are pre-defined in KokkosBatched but a user can
// control the vector size for CUDA.
// vl: Host 8, Cuda 16, Vector length used for compact data format.
// il: Host 8, Cuda 2, Internal vector length used in a functor.
//
//          Using internal vector type, we make sure
//          using 128bit memory instructions.
constexpr int vl = DefaultVectorLength<exec_space>::value;
constexpr int il = DefaultInternalVectorLength<exec_space>::value;

typedef double scalar_type
typedef Vector<scalar_type,vl> vector_type;
typedef Vector<scalar_type,il> internal_vector_type;

// Av: Host 8x4x4, Cuda 4x4x4, Rank 3 Multi-dimensional array
//          of vector_type.
View<vector_type*** ,LayoutRight> Av('A', N/vl, b, b);
```

*Example: Set identity of  $N$  matrices  $b \times b$*



# Kokkos Unified Programming Approach

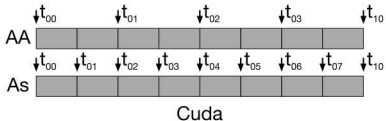
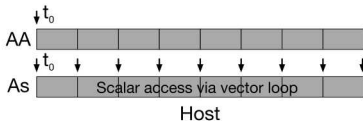
```
using namespace Kokkos;
// Suppose that N = 64, b = 4

// Av: Host 8x4x4, Cuda 4x3x3, Rank 3 Multi-dimensional array
//      of vector_type.
View<vector_type***, LayoutRight> Av('A', N/vl, b, b);

// Reinterpret the rank 3 vector view with the rank 4 internal vector type.
// AA: Host 8x4x4x1, Cuda 4x4x4x8
View<internal_vector_type****, LayoutRight> AA(Av.data(), N/vl, b, b, vl/il);

// One can also reinterpret the view with rank 4 scalar type if
// accessing individual scalar entries is necessary.
// As: Host 8x4x4x8, Cuda 4x4x4x16
View<scalar_type****, LayoutRight> As(Av.data(), N/vl, b, b, vl);

// For this example, we use a view with internal vector type (AA).
```



*Example: Set identity of  $N$  matrices  $b \times b$ ; figure describes different data access patterns for a vector length 8.*

```
using namespace Kokkos;
// Suppose that N = 64, b = 4

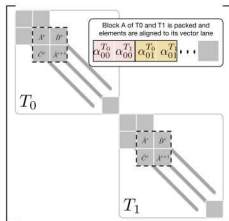
// Reinterpret the rank 3 vector view with the rank 4 internal vector type.
// AA: Host 8x4x4x1, Cuda 4x4x4x8
View<internal_vector_type****,LayoutRight> AA(Av.data(), N/vl, b, b, vl/il);

// Host policy(8,1,1), Cuda policy(gridSize=8,blockDim.y=8, blockDim.x=4)
// You can consider league_size equivalent to the number of independent works.
TeamPolicy<exec_space> policy(league_size = N/vl,
                             team_size = AUTO,
                             vector_loop_size = vl/il);
parallel_for(policy, KOKKOS_LAMBDA(const member_type member) {
    int p = member.league_rank();
    // CUDA: for (int v=threadIdx.x;v<vector_loop_size;v+=blockDim.x)
    parallel_for(ThreadVectorRange(member, vector_loop_size), [&](int v) {
        // Host A 4x4, Internal vector length 8.
        // Cuda A 4x4, Internal vector length 2.
        // Consecutive indexing of v allows perfectly coalesced
        // memory access along 16 doubles.
        auto A = subview(AA, p, ALL, ALL, v);
        // CUDA: for (int ij=threadIdx.y;ij<b*b;ij+=blockDim.y)
        Kokkos::parallel_for(TeamThreadRange(member, b*b), [&](int ij) {
            int i = ij/b, j = ij%b;
            AA(i,j) = i == j ? one : zero;
        }); }); });
```

*Example: Set identity of  $N$  matrices  $b \times b$*

# Line Implicit Solver Implementation

- KokkosKernels provides serial/team level dense linear algebra components so that users can compose their own batch operations.
- Solve phase of block tridiagonal matrices can be done in a similar fashion.



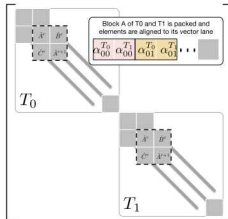
```

1  for a pair  $T$  in  $\{(T_0, T_1), (T_2, T_3), \dots, (T_{m \times n - 2}, T_{m \times n - 1})\}$  do in parallel
2      for  $r \leftarrow 0$  to  $k - 2$  do
3           $\hat{A}^r := LU(\hat{A}^r);$ 
4           $\hat{B}^r := L^{-1} \hat{B}^r;$ 
5           $\hat{C}^r := \hat{C}^r U^{-1};$ 
6           $\hat{A}^{r+1} := \hat{A}^{r+1} - \hat{C}^r \hat{B}^r;$ 
7           $\hat{A}^{k-1} := LU(\hat{A}^{k-1});$ 

```

# Line Implicit Solver Implementation

- KokkosKernels provides serial/team level dense linear algebra components so that users can compose their own batch operations.
- Solve phase of block tridiagonal matrices can be done in a similar fashion.



```
// Problem domain ( M x N ) x K, blocksize b
// - M x N : the number of block tridiagonals (concurrency)
// - K : the length of block tridiagonals (workload per team)
View<vector_type*****,LayoutRight> Av('A', M*N/vl, K, 3, b, b);
View<internal_vector_type*****,LayoutRight>
    AA(Av.data(), M*N/vl, K, 3, b, b, vl/il);
```

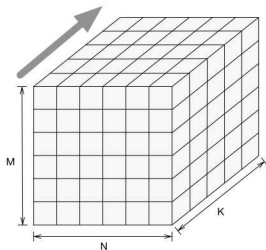
```
TeamPolicy<exec_space> policy(league_size = M*N/vl,
                              team_size = AUTO,
                              vector_loop_size = vl/il);
parallel_for(policy, KOKKOS_LAMBDA(const member_type member) {
    int p = member.league_rank();
    parallel_for(ThreadVectorRange(member, vector_loop_size), [&](int v) {
        for (int r=0;r<(K-1);++r) {
            auto A = Kokkos::subview(AA, p, r, 1, ALL(), ALL(), v);
            auto B = Kokkos::subview(AA, p, r, 2, ALL(), ALL(), v);
            auto C = Kokkos::subview(AA, p, r, 0, ALL(), ALL(), v);
            auto D = Kokkos::subview(AA, p, r+1, 1, ALL(), ALL(), v);
            TeamLU(member, A);
            TeamTrsm(member, 1.0, A, B);
            TeamTrsm(member, 1.0, A, C);
            TeamGemm(member, -1.0, C, B, 1.0, D);
        }
        auto A = Kokkos::subview(AA, p, K-1, 1, ALL(), ALL(), v);
        TeamLU(member, A);
    }); });
```

## Test Machines

- Intel Xeon Phi
- NVIDIA Tesla V100

## Test Problems

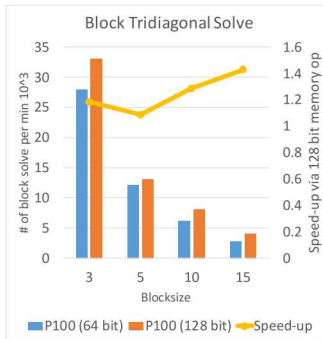
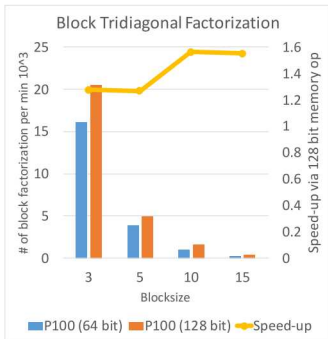
- With limited device memory (16 GB), the following test problems are evaluated.



Blocksize	$M \times N \times K$
3	$128 \times 128 \times 128$
5	$128 \times 128 \times 128$
10	$64 \times 64 \times 128$
15	$64 \times 64 \times 128$

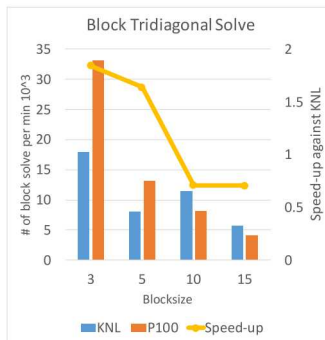
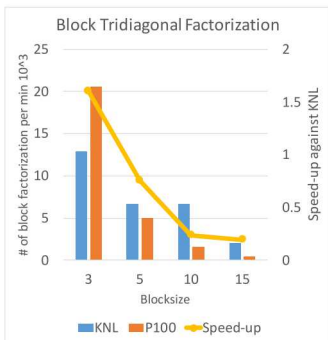
- Both codes shows 100% memory store/load efficiency and performance is limited by L2.
- Bigger blocksize (more computational intensity) shows more benefits from 128 bit memory instructions.

Team (blockDim)	$2 \times 16$	$4 \times 8$
Global Store/Load Throughput (GB/s)	215.31 / 641.38	287.68 / 762.08
L2 Throughput Writes/Reads (GB/s)	215.31 / 530.37	287.68 / 677.41
Global Store/Load Transactions Per Request (Byte)	7.69 / 31.22	14.27 / 60.81



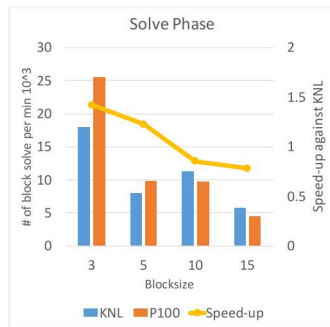
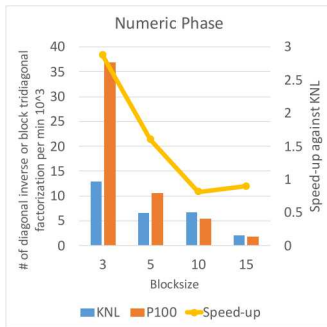
- Block tridiagonal factorization/solves are inherently sequential.
- With increasing blocksize, the number of concurrent work (tridiagonals) is reduced while the workload per team increases.
- Need more parallelism for bigger problems.

Cube Size( $M \times N \times K$ )	$128 \times 128 \times 128$	$64 \times 64 \times 128$
Multiprocessor Activity (%)	98.31	95.75
Achieved Occupancy	0.29	0.07



# Block Jacobi Iterations for Solving Block Tridiagonals

- In transient flow, application (CFD) uses tiny time steps with relaxed error threshold - more nonlinear iterations and less number of iterations for solves.
- Apply a small number of Jacobi iterations instead of direct solutions (factorize, forward/backward solves).
- Improve performance on numeric phase (inverse diagonals).
- With increasing data movement during Jacobi iterations, there is not much benefit in the solve phase.



*Comparison of Jacobi block tridiagonal solves (threshold  $10^{-4}$ ) on P100 against direct block tridiagonal solves on KNL*



- Demonstrated performance portable SIMD approach for line implicit solver using Kokkos framework.
  - Prepare data for multidimensional array of SIMD type.
  - Reinterpret the multidimensional array based on execution space.
  - Construct `Kokkos::TeamPolicy` accordingly to exploit both compiler vectorization and dynamic vectorization.
- Improved GPU performance (10 - 50 %) by using 128 bit memory instructions (`double2` or `float4`); performance improvement is higher when computational intensity increases.
- Compared to highly optimized KNL code (same code), GPUs perform better for small block sizes (3 and 5)
- Due to device memory constraint (lowering parallelism), it is difficult to use a GPU efficiently for bigger block sizes (10 and 15).
- Easy solution: a GPU with a bigger device memory (32 or 64 GB).