

# Modeling Information Spread Through Large Networks

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# Outline

- Motivation
  - Predict information spread through large networks
- Monte Carlo methods
  - Independent cascade (IC)
- Deterministic methods
  - Belief propagation (BP)
- Extending BP to directed graphs
  - Directed belief propagation (DP)
- Applications

# Objective and Motivation

- Given a graph
  - Predict information spread
  - Identify the most influential nodes
- Applications
  - Business
  - Elections
  - National security
- Requirements
  - Estimate model parameters from real-world data
  - Results should match real-world data when available
  - Computationally efficient for large networks



# Monte Carlo Methods

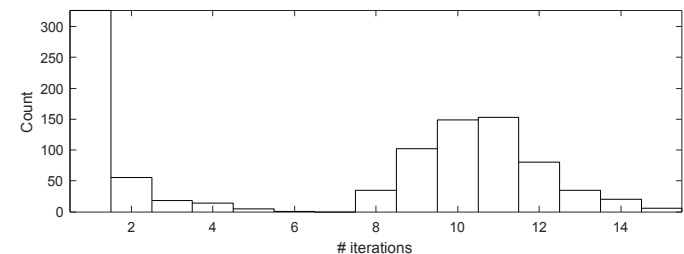
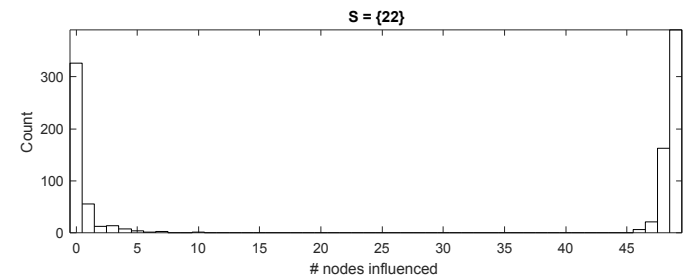
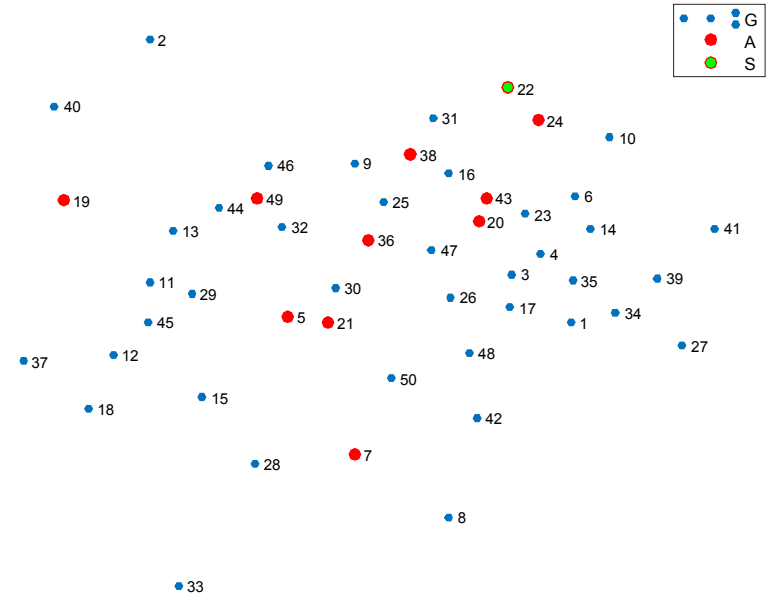
- Pros
  - Simple to understand and implement
  - Applicable to general graphs
  - Time-dependency is represented
  - A variety of metrics can be studied (e.g., joint distributions)
  - Accuracy can be estimated
- Cons
  - Convergence is slow
  - Approach can be infeasible for large graphs
- One approach: Independent cascade (IC)
  - Kempe et al. (2003)
  - Model parameters: influence probabilities

# Independent Cascade (IC)

- Let  $G(\mathcal{V}, \mathcal{E})$  be a graph with node set  $\mathcal{V} = \{1, \dots, n\}$  and edge set  $\mathcal{E} = \{(i, j), 1 \leq i, j \leq n\}$ .
- Each node  $i$  in the graph has a state  $X_i(t) \in \{0, 1\}$  at time  $t$ , where  $X_i(t) = 1$  and  $X_i(t) = 0$  denotes that node  $i$  is active and inactive. Node  $j$  in the network exhibits influence on its neighbors, that is, on the set of nodes  $\mathcal{N}(j) = \{i \in \mathcal{V}, i \neq j : (i, j) \in \mathcal{E}\}$ .
- We say that node  $i \in \mathcal{V}$  is “contagious” at time  $t$  if  $X_i(t) = 1$  and  $X_i(t - 1) = 0$ , that is, if it is active at time  $t$  but was inactive at time  $t - 1$ .
- Assuming node  $i$  is contagious at time  $t$ , it attempts to influence its neighbors  $j \in \mathcal{N}(i)$  with probability  $p(i, j) \in [0, 1]$ . If successful, node  $j$  becomes active at time  $t + 1$ ; otherwise node  $j$  is not active. At time  $t + 1$ , node  $i$  is no longer contagious and is unable to influence any more nodes.
- To start the spread of influence through  $G$ , there must be an initial condition. Let  $S \subset \mathcal{V}$  be a collection of “seed” nodes that are assumed contagious at time  $t = 0$ . At each time step, the contagious nodes are allowed to activate their neighboring nodes with probabilities  $p(i, j) \in [0, 1]$ . The simulation ends when no additional node become contagious or all nodes have been activated.

# IC Example

- E-R graph with  $n = 50$  nodes and  $p(i,j) = 0.1$ . Seed node  $S = 22$  (green).
- At convergence, 11 (red) nodes influenced; 39 (blue) nodes remain inactive
- IC model is repeated 1,000 times with identical seed node
- Distribution of number of nodes influenced (top)
- Number of iterations of IC needed for convergence



# A Deterministic Method

- Belief propagation (BP)
  - J. Pearl (1982), (Yedidia, 2001)
  - Adaptable to iterative solver for linear systems
- Pros
  - Solve for marginal distributions directly
  - Minimal computational effort
- Cons
  - More difficult to understand and implement
  - Strict assumptions on graph structure (e.g., polytree)
  - Used in general graphs (“loopy BP”); consequences of this are largely unknown
  - Not really appropriate for directed graphs
- Model parameters
  - Initial beliefs and potential functions

# Belief Propagation (BP)

- Let  $G(\mathcal{V}, \mathcal{E})$  be a graph with node set  $\mathcal{V} = \{1, \dots, n\}$  and edge set  $\mathcal{E} = \{(i, j), 1 \leq i, j \leq n\}$ . We will assume  $G$  contains no directed cycles, that is, there exist no closed paths in the graph.
- Let  $\phi_i(x_i)$  quantify node  $i$ 's belief it has value  $x_i$ ; this quantity represents previous knowledge about node  $i$  by, for example, evidence gained from observations. Further, let  $\psi_{ij}(x_i, x_j)$  quantify the probability that node  $j$  believes it has value  $x_j$  given that node  $i$  believes it has value  $x_i$ . Then

$$m_{ij}(x_j) = \sum_{x_i} \phi_i(x_i) \psi_{ij}(x_i, x_j) \prod_{k \in \mathcal{N}(i) \setminus \{j\}} m_{ki}(x_i)$$

denotes a message passed from node  $i$  to node  $j$  describing node  $i$ 's belief that node  $j$  has value  $x_j$ , where  $\mathcal{N}(j) = \{i \in \mathcal{V}, i \neq j : (i, j) \in \mathcal{E}\}$  defines the neighborhood of node  $j$ .

- The marginal probability at node  $k$  is then the product of all incoming messages at that node multiplied by its prior belief, i.e.,

$$p(x_k) \propto \phi_k(x_k) \prod_{j \in \mathcal{N}(k)} m_{jk}(x_k)$$

# Some Additional Details

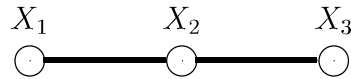
- Assume the state of each node is  $x_i \in \{0, 1\}$ .
- $\phi_i(x_i)$  is a  $2 \times 1$  vector representing the initial belief that node  $i$  is in state 0 or 1.
  - If the state of node  $i$  is known, this is generally  $(0.999, 0.001)^T$ .
  - If the state of node  $i$  is unknown, this is generally set to  $(0.5, 0.5)^T$ .
- $\psi_{ij}(x_i, x_j)$  is a  $2 \times 2$  matrix representing the probability that node  $j$  believes it has value  $x_j$  given that node  $i$  believes it has value  $x_i$ , for all possible values of  $x_i$  and  $x_j$ .

$$\text{Homophily: } \begin{bmatrix} 0.99 & 0.01 \\ 0.01 & 0.99 \end{bmatrix} \quad \text{Heterophily: } \begin{bmatrix} 0.01 & 0.99 \\ 0.99 & 0.01 \end{bmatrix}$$

- $\prod_{k \in \mathcal{N}(i) \setminus \{j\}}$  ensures that node  $j$ 's message to node  $i$  don't come right back.
- Message passing repeats until convergence is attained.

# BP Example

- Consider a simple graph with 3 nodes, 2 edges



- We assume random variables  $x_1$ ,  $x_2$ , and  $x_3$  each take values in  $\{0, 1\}$  with potential functions

$$\psi_{1,2}(x_1, x_2) = \begin{pmatrix} 1 & 0.9 \\ 0.9 & 1 \end{pmatrix} \quad \text{and} \quad \psi_{2,3}(x_2, x_3) = \begin{pmatrix} 0.1 & 1 \\ 1 & 0.1 \end{pmatrix}$$

and initial beliefs

$$\phi_1(x_1) = \phi_3(x_3) = \begin{pmatrix} 0.5 \\ 0.5 \end{pmatrix} \quad \text{and} \quad \phi_2(x_2) = \begin{pmatrix} 1 \\ 0.1 \end{pmatrix}$$

- Node 1 has a mild preference to be like node 2
- Node 3 strongly wants to be the opposite of node 2
- Node 2 has strong initial belief that it is in state 0
- Nodes 1 and 3 initially believe they are in state 0 or 1 with equal probability

# Messages

$$\begin{aligned} m_{12}(x_2) &= \sum_{x_1} \phi_1(x_1) \psi_{1,2}(x_1, x_2) \prod_{k \in \mathcal{N}(1) \setminus \{2\}} m_{k1}(x_1) = \sum_{x_1} \phi_1(x_1) \psi_{1,2}(x_1, x_2) \\ &= 0.5 \begin{pmatrix} 1 \\ 0.9 \end{pmatrix} + 0.5 \begin{pmatrix} 0.9 \\ 1 \end{pmatrix} = \begin{pmatrix} 0.95 \\ 0.95 \end{pmatrix} \end{aligned}$$

$$\begin{aligned} m_{32}(x_2) &= \sum_{x_3} \phi_3(x_3) \psi_{3,2}(x_3, x_2) \prod_{k \in \mathcal{N}(3) \setminus \{2\}} m_{k3}(x_3) = \sum_{x_3} \phi_3(x_3) \psi_{3,2}(x_3, x_2) \\ &= 0.5 \begin{pmatrix} 0.1 \\ 1 \end{pmatrix} + 0.5 \begin{pmatrix} 1 \\ 0.1 \end{pmatrix} = \begin{pmatrix} 0.75 \\ 0.75 \end{pmatrix} \end{aligned}$$

$$\begin{aligned} m_{21}(x_1) &= \sum_{x_2} \phi_2(x_2) \psi_{2,1}(x_2, x_1) \prod_{k \in \mathcal{N}(2) \setminus \{1\}} m_{k2}(x_2) = \sum_{x_2} \phi_2(x_2) \psi_{2,1}(x_2, x_1) m_{32}(x_2) \\ &= 1.0 \begin{pmatrix} 1 \\ 0.9 \end{pmatrix} (0.75) + 0.1 \begin{pmatrix} 0.9 \\ 1 \end{pmatrix} (0.75) = \begin{pmatrix} 0.8175 \\ 0.75 \end{pmatrix} \end{aligned}$$

$$\begin{aligned} m_{23}(x_3) &= \sum_{x_2} \phi_2(x_2) \psi_{2,3}(x_2, x_3) \prod_{k \in \mathcal{N}(2) \setminus \{3\}} m_{k2}(x_2) = \sum_{x_2} \phi_2(x_2) \psi_{2,3}(x_2, x_3) m_{12}(x_2) \\ &= 1.0 \begin{pmatrix} 0.1 \\ 1 \end{pmatrix} (0.95) + 0.1 \begin{pmatrix} 1 \\ 0.1 \end{pmatrix} (0.95) = \begin{pmatrix} 0.19 \\ 0.9595 \end{pmatrix} \end{aligned}$$

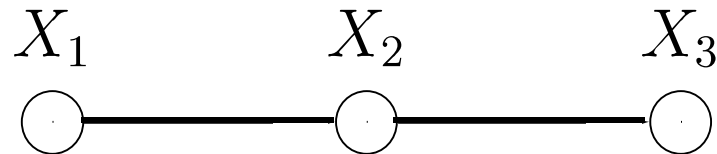
# Final Result

- Marginal probabilities for the state of each node

$$p(x_1) \propto \phi_1(x_1) m_{21}(x_1) = \begin{pmatrix} 0.40875 \\ 0.375 \end{pmatrix} \implies p(x_1) \approx \begin{pmatrix} 0.522 \\ 0.478 \end{pmatrix}$$

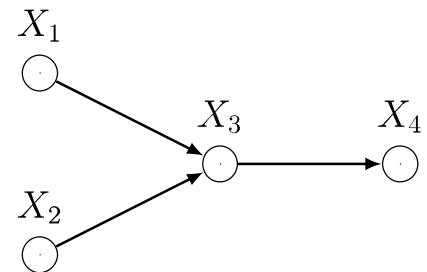
$$p(x_2) \propto \phi_2(x_2) m_{12}(x_2) m_{32}(x_2) = \begin{pmatrix} 0.7125 \\ 0.07125 \end{pmatrix} \implies p(x_2) \approx \begin{pmatrix} 0.909 \\ 0.091 \end{pmatrix}$$

$$p(x_3) \propto \phi_3(x_3) m_{23}(x_3) = \begin{pmatrix} 0.095 \\ 0.47975 \end{pmatrix} \implies p(x_3) \approx \begin{pmatrix} 0.165 \\ 0.835 \end{pmatrix}$$



# BP Extended to Digraphs

- A modification of Belief Propagation that preserves directed influence
  - BP passes update messages in both directions along edges in a graph
  - In directed propagation, messages pass only downstream
- Requires per-node and per-edge functions
  - Per-node may be learned from real-world data
  - Per-edge based on node in-degree
- Results in each node's likelihood of adoption
  - Diffusion score is sum of all nodes' likelihoods
- Implementation details in the paper
  - Available at <https://github.com/algorithmfoundry/Foundry>
  - Details on BP (Yedidia, 2001)



# Directed BP – Standard BP on a Collection of Subgraphs

- Let  $G = (\mathcal{V}, \mathcal{E})$  be a directed tree or polytree with node set  $\mathcal{V} = \{1, \dots, n\}$  and edge set  $\mathcal{E} = \{(i, j) : 1 \leq i, j \leq n\}$ .
- Suppose a directed path of two or more nodes exists from node  $u \in \mathcal{V}$  to node  $v \in \mathcal{V}$  which we denote by  $u \rightarrow v$ ; because  $G$  is assumed a polytree we know each  $u \rightarrow v$  path is unique.
- We define a collection of subgraphs  $G_k = (\mathcal{V}_k, \mathcal{E}_k) \subset G$ ,  $k = 1, \dots, n$ , where

$$\mathcal{V}_k = \{k\} \cup \{u \in \mathcal{V} : \text{there exists a path } u \rightarrow k\}$$

$$\mathcal{E}_k = \{(i, j) : i, j \in \mathcal{V}_k \text{ and } (i, j) \in \mathcal{E}\}$$

denote the node set and edge set, respectively, of subgraph  $G_k$ .

# More on DP

- Then

$$f_k(x_k) = \sum_{x_j: j \in \mathcal{V}_k \setminus \{k\}} p(\mathbf{x}_{\mathcal{V}_k})$$

is the marginal pdf of node  $k$  with respect to graph  $G_k$ , which can be related to the marginal pdf with respect to the original graph  $G$  via

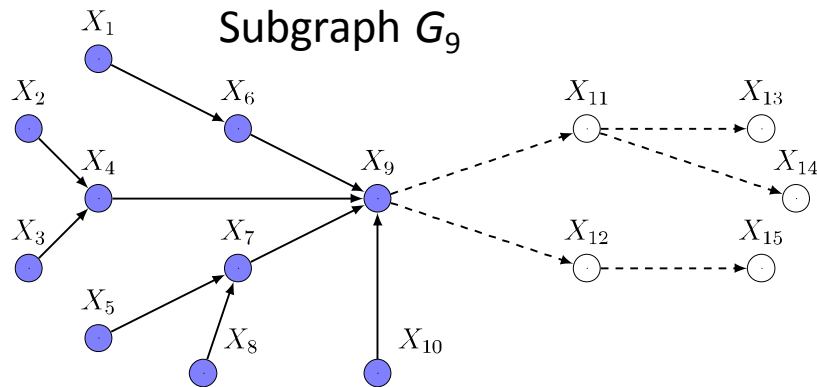
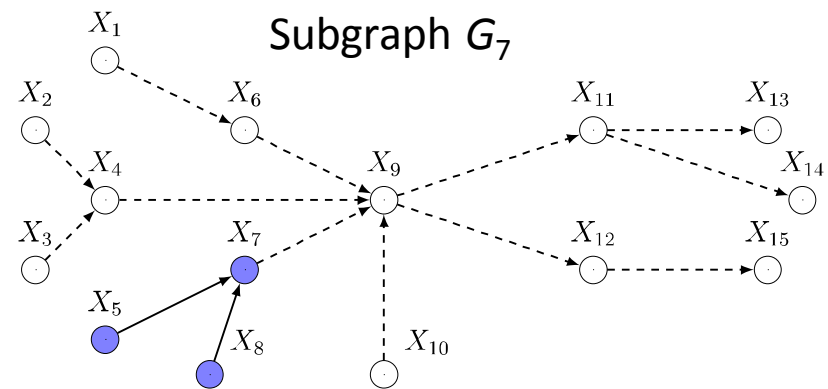
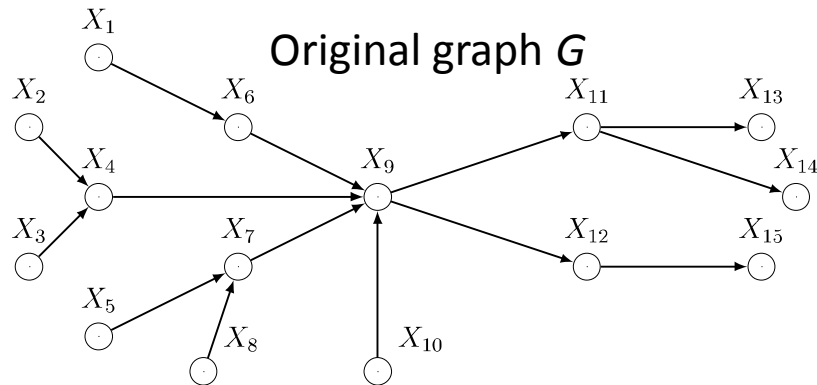
$$p(x_k) = f_k(x_k) \sum_{x_j: j \in \mathcal{V} \setminus \mathcal{V}_k \setminus \{k\}} p(\mathbf{x}_{\mathcal{V} \setminus \mathcal{V}_k})$$

- Then, assuming nodes  $i, j \in \mathcal{V}_k$ ,

$$m_{ij}(x_j) = \sum_{x_i} \phi_i(x_i) \psi_{i,j}(x_i, x_j) \prod_{k \in (\mathcal{N}(i) \cap \mathcal{V}_k) \setminus \{j\}} m_{ki}(x_i)$$

denotes a message passed from node  $i$  to node  $j$  on subgraph  $G_k$  describing node  $i$ 's belief that node  $j$  has value  $x_j$ . Only those nodes that can have influence on node  $j$  are considered; this is equivalent to computing only the “forward messages” to node  $j$ , and ignoring any “backwards messages”.

# An Example



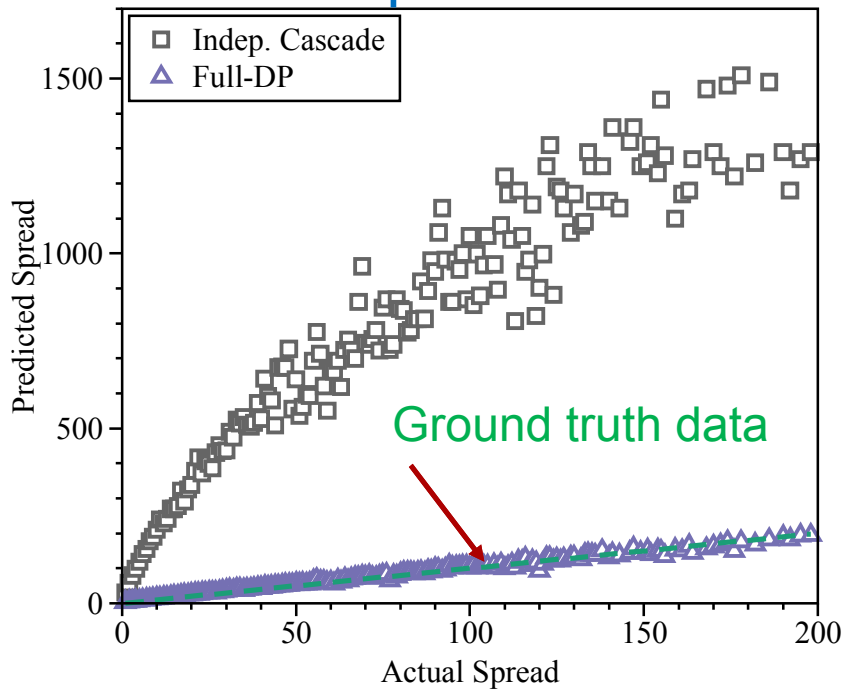
# Application #1 – Influence Spread

- Data sets
  - Flixster movie review propagation
    - 800K nodes; 12M edges
  - Epinions product review propagation
    - 18K nodes; 1.2M edges
  
- Diffusion algorithms
  - Independent cascade (IC)
  - Directed belief propagation (DP)
    1. Degree-weight per-edge; learned per-node (Full-DP)
    2. Degree-weight per-edge; constant per-node (Edge-DP)
    3. Constant-weight per-edge; learned per-node (Node-DP)

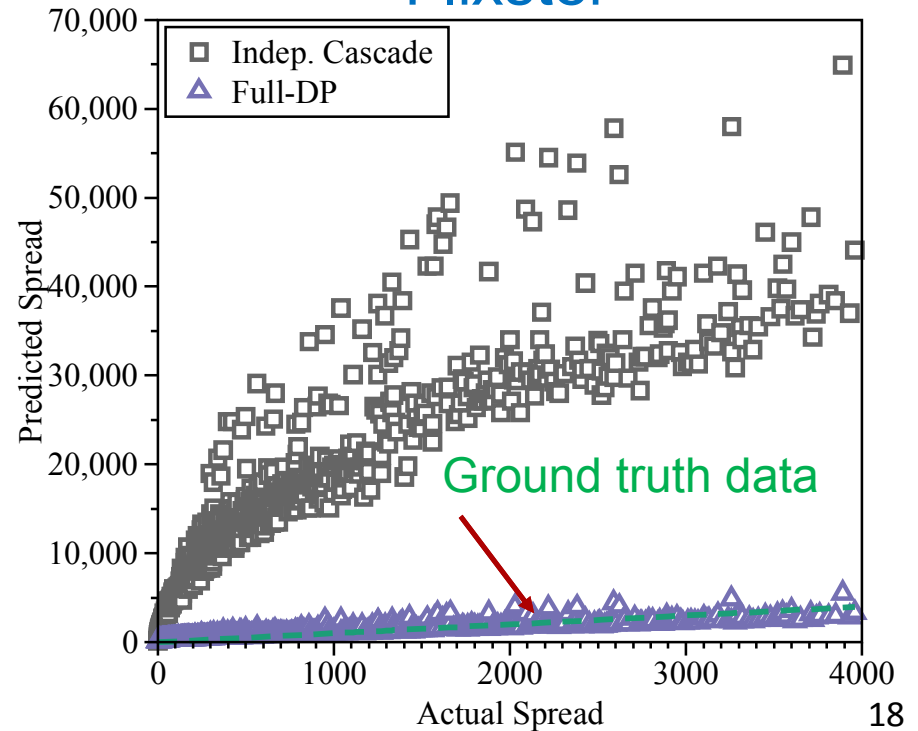
# Predicting Influence Spread (1 of 2)

- Directed propagation (DP) matches real-world data fairly well
- Independent cascade (IC) does not

## Epinions

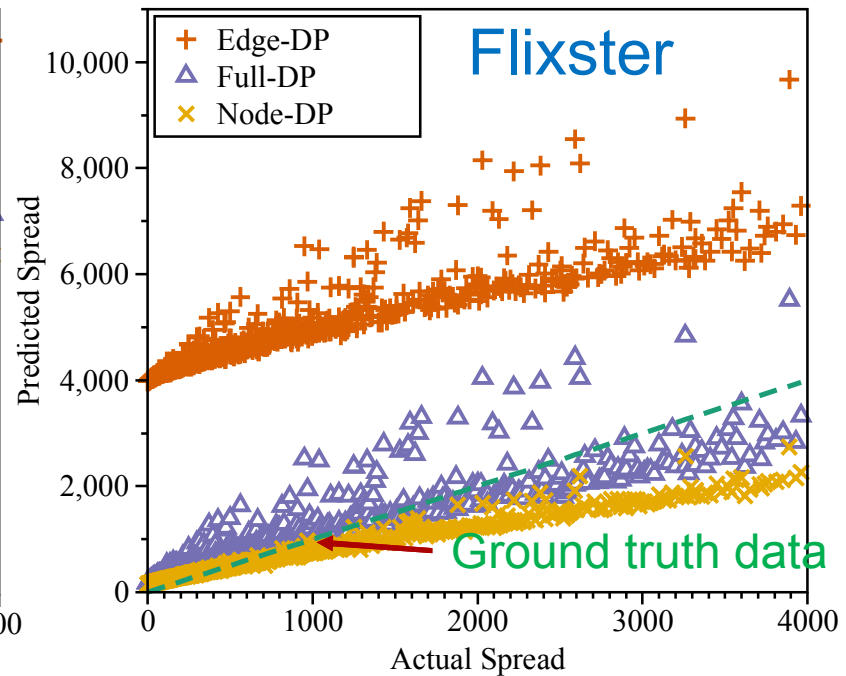
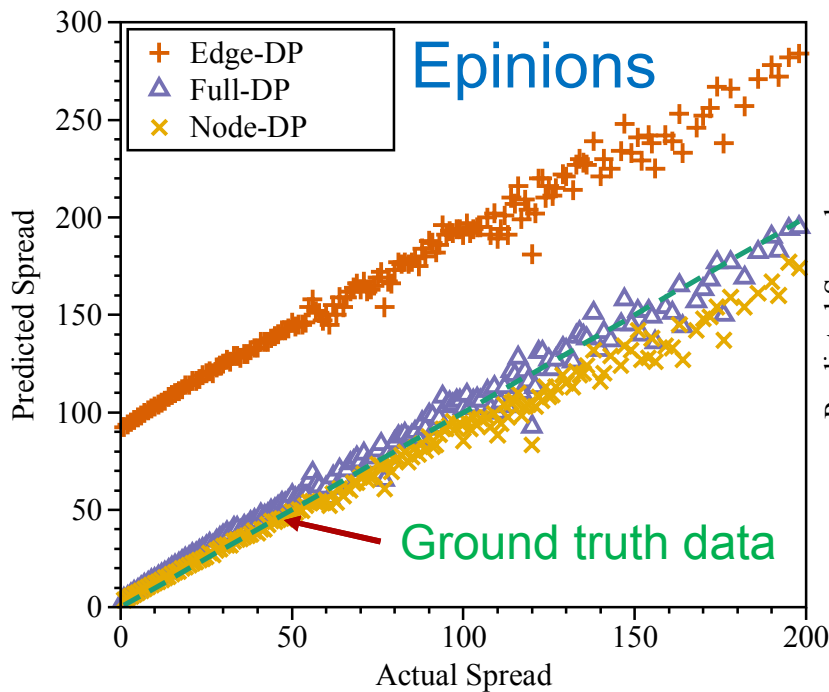


## Flixster



# Predicting Influence Spread (2 of 2)

- Three versions of DP
  1. Degree-weight per-edge; learned per-node (Full-DP)
  2. Degree-weight per-edge; constant per-node (Edge-DP)
  3. Constant-weight per-edge; learned per-node (Node-DP)
- Performs best when trained with minimal real-world data

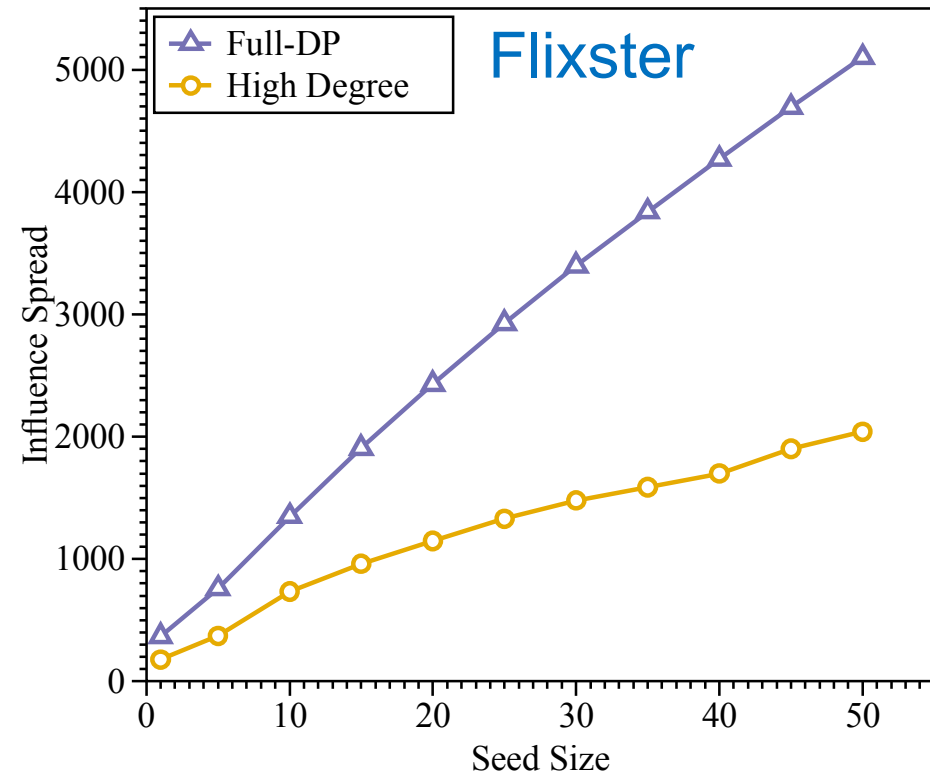
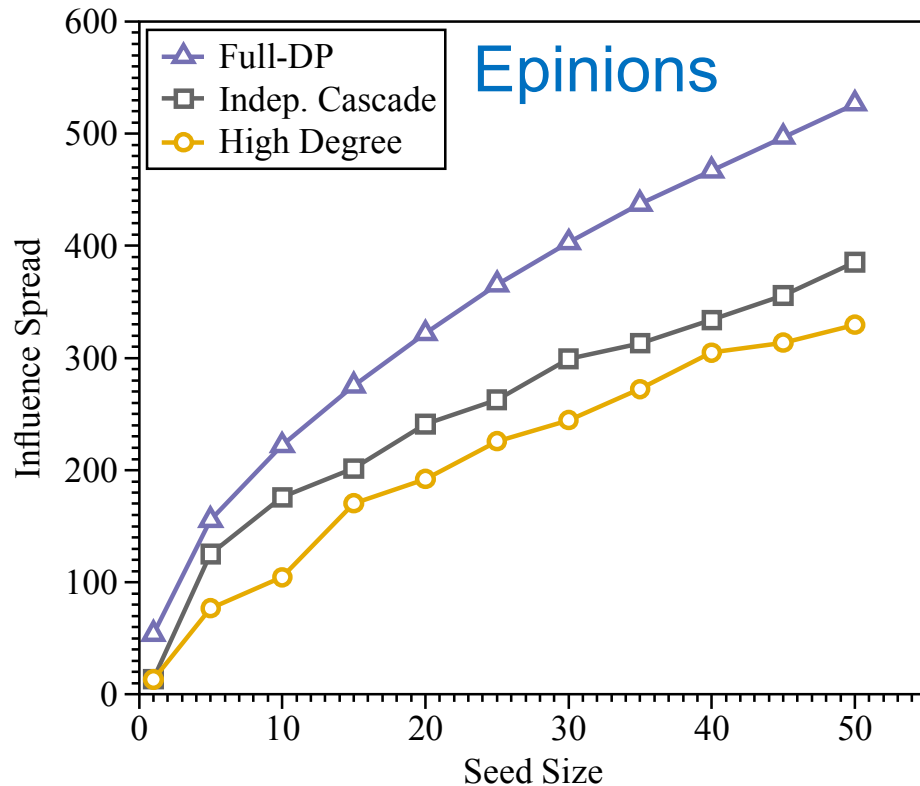


# Application #2 – Maximal Influencers

- Full  $k$ -seed influence maximization is NP-Hard
  - Greedy algorithm widely used (Kempe, 2003)
  - CELF gives same set; more efficient (Leskovec, 2007)
- Diffusion algorithms tested
  - IC (Epinions only)
  - Full-DP
  - Edge-DP
  - High Degree\*
- Procedure
  - For given diffusion method, compute 50 most influential nodes using CELF
  - For common comparison, compute diffusion spread for those seeds using Full-DP

\*High Degree selects nodes based on degree – does not require CELF runs

# Influence Maximization (1 of 2)



Note: Edge-DP not shown as results are nearly identical to Full-DP

# Influence Maximization (2 of 2)

- Overlap between identified seed sets

Epinions	10	20	30	40	50
High Degree	3	6	10	16	18
Edge-DP	10	18	29	36	45
Indep. Cascade	6	9	14	18	24

Flixster	10	20	30	40	50
High Degree	0	1	3	4	5
Edge-DP	10	20	30	40	49

- How do the selected seeds differ?
  - Community Detection*: Full-DP and IC chose seeds in separate communities more than Max Degree
  - Average Distance*: Full-DP chose seeds further apart than IC which chose further apart than Max Degree
  - Node Degree*: Full-DP chose lower degree nodes than IC which chose lower degree nodes than Max Degree
    - Full-DP and IC chose nodes well above average degree
- Balance between higher degree and distance between seeds

# Compute Resources

- Full DP (Flixster maximization)
  - Initial computation for each node as seed: 12 hours on 60 compute nodes
  - CELF identification of 50 top nodes: 16 minutes on workstation
  - Average propagation: 4 seconds
    - Contrast IC with 10,000 MC simulations: 6 minutes

# Summary and Conclusions

- Objective: Predict information spread through large networks
- Monte Carlo methods
  - Independent cascade
- Deterministic methods
  - Belief propagation (BP)
  - Extension: Directed belief propagation (DP)
- Applications
  - Influence spread
  - Influence maximization
- BP and DP
  - More accurate to real-world datasets
  - Computationally efficient enough for large-scale datasets