

Profiling Kokkos Applications

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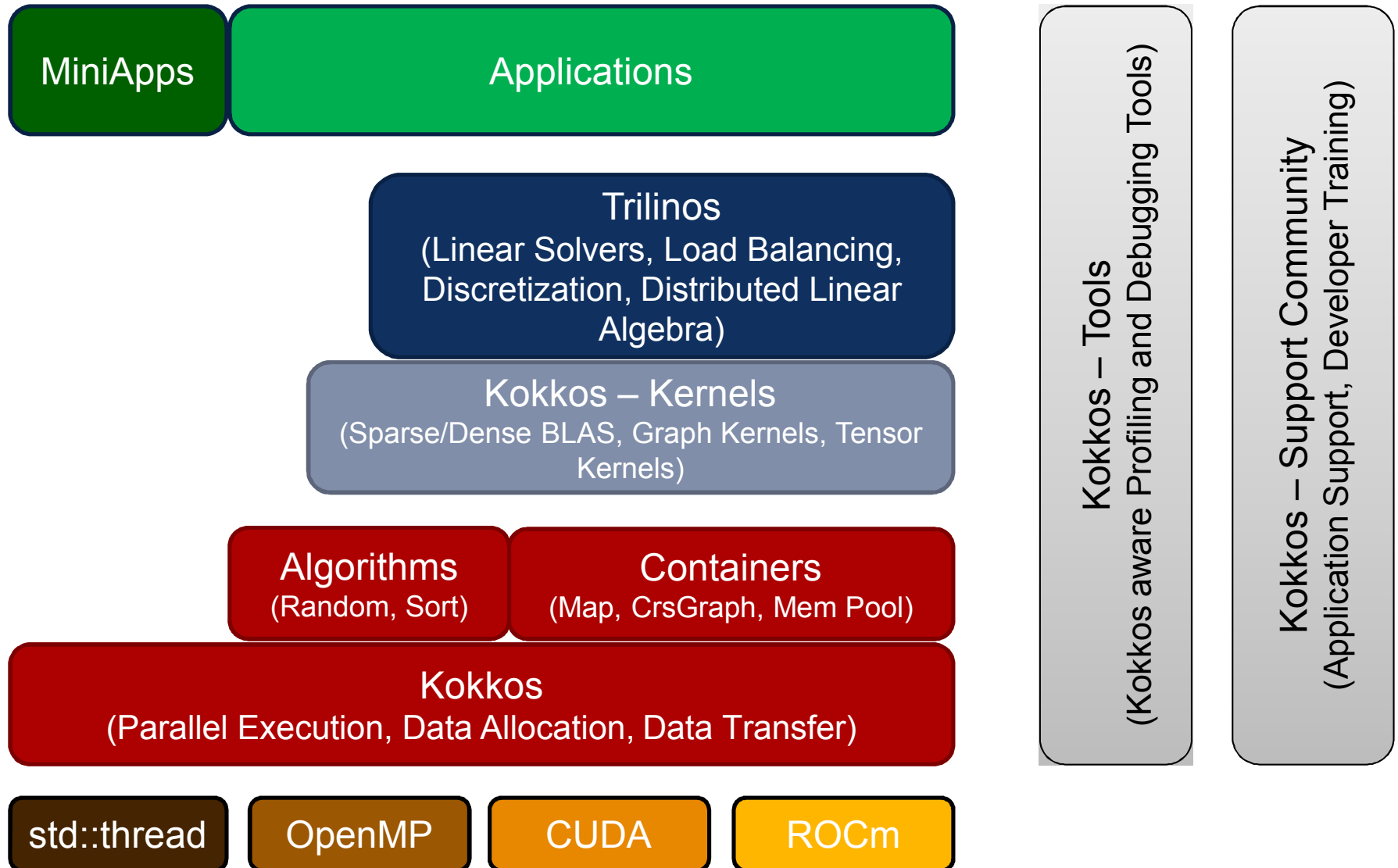
C++ The Bane of Profiling Tools

- It is hard to understand C++ code for a compiler
 - Template Metaprogramming
 - Function Pointers
 - Inheritance
 - Arbitrary aliasing
- It is even harder for a Performance Analysis Tool
 - Most come from a Fortran history
- Abstraction Models make it all worse
 - E.g. only one place where the actual OpenMP loop is
 - Really complex type names which may exceed internal typename length limits
- And we want this across all Platforms ...

Abstractions to the Win ?!

- Abstractions can also help us: Instrumentation
- KokkosTools provide built-in instrumentation for Kokkos applications
 - By default enabled on most platforms
- This Instrumentation knows about Kokkos abstractions
 - Get information organized by Kokkos constructs (Parallel Regions, Allocations in Memory Spaces, etc.)
- Enables Meta Instrumentation for Third Party Tools
 - Provide information to Vtune, Nsight, ...
- Easy to use Tools Provide Basic Information accross all Platforms
 - Kernel and Region Times, Memory Utilization, Allocation and Deallocation Frequency, ...

Building an EcoSystem



- Utilities
 - **KernelFilter:** Enable/Disable Profiling for a selection of Kernels
- Kernel Inspection
 - **KernelLogger:** Runtime information about entering/leaving Kernels and Regions
 - **KernelTimer:** Postprocessing information about Kernel and Region Times
- Memory Analysis
 - **MemoryHighWater:** Maximum Memory Footprint over whole run
 - **MemoryUsage:** Per Memory Space Utilization Timeline
 - **MemoryEvents:** Per Memory Space Allocation and Deallocations
- Third Party Connector
 - **VTune Connector:** Mark Kernels as Frames inside of Vtune
 - **VTune Focused Connector:** Mark Kernels as Frames + start/stop profiling

How to Use the Tools

- Checkout from github.com/kokkos/kokkos-tools
- Documentation in a Wiki
- Go to `src/tools/TOOLNAME` and build the tool
 - On most platforms typing "make" is enough
- Before running your code set environment variable
 - `export KOKKOS_PROFILE_LIBRARY=/PATH/TO/TOOLS/LIBRARY`
- Analyze output
 - Some tools print to screen, some write per-process files
 - Some tools have readers for binary output files
- How does it work internally?
 - Instrumentation is always active, but internal function pointers are NULL
 - At `Kokkos::initialize` tool library is dynamically loaded, and function pointers are set

Typical Approach

- Run KernelTimer
 - Check if majority of time is in Kernels
 - Check for HotSpot Kernels
- Run MemoryUsage
 - Check where your memory utilization is coming from
 - Check total number of entries to see if frequent alloc/dealloc could be an issue
- If frequent allocations are an issue: run MemoryEvents
 - Figure out which allocations are causing the issue
 - Less than 1000 per second per socket is usually no issue
- To find unaccounted time: put region markers into code
 - Compare region times with kernel times
- Use Connector tools to help investigate individual kernels

The Tools

- See Wiki

Exercise

- Use new miniApp: ExaMiniMD
 - Note this is NOT a full OpenSource release
 - This can be shared within ECP or under vendor NDA, but ask first
 - Do not share with anyone else; OpenSource license expected this fall
 - git clone [git@github.com:crtrott/ExaMiniMD](https://github.com:crtrott/ExaMiniMD)
 - git checkout profiling-exercise
- Basic Molecular Dynamics Importance of Kernels:
 - Force Calculation
 - NeighborList Construction
 - Communication
 - Other stuff: (Integration, Particle Sorting etc.)
- This variant has a problem hidden
 - Use Kokkos Tools to find the issue
 - Follow the typical approach lined out before

(NVIDIA) GPU Profiling

- Dominant Performance Bottlenecks:
 - Occupancy
 - Memory Bandwidth
 - Memory Efficiency
 - Memory Load/Store Slots
 - Availability of Instruction Parallelism
- Visual Profiler (nvvp or as part of nsight)
 - Guided Analysis
- Use nvprof to collect data on commandline
 - Generally same information as in the Visual Profiler
 - `nvprof [OPTIONS] ./Executable [OPTIONS]`
 - `--print-gpu-summary` : Summary of Kernels, and data transfers
 - `--print-gpu-trace`: timeline of kernels and data transfers
 - `--query-metrics`: list of collectable events
 - `-m [EVENTS]`: set events to be collected
 - `--kernels [KERNELS]`: restrict profiling to specified kernels

Important Metrics – Occupancy/Mem

- **achieved_occupancy:**
 - Actually reached occupancy
 - Cause1: high register pressure (check with --print-gpu-trace)
 - Cause2: high shared memory usage (check with --print-gpu-trace)
 - Cause3: low total available parallelism (too few blocks)
- ****_throughput:** Bandwidth for different parts of the memory subsystem
 - **dram_[read/write]:** device memory traffic including ECC
 - **[gld/gst]:** global memory access, could be cached (this is larger than requested due to efficiency)
 - **[gld/gst]_requested:** the memory throughput of things the code actually wants
 - **l2_l1_read, l2_tex_[read/write], l2_atomic:** L2 Cache Throughput by Source
 - **local_[load/store]:** data traffic due to register spilling
 - **shared_[load/store]:** shared memory (team scratch level 0)
- ****_efficiency:** different efficiency metrics
 - **[gld/gst]:** global memory access (coalesced access = 100%)
 - **shared:** shared memory loads
- ****_hit_rate:** Cache hit rates
 - **l1_cache_[global/local]:** Hit rate in L1 Cache due to global/local load store
 - **tex_cache:** Hit rate for texture fetches
 - **l2_l1_[read]:** Hit rate for all L1 misses in L2
 - **l2_tex_read:** Hit rate for texture misses in L2

Important Metrics – Compute

- ****_efficiency:**
 - **sm:** multiprocessors are active (small means not enough work launched in a kernel)
 - **warp_execution:** active threads vs non-active threads due to branching (small means divergence)
 - **branch:** non-divergent vs total branches (small means divergence in kernel)
 - **flop_[sp/dp]:** achieved single/double precision peak flop/s fraction
- **stall_**:** Reasons why a warp does not execute an instruction
 - **inst_fetch:** instructions are not yet loaded,
 - very unlikely to be an issue, if it is think about breaking up kernel into smaller ones
 - **memory_dependency:** waiting for a load
 - typical sign for memory bandwidth limitation
 - Use less memory, spread out loads more if possible to overlap with compute
 - **exec_dependency:** can't execute because prior instruction not done
 - find and expose instruction parallelism
 - **memory_throttle:** no load/store slots available
 - If this happen without memory_dependency data access patterns are usually bad
 - **sync:** warps waiting for other warps at a barrier
 - Check if barriers are necessary
 - **pipe_busy:** compute operation pipe is busy
 - Rarely an issue, except when making haevy use of special function units
 - **not_selected:** The “good” stall, more ready threads are availabe than slots are to be filled
 - If you see this as the primary stall reason you have a code which is either artificial or a Gordon Bell candidate

A Case Study I

- A Parametrized Benchmark Code

- Very fancy vector addition with additional math

- Parameters:

- N: control total work (RT)
 - M: control data reuse origin (RT)
 - S: control data stride (RT)
 - R: control data reuse (RT)
 - K: control instruction parallelism (CT)
 - F: control Flops/Bytes ratio (RT)

```
View<double***> a("A",N,M,S), b("B",N,M,S),
                c("C",N,M,S);
// Loop over blocks
for(int i=0; i<N; i++) {
    // Repeat work on block to control cache reuse
    for(int r=0; r<R; r++) {
        // Loop over block
        for(int j=0; j<M; j++) {
            a_1 = a(i,j,0);
            b_1 = b(i,j,0);
            // Repeat for instruction parallelism
            a_K = a(i,j,0);
            b_K = a(i,j,0);
            // Loop to add more flops
            for(int f=0; f<F; f++) {
                a_1 += b_1;
                // Repeat for instruction parallelism
                a_K += b_K;
            }
            c(i,j,0) = a_1 + /*...*/ a_K;
        }
    }
}
```

A Case Study II

- Exercise I
 - Find settings to measure hardware global and cache bandwidth
- Exercise II
 - Find setting to maximize flop rate
- Exercise III
 - Find settings to make each of the stall reasons the primary stall reasons



1. *Journal of the American Medical Association*, 2000; 283: 2686-2692.

<http://www.github.com/kokkos>