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Aztec User's Guide

Version 1.0

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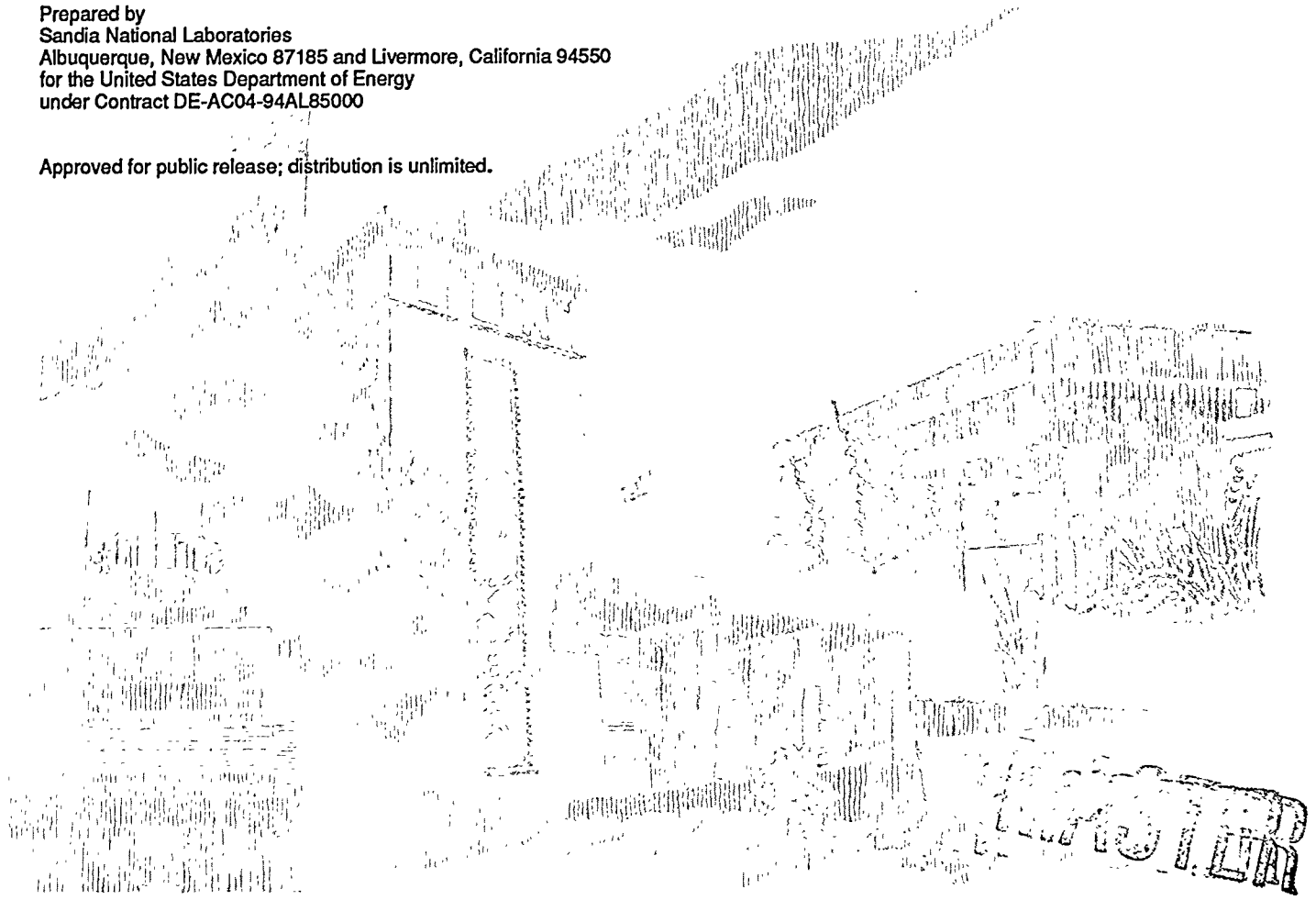

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Aztec User's Guide*
Version 1.0

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Abstract

Aztec is an iterative library that greatly simplifies the parallelization process when solving the linear systems of equations $Ax = b$ where A is a user supplied $n \times n$ sparse matrix, b is a user supplied vector of length n and x is a vector of length n to be computed. **Aztec** is intended as a software tool for users who want to avoid cumbersome parallel programming details but who have large sparse linear systems which require an efficiently utilized parallel processing system. A collection of data transformation tools are provided that allow for easy creation of distributed sparse unstructured matrices for parallel solution. Once the distributed matrix is created, computation can be performed on any of the parallel machines running **Aztec**: nCUBE 2, IBM SP2 and Intel Paragon, MPI platforms as well as standard serial and vector platforms.

Aztec includes a number of Krylov iterative methods such as conjugate gradient (CG), generalized minimum residual (GMRES) and stabilized biconjugate gradient (BiCGSTAB) to solve systems of equations. These Krylov methods are used in conjunction with various preconditioners such as polynomial or domain decomposition methods using LU or incomplete LU factorizations within subdomains. Although the matrix A can be general, the package has been designed for matrices arising from the approximation of partial differential equations (PDEs). In particular, the **Aztec** package is oriented toward systems arising from PDE applications.

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Contents.

Notation Conventions

Different fonts are used to indicate program fragments, keys words, variables, or parameters in order to clarify the presentation. The table below describes the meaning denoted by these different fonts.

Convention	Meaning
`typewriter`	File names, code examples and code fragments.
sans serif	C language elements such as function names and constants when they appear embedded in text or in function definition syntax lines.
italics	Parameter and variable names when they appear embedded in text or function definition syntax lines.
AZ_	C language elements such as function names and constants which are supplied by the **Aztec** library.

Code Distribution

Aztec is publicly available for research purposes and may be licensed for commercial application. The code is distributed along with technical documentation, example C and Fortran driver routines and sample input files via the internet. It may be obtained by contacting one of the authors listed on page i of this report.

1. Overview. **Aztec** is an iterative library that greatly simplifies the parallelization process when solving the linear system of equations

$$Ax = b$$

where A is a user supplied $n \times n$ sparse matrix, b is a user supplied vector of length n and x is a vector of length n to be computed. **Aztec** is intended as a software tool for users who want to avoid cumbersome parallel programming details but who have large sparse linear systems requiring efficient use of a parallel processing system. The most complicated parallelization task for an **Aztec** user is the distributed matrix specification for the particular application. Although this may seem difficult, a collection of data transformation tools are provided that allow creation of distributed sparse unstructured matrices for parallel solution with ease of effort that is similar to a serial implementation. Background information regarding the data transformation tools can be found in [5]. Once the distributed matrix is created, computation can occur on any of the parallel machines running **Aztec**: nCUBE 2, IBM SP2, Intel Paragon, and MPI platforms. In addition, **Aztec** can be used on standard serial and vector platforms such as SUN, SGI and CRAY computers.

Aztec includes a number of Krylov iterative methods such as conjugate gradient (CG), generalized minimum residual (GMRES) and stabilized biconjugate gradient (BiCGSTAB) to solve systems of equations. These Krylov methods are used in conjunction with various preconditioners such as polynomial preconditioners or domain decomposition using LU or incomplete LU factorizations within subdomains. Background information concerning the iterative methods and the preconditioners can be found in [4]. Although the matrix A can be general, the package has been designed for matrices arising from the approximation of partial differential equations (PDEs). In particular, the preconditioners, iterative methods and parallelization techniques are oriented toward systems arising from PDE applications. Lastly, **Aztec** can use one of two different sparse matrix notations – either a point-entry modified sparse row (MSR) format or a block-entry variable block row (VBR) format. These two formats have been generalized for parallel implementation and, as such, are referred to as "distributed" yielding DMSR and DVBR references.

The remainder of this guide describes how **Aztec** is invoked within an application. **Aztec** is written in ANSI-standard c and as such, all arrays in the descriptions which follow begin indexing with 0. Also, all function prototypes (loosely, descriptions) are presented in ANSI c format. Section 2 discusses iterative method, preconditioning and convergence options. Section 3 explains vectors and sparse matrix formats supported by **Aztec**. In Section 4 we discuss the data transformation tool for creating distributed vectors and matrices. A concrete detailed programming example using this tool is given in Section 5 and some advance topics are discussed in Section 6. Finally, Section 7 gives a glossary of **Aztec** functions available to users.

2. Aztec: High Level View. The following tasks must be performed to successfully invoke **Aztec**:

- describe the parallel machine (e.g. number of processors).
- initialize matrix and vector data structures.
- choose iterative methods, preconditioners and the convergence criteria.
- initialize the right hand side and initial guess.
- invoke the solver.

Example

```
#include "az_aztec.h"

void main(void) {

  AZ_processor_info(proc_config);

  init_matrix_vector_structures(bindx, val, update, external,
                                update_index, extern_index, data_org);
  init_options(options, params);

  init_guess_and_rhs(x, b, data_org, update, update_index);

  AZ_solve(x, b, options, params, bindx, val, data_org, status,
           proc_config);
}
```

FIG. 1. *High level code for* **Aztec** *application.*

A sample C program is shown in Figure 1 omitting declarations and some parameters[1]. The functions init_matrix_vector_structures, init_options, and init_guess_and_rhs are supplied by the user. In this section, we give an overview of **Aztec**'s features by describing the user input arrays, *options* and *params*, that are set by the user in the function init_options. A discussion of the other subroutines is deferred to Sections 4 and 5.

2.1. Aztec Options. *options* is an integer array of length AZ_OPTIONS_SIZE set by the user. It is used (but not altered) by the function AZ_solve to choose between iterative solvers, preconditioners, etc. Below we discuss each of the possible options. In some of these descriptions, reference is made to a user-defined *options* or *params* value which is yet be introduced. These descriptions will follow but the reader may wish to "jump ahead" and read the descriptions if the immediate context is not clear.

Specifications

options[AZ_solver]	Specifies solution algorithm. DEFAULT: AZ_gmres.
AZ_cg	Conjugate gradient (only applicable to symmetric positive definite matrices).
AZ_gmres	Restarted generalized minimal residual.
AZ_cgs	Conjugate gradient squared.
AZ_tfqmr	Transpose-free quasi-minimal residual.

[1] The entire main program with specific sample problems is distributed with the package in the file az_main.c

AZ_bicgstab	Bi-conjugate gradient with stabilization.
AZ_lu	Sparse direct solver (single processor only).
*options[*AZ_scaling*]*	Specifies scaling algorithm. The entire matrix is scaled (overwriting the old matrix). Additionally, the right hand side, the initial guess and the final computed solution are scaled if necessary. DEFAULT: AZ_none.
AZ_none	No scaling.
AZ_Jacobi	Point Jacobi scaling.
AZ_BJacobi	Block Jacobi scaling where the block size corresponds to the VBR blocks. Point Jacobi scaling is performed when using the MSR format.
AZ_row_sum	Scale each row so the magnitude of its elements sum to 1.
AZ_sym_diag	Symmetric scaling so diagonal elements are 1.
AZ_sym_row_sum	Symmetric scaling using the matrix row sums.
*options[*AZ_precond*]*	Specifies preconditioner. DEFAULT: AZ_none.
AZ_none	No preconditioning.
AZ_Jacobi	k step Jacobi (block Jacobi for DVBR matrices where each block corresponds to a VBR block). The number of Jacobi steps, k, is set via *options[*AZ_poly_ord*]*.
AZ_Neumann	Neumann series polynomial where the polynomial order is set via *options[*AZ_poly_ord*]*.
AZ_ls	Least-squares polynomial where the polynomial order is set via *options[*AZ_poly_ord*]*.
AZ_lu	Domain decomposition preconditioner (additive Schwarz) using a sparse LU factorization in conjunction with a drop tolerance *params[*AZ_drop*]* on each processor's submatrix. The treatment of external variables in the submatrix is determined by *options[*AZ_overlap*]*. The current sparse lu factorization is provided by the package y12m [6].
AZ_ilu	Similar to AZ_lu using ilu(0) instead of LU.
AZ_bilu	Similar to AZ_lu using block ilu(0) instead of LU where each block corresponds to a VBR block.

AZ_sym_GS

Non-overlapping domain decomposition (additive Schwarz) k step symmetric Gauss-Siedel. In particular, a symmetric Gauss-Siedel domain decomposition procedure is used where each processor independently performs one step of symmetric Gauss-Siedel on its local matrix, followed by communication to update boundary values before the next local symmetric Gauss-Siedel step.The number of steps, k, is set via *options[*AZ_poly_ord*]*.

*options[*AZ_conv*]*

Determines the residual expression used in convergence checks and printing. DEFAULT: AZ_r0. The iterative solver terminates if the corresponding residual expression is less than *params[*AZ_tol*]*:

AZ_r0

$\|r\|_2/\|r^{(0)}\|_2$

AZ_rhs

$\|r\|_2/\|b\|_2$

AZ_Anorm

$\|r\|_2/\|A\|_\infty$

AZ_sol

$\|r\|_\infty/(\|A\|_\infty * \|x\|_1 + \|b\|_\infty)$

AZ_weighted

$\|r\|_{WRMS}$
where $\|\cdot\|_{WRMS} = \sqrt{(1/n)\sum_{i=1}^{n}(r_i/w_i)^2}$, n is the total number of unknowns, w is a weight vector provided by the user via *params[*AZ_weights*]* and $r^{(0)}$ is the initial residual.

*options[*AZ_output*]*

Specifies information (residual expressions - see *options[*AZ_conv*]*) to be printed. DEFAULT: 1.

AZ_all

Print out the matrix and indexing vectors for each processor. Print out all intermediate residual expressions.

AZ_none

No intermediate results are printed.

AZ_last

Print out only the final residual expression.

> 0

Print residual expression every *options[*AZ_output*]* iterations.

*options[*AZ_pre_calc*]*

Indicates whether to use factorization information from previous calls to AZ_solve. DEFAULT: AZ_calc.

AZ_calc

Use no information from previous AZ_solve calls.

AZ_recalc

Use preprocessing information from a previous call but recalculate preconditioning factors. This is primarily intended for factorization software which performs a symbolic stage.

AZ_reuse — Use preconditioner from a previous AZ_solve call, do not recalculate preconditioning factors. Also, use scaling factors from previous call to scale the right hand side, initial guess and the final solution.

*options[*AZ_max_iter*]* — Maximum number of iterations. DEFAULT: 500.

*options[*AZ_poly_ord*]* — The polynomial order when using polynomial preconditioning. Also, the number of steps when using Jacobi or symmetric Gauss-Seidel preconditioning. DEFAULT: 3.

*options[*AZ_overlap*]* — Determines the submatrices factored with the domain decomposition algorithms: AZ_lu, AZ_ilu, AZ_bilu. DEFAULT: AZ_none.

- AZ_none — Factor the local submatrix defined on this processor discarding column entries that correspond to external elements.
- AZ_diag — Factor the local submatrix defined on this processor augmented by a diagonal (block diagonal for VBR format) matrix. This diagonal matrix corresponds to the diagonal entries of the matrix rows (found on other processors) associated with external elements. This can be viewed as taking one Jacobi step to update the external elements and then performing domain decomposition with AZ_none on the residual equations.
- AZ_full — Factor the local submatrix defined on this processor augmented by the rows (found on other processors) associated with external variables (discarding column entries associated with variables not defined on this processor). The resulting procedure is an overlapped additive Schwarz procedure.

*options[*AZ_kspace*]* — Krylov subspace size for restarted GMRES. DEFAULT: 30.

*options[*AZ_orthog*]* — GMRES orthogonalization scheme. DEFAULT: AZ_classic.

- AZ_classic — Classical Gramm-Schmidt orthogonalization.
- AZ_modified — Modified Gramm-Schmidt orthogonalization.

*options[*AZ_aux_vec*]* — Determines $\tilde{r}$ (a required vector within some iterative methods). The convergence behavior varies slightly depending on how this is set. DEFAULT: AZ_resid.

- AZ_resid — $\tilde{r}$ is set to the initial residual vector.

AZ_rand	$\tilde{r}$ is set to random numbers between -1 and 1. NOTE: When using this option, the convergence depends on the number of processors (i.e. the iterates obtained with x processors differ from the iterates obtained with y processors if x $\neq$ y).

2.2. Aztec parameters. *params* is a double precision array set by the user and normally of length AZ_PARAMS_SIZE. However, when a weight vector is needed for the convergence check (i.e. *options*[AZ_conv] = AZ_weighted), it is embedded in *params* whose length must now be AZ_PARAMS_SIZE + # of elements updated on this processor. In either case, the contents of *params* are used (but not altered) by the function AZ_solve to control the behavior of the iterative methods. The array elements are specified as follows:

Specifications

params[AZ_tol]	Specifies tolerance value used in conjunction with convergence tests. DEFAULT: 10^{-6}.
params[AZ_drop]	Specifies drop tolerance used in conjunction with LU preconditioner. DEFAULT: 0.0.
params[AZ_weights]	When *options*[AZ_conv] = AZ_weighted, the i'th local component of the weight vector is stored in the location *params*[AZ_weights+i].

Figure 2 illustrates a sample function init_options where the Aztec function AZ_defaults sets the default options.

Example

```
void init_options(int options[AZ_OPTIONS_SIZE],
                  double params[AZ_PARAMS_SIZE])
{
  AZ_defaults(options,params);
  options[AZ_solver]     = AZ_cgs;
  options[AZ_scaling]    = AZ_none;
  options[AZ_precond]    = AZ_ls;
  options[AZ_output]     = 1;
  options[AZ_max_iter]   = 640;
  options[AZ_poly_ord]   = 7;
  params[AZ_tol]         = 0.0000001;
}
```

FIG. 2. *Example option initialization routine (*init_options*).*

2.3. Return status. *status* is a double precision array of length AZ_STATUS_SIZE returned from AZ_solve[2]. The contents of *status* are described below.

Specifications

status[AZ_its]	Number of iterations taken by the iterative method.
status[AZ_why]	Reason why AZ_solve terminated.
AZ_normal	User requested convergence criteria is satisfied.
AZ_param	User requested option is not available.
AZ_breakdown	Numerical breakdown occurred.
AZ_loss	Numerical loss of precision occurred.
AZ_maxits	Maximum iterations taken without convergence.
status[AZ_r]	The true residual norm corresponding to the choice *options*[AZ_conv] (this norm is calculated using the computed solution).
status[AZ_scaled_r]	The true residual ratio expression as defined by *options*[AZ_conv].
status[AZ_rec_r]	Norm corresponding to *options*[AZ_conv] of final residual or estimated final residual (recursively computed by iterative method). Note: When using the 2-norm, **tfqmr** computes an estimate of the residual norm instead of computing the residual.

When AZ_solve returns abnormally, the user may elect to restart using the current computed solution as an initial guess.

3. Data Formats. In this section we describe the matrix and vector formats used internally by **Aztec**. In Section 4 we discuss a tool that transforms data from a simpler format to this format. Here, the terms "element" and "component" are used interchangeably to denote a particular entry of a vector.

The sparse matrix-vector product, $y \leftarrow Ax$, is the major kernel operation of **Aztec**. To perform this operation in parallel, the vectors x and y as well as the matrix A must be distributed across the processors. The elements of any vector of length n are assigned to a particular processor via some partitioning method (e.g. **Chaco** [2]). When calculating elements in a vector such as y, a processor computes only those elements in y which it has been assigned. These vector elements are explicitly stored on the processor and are defined by a set of indices referred to as the processor's *update* set. The *update* set is further divided into two subsets: *internal* and *border*. A component corresponding to an index in the *internal* set is updated using only information on the

[2] All integer information returned from AZ_solve is cast into double precision and stored in *status*.

current processor. As an example, the index i is in *internal* if, in the matrix-vector product kernel, the element y_i is updated by this processor and if each j defining a nonzero A_{ij} in row i is in *update*. The *border* set defines elements which would require values from other processors in order to be updated during the matrix vector product. For example, the index i is in *border* if, in the matrix-vector product kernel, the element y_i is updated by this processor and if there exists at least one j associated with a nonzero A_{ij} found in row i that is not in *update*. In the matrix-vector product, the set of indices which identify the off-processor elements in x that are needed to update components corresponding to *border* indices is referred to as *external*. They are explicitly stored by and are obtained from other processors via communication whenever a matrix-vector product is performed. Figure 3 illustrates how a set of vertices in a partitioning of a grid would be used to define these sets. Since these sets of indices are used exclusively

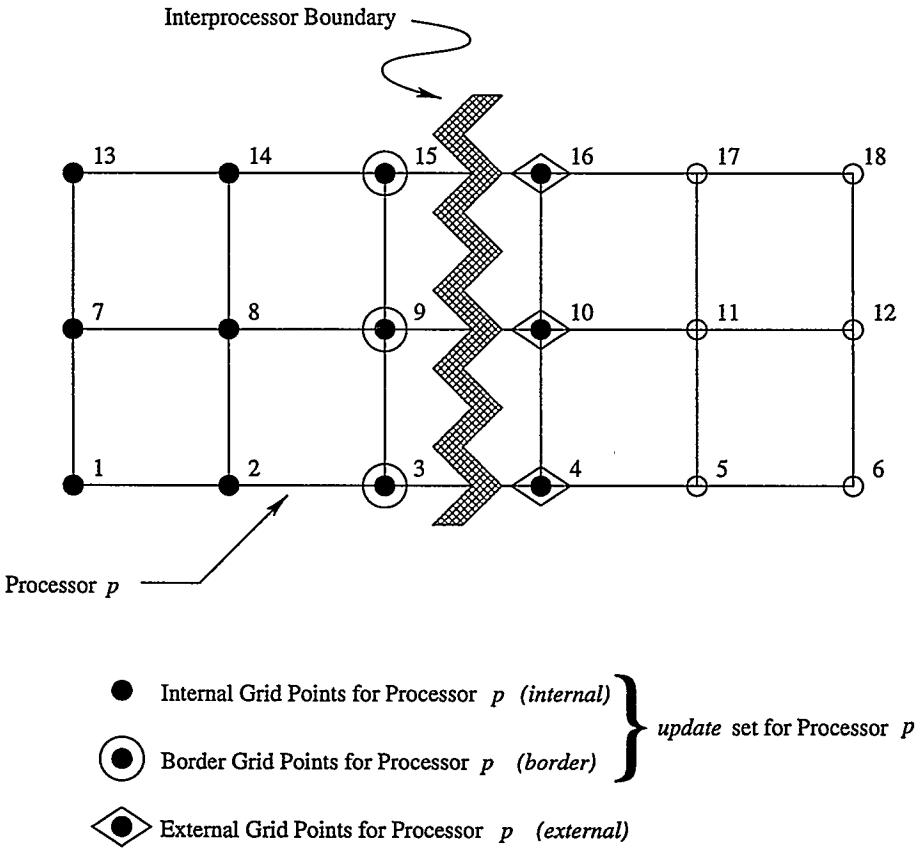

FIG. 3. *Example partitioning of a finite element grid.*

to reference specific vector components, the same names (i.e., *update*, *internal*, *border* and *external*) are sometimes used below to describe the vector elements themselves. Having generalized these labels, the three types of vector elements are distinguished by locally storing the *internal* components first, followed by the *border* components and finally by the *external* components. In addition, all *external* components received from the same processor are stored consecutively. Below we summarize the nomenclature for a processor with N total elements where *N_internal*, *N_border*, and *N_external* elements are distributed over the sets *internal*, *border* and *external* respectively.

set	description	local numbering
internal	updated w/o communication	0 to $N\_internal - 1$.
border	updated with communication	$N\_internal$ to $N\_internal+$ $N\_border - 1$.
external	not updated but used to update *border*	$N\_internal + N\_border$ to $N - 1$. elements received from the same processor are numbered consecutively.

Similar to vectors, a subset of matrix non-zeros is stored on each processor. In particular, each processor stores only those rows which correspond to its *update* set. For example, if vector element i is updated on processor p, then processor p also stores all the non-zeros of row i in the matrix. Further, the local numbering of vector elements on a specific processor induces a local numbering of matrix rows and columns. For example, if vector element k is locally numbered as k_l, then all references to row k or column k in the matrix would be locally numbered as k_l. Thus, each processor contains a submatrix whose row and column entries correspond to variables defined on this processor.

The remainder of this section describes the two sparse matrix formats that are used to store the local renumbered submatrix. These two sparse matrix formats correspond to common formats used in serial computations.

3.1. Distributed Modified Sparse Row (DMSR) Format. The DMSR format is a generalization of the MSR format [3]. The data structure consists of an integer vector *bindx* and a double precision vector *val* each of length *N_nonzeros + 1* where *N_nonzeros* is the number of nonzeros in the local submatrix. For a submatrix with m rows the DMSR arrays are as follows:

bindx :

bindx[0]	= m + 1
bindx[k+1] - bindx[k]	= number of nonzero off-diagonal elements in k'th row, $k < m$
bindx[k_s...k_e]	= column indices of the off-diagonal nonzeros in row k, where k_s = *bindx[k]* and k_e = *bindx[k+1]-1.*

val :

val[k]	= $A_{kk}, k < m$
val[k_i]	= the (k, *bindx[k_i]*)'th matrix element where $k_s \leq k_i \leq k_e$ with k_s and k_e as defined above.

Note: *val[m]* is not used. See [1] for a detailed discussion of the MSR format.

3.2. Distributed Variable Block Row (DVBR) Format. The Distributed Variable Block Row (DVBR) format is a generalization of the VBR format [1]. The data structure consists of a double precision vector *val* and five integer vectors: *indx*, *bindx*, *rpntr*, *cpntr* and *bpntr*. The format is best suited for sparse block matrices of

the form

$$A = \begin{pmatrix} A_{00} & A_{01} & \cdots & A_{0k} \\ A_{10} & A_{11} & \cdots & A_{1k} \\ \vdots & & \ddots & \vdots \\ A_{m0} & \cdots & \cdots & A_{mk} \end{pmatrix}$$

where A_{ij} denotes a block (or submatrix). In a sparse block matrix, some of these blocks would be entirely zero while others may be dense. The DVBR vectors are described below for a matrix with $M \times K$ blocks.

rpntr[0 ... M] :

rpntr[0] = 0
rpntr[k+1] - rpntr[k] = number of rows in k'th block row

cpntr[0 ... K] :

cpntr[0] = 0
cpntr[k+1] - cpntr[k] = number of columns in k'th block column

bpntr[0 ... M] :

bpntr[0] = 0
bpntr[k+1] - bpntr[k] = number of nonzero blocks in the k'th block row

bindx[0 ... bpntr[M]] :

bindx[k_s...k_e] = block column indices of nonzero blocks in block row k where k_s = *bpntr[k]* and k_e = *bpntr[k+1]*-1

indx[0 ... bpntr[M]] :

indx[0] = 0
indx[k_i+1] - indx[k_i] = number of nonzeros in the (k, *bindx[k_i]*)'th block where $k_s \leq k_i \leq k_e$ with k_s and k_e as defined above.

val[0 ... indx[bpntr[M]]] :

val[i_s...i_e] = nonzeros in the (k, *bindx[k_i]*)'th block stored in column major order where k_i is as defined above, i_s = *indx[k_i]* and i_e = *indx[k_i+1]*-1

See [1] for a detailed discussion of the VBR format.

4. High Level Data Interface. Setting up the distributed format described in Section 3 for the local submatrix on each processor can be quite cumbersome. In particular, the user must determine a mapping between the global numbering scheme and a local scheme which facilitates proper communication. Further, a number of additional variables must be set for communication and synchronization (see Section 6). In this section we describe a simpler data format that is used in conjunction with a transformation function to generate data structures suitable for **Aztec**. The new format allows the user to specify the rows in a natural order as well as to use global column numbers in the *bindx* array. To use the transformation function the user supplies the

update set and the submatrix for each processor. Unlike the previous section, however, the submatrix is specified using the global coordinate numbering instead of the local numbering required by **Aztec**. This procedure greatly facilitates matrix specification and is the main advantage of the transformation software.

On a given processor, the *update* set (i.e. vector element assignment to processors) is defined by initializing the array *update* on each processor so that it contains the global index of each element assigned to the processor. The *update* array must be sorted in ascending order (i.e. $i < j \Rightarrow update[i] < update[j]$). This sorting can be performed using the **Aztec** function AZ_sort. Matrix specification occurs using the arrays defined in the previous section. However, now the local rows are defined in the same order as the *update* array and column indices (e.g. *bindx*) are given as global column indices. To illustrate this in more detail, consider the following example matrix:

$$A = \begin{pmatrix} a_{00} & a_{01} & & a_{03} & a_{04} & \\ a_{10} & a_{11} & & a_{13} & & \\ & & a_{22} & a_{23} & a_{24} & a_{25} \\ a_{30} & a_{31} & a_{32} & a_{33} & a_{34} & a_{35} \\ a_{40} & & a_{42} & a_{43} & a_{44} & \\ & & a_{52} & a_{53} & & a_{55} \end{pmatrix}.$$

Figure 4 illustrates the information corresponding to a particular matrix partitioning that is specified by the user as input to the data transformation tool. Using this

Examples

proc 0:

N_update:	3													
update:	0	1	3											
bindx:	4	7	9	14	1	3	4	0	3	1	0	4	2	5
val:	a_{00}	a_{11}	a_{33}	-	a_{01}	a_{03}	a_{04}	a_{10}	a_{13}	a_{30}	a_{31}	a_{32}	a_{34}	a_{35}

proc 1:

N_update:	1				
update:	4				
bindx:	2	5	0	3	2
val:	a_{44}	-	a_{40}	a_{42}	a_{43}

proc 2:

N_update:	2							
update:	2	5						
bindx:	3	6	8	4	3	5	3	2
val:	a_{22}	a_{55}	-	a_{23}	a_{24}	a_{25}	a_{52}	a_{53}

FIG. 4. *User input (MSR format) to initialize the sample matrix problem.*

information, AZ_transform

- determines the sets *internal*, *border* and *external*.
- determines the local numbering: *update_index[i]* is the local numbering for *update[i]* while *extern_index[i]* is the local numbering for *external[i]*.

- permutes and renumbers the local submatrix rows and columns so that they now correspond to the new ordering.
- computes additional information (e.g. the number of internal, border and external components on this processor) and stores this in *data_org* (see Section 6).

A sample transformation is given in Figure 5 and is found in the file `az_app_utils.c`.

Example

```
init_matrix_vector_structures(bindx, val, update, external,
                              update_index, extern_index, data_org);
{
  AZ_read_update(update, N_update);
  create_matrix(bindx, val, update, N_update);
  AZ_transform(bindx, val, update, external, update_index,
               extern_index, data_org, N_update);
}
```

FIG. 5. init_matrix_vector_structures.

AZ_read_update is an **Aztec** utility which reads a file and assigns elements to *update*. The user supplied routine create_matrix creates an MSR or VBR matrix using the global numbering. Once transformed the matrix can now be used within **Aztec**.

5. Examples. A sample program is described by completing the program fragments given earlier (Figures 1, 2 and 5). In Figure 1, AZ_processor_info is an **Aztec** utility which initializes the array *proc_config* to reflect the number of processors being used and the node number of this processor. The function AZ_solve is also supplied by **Aztec** to solve the user supplied linear system. Thus, the only functions that the user must supply which have not already been discussed include: init_guess_and_rhs in Figure 1 and create_matrix in Figure 5.

The function init_guess_and_rhs initializes the initial guess and the right hand side.

Example

```
void init_guess_and_rhs(x, rhs, data_org, update, update_index)
{
  N_update = data_org[AZ_N_internal] + data_org[AZ_N_border];
  for (i = 0; i < N_update ; i = i + 1) {
    rhs[update_index[i]] = (double) update[i];
    x[i] = 0.0;
  }
}
```

FIG. 6. init_guess_and_rhs.

In Figure 6, a sample routine is given which sets the initial guess vector to zero and sets the right hand side vector equal to the global indices (where the local element *update_index[i]* corresponds to global element *update[i]*, see Section 4).

A create_matrix function to initialize an MSR matrix is illustrated in Figure 7. Different matrix problems can be implemented by changing the function add_row which computes the MSR entries corresponding to a new row of the matrix. The specific

Example

```
void create_matrix(bindx, val, update, N_update)

{
  N_nonzeros = N_update + 1;
  bindx[0] = N_nonzeros;

  for (i = 0; i < N_update; i = i + 1)
    add_row(update[i], i, val, bindx);

}
```

FIG. 7. create_matrix.

add_row function for implementing a 5-point 2D Poisson operator on an $n \times n$ grid is shown in Figure 8 (n is a global variable set by the user). With these few lines of code

Example

```
void add_row(row, location, val, bindx)
{
   k = bindx[location];

   /* check neighboring points in each direction and add nonzero */
   /* entry if neighbor exists.                                  */

   bindx[k]  = row + 1;   if (row%n !=     n-1) val[k++] = -1.;
   bindx[k]  = row - 1;   if (row%n !=       0) val[k++] = -1.;
   bindx[k]  = row + n;   if ((row/n)%n != n-1) val[k++] = -1.;
   bindx[k]  = row - n;   if ((row/n)%n !=   0) val[k++] = -1.;

   bindx[location+1] = k;
   val[location] = 4.;        /* matrix diagonal */
}
```

FIG. 8. add_row *for a 2D Poisson problem*

and the functions described earlier, the user initializes and solves a 2D Poisson problem. While for simplicity of presentation this specific example is structured the **Aztec** library does not assume any structure in the sparse matrix. All the communication and variable renumbering is done automatically without the assumption of structured communication.

Other add_row functions corresponding to a 3D Poisson equation and a high order 2D Poisson equation are distributed with **Aztec** (file `az_examples.c`). We recommend that potential users review at these examples. In many cases, new applications can be written by simply editing these programs. The interested reader should note that only a few lines of code are different between the functions for the 5-pt Poisson, the high order Poisson and the 3D Poisson codes. Further, the add_row routines are essentially identical to those that would be used to set up sparse matrices in serial applications and that there are no references to processors, communications or anything specific to parallel programming.

While **Aztec** simplifies the parallel coding associated with structured problems, it is for unstructured problems that **Aztec** makes a significant programming difference. To illustrate this, a 2D finite element example is given where the underlying grid is a triangulation of a complex geometry. Unlike the previous example create_matrix defines a sparsity pattern (i.e. *bindx*) but not the actual nonzero entries (i.e. *val*) as interprocessor communication is required before they can be computed. Thus, in this example **AZ_transform** takes the sparsity pattern and initializes the communication data structures. Using these structures, communication can be performed at a later stage in computing the matrix nonzeros.

Figure 9 depicts create_matrix while Figure 10 depicts an additional function ma-

Example

```
void create_matrix(bindx, val, update, N_update);
{
  read_triangles(T, N_triangles);
  init_msr(val, bindx, N_update);

  for (triangle = 0; triangle < N_triangles; triangle = triangle + 1)
    for (i = 0; i < 3; i = i + 1) {
      row = AZ_find_index(T[triangle][i], update,N_update);
      for (j = 0; j < 3; j = j + 1) {
        if (row != NOT_FOUND)
          add_to_element(row, T[triangle][j], 0.0, val, bindx, i==j);
      }
    }
  }
  compress_matrix(val, bindx, N_update);
}
```

FIG. 9. create_matrix *for the Poisson finite element problem.*

trix_fill that must be included before AZ_solve is invoked in Figure 1. We have not made any effort to optimize these routines. In both figures the new lines that have been added specifically for a parallel implementation are underlined. That is, create_matrix and matrix_fill have been created by taking a serial program that creates the finite element discretization, splitting this program over the two functions and adding a few new lines necessary for the parallel implementation. The only additional change is to replace the single data file containing the triangle connectivity read using read_triangles

by a set of data files containing the triangle connectivity for each processor. We do not discuss the details of this program but only wish to draw the readers attention to the small number of lines that need changing to convert the serial unstructured application to parallel. Most of the main routines such as setup_Ke which computes the element contributions and add_to_element which stores the element contributions in the MSR data structures remain the same. In fact, almost all the new lines of code correspond to adding the communication (AZ_exchange_bdry) (which was the main reason that the calculation of the matrix nonzeros was deferred) and the conversion of global index values by local index values with the help of AZ_find_index. As in the Poisson example, all of the details with respect to communication are hidden from the user.

6. Advanced Topics.

6.1. Data Layout. The **Aztec** function AZ_transform initializes the integer array *data_org*. This array specifies how the matrix is set up on the parallel machine. In many cases, the user need not be concerned with the contents of this array. However, in some situations it is useful to initialize these elements without the use of AZ_transform, to access these array elements (e.g. determine how many *internal* components are used), or to change these array elements (e.g. when reusing factorization information, see Section 6.2). When using the transformation software, the user can ignore the size of *data_org* as it is allocated in AZ_transform. However, when this is not used, *data_org* must be allocated of size AZ_COMM_SIZE + number of vector elements sent to other processors during matrix-vector multiplies. The contents of *data_org* are as follows:

Specifications

data_org[AZ_matrix_type]	Specifies matrix format.
AZ_VBR_MATRIX	Matrix corresponds to VBR format.
AZ_MSR_MATRIX	Matrix corresponds to MSR format.
data_org[AZ_N_internal]	Number of elements updated by this processor that can be computed without information from neighboring processors (*N_internal*). This also corresponds to the number of internal rows assigned to this processor.
data_org[AZ_N_border]	Number of elements updated by this processor that use information from neighboring processors (*N_border*).
data_org[AZ_N_external]	Number of *external* components needed by this processor (*N_external*).
data_org[AZ_N_int_blk]	Number of internal VBR block rows owned by this processor. Set to data_org[AZ_N_internal] for MSR matrices.
data_org[AZ_N_bord_blk]	Number of border VBR block rows owned by this processor. Set to data_org[AZ_N_border] for MSR matrices.

Example

```
void matrix_fill(bindx, val, N_update, update, update_index,
                 N_external, external, extern_index)

  /* read the x and y coordinates from an input file */

  for (i = 0; i < N_update; i = i + 1){
    read_from_file(x[update_index[i]], y[update_index[i]]);
  }
  AZ_exchange_bdry(x);
  AZ_exchange_bdry(y);

  /* Locally renumber the rows and columns of the new sparse matrix */

  for (triangle = 0; triangle < N_triangles; triangle = triangle + 1)
    for (i = 0; i < 3; i = i + 1) {
      row = AZ_find_index(T[triangle][i], update, N_update);
      if (row == NOT_FOUND) {
        row = AZ_find_index(T[triangle][i], external, N_external);
        T[triangle][i] = extern_index[row];
      }
      else T[triangle][i] = update_index[row];
    }
  }

  /* Fill the element stiffness matrix Ke */

  for (triangle = 0; triangle < N_triangles; triangle = triangle + 1){
    setup_Ke(Ke, x[T[triangle][0]], y[T[triangle][0]],
                 x[T[triangle][1]], y[T[triangle][1]],
                 x[T[triangle][2]], y[T[triangle][2]]);

  /* Fill the sparse matrix by scattering Ke to appropriate locations */

    for (i = 0; i < 3; i = i + 1) {
      for (j = 0; j < 3; j = j + 1){
        if (T[triangle][i] < N_update){
          add_to_element(T[triangle][i], T[triangle][j], Ke[i][j],
                         val, bindx, i==j);
        }
      }
    }
  }
```

FIG. 10. matrix_fill *for the Poisson finite element problem.*

*data_org[*AZ_N_ext_blk*]*	Number of external VBR block rows on this processor. Set to data_org[AZ_N_external] for MSR matrices.
*data_org[*AZ_N_neigh*]*	Number of processors with which we exchange information (send or receive) in performing matrix-vector products.
*data_org[*AZ_total_send*]*	Total number of vector elements sent to other processors during matrix-vector products.
*data_org[*AZ_name*]*	Name of the matrix. This name is utilized when deciding which previous factorization to use as a preconditioner (see Section 6.2). (positive integer value).
*data_org[*AZ_neighbors*]*	Start of vector containing node i.d.'s of neighboring processors. That is, *data_org[*AZ_neighbors*+i]* gives the node i.d. of the *(i+1)*'th neighbor.
*data_org[*AZ_rec_length*]*	Start of vector containing the number of elements to receive from each neighbor. We receive from the *(i+1)*'th neighbor *data_org[*AZ_rec_length*+i]* elements.
*data_org[*AZ_send_length*]*	Start of vector containing the number of elements to send to each neighbor. We send to the *(i+1)*'th neighbor *data_org[*AZ_rec_length*+i]* elements.
*data_org[*AZ_send_list*]*	Start of vector indicating the elements that we will send to other processors during communication. The first *data_org[*AZ_send_length*]* components correspond to the elements for the first neighbor and the next *data_org[*AZ_send_length*+1]* components correspond to element indices for the second neighbor, and so on.

6.2. Reusing factorizations. When solving a problem, **Aztec** may create certain information that can be reused later. In most cases, this information corresponds to either matrix scaling factors or preconditioning factorization information for LU or ILU. This information is saved internally and referenced by the matrix name given by *data_org[*AZ_name]. By changing *options[*AZ_pre_calc*]* and *data_org[*AZ_name*]* a number of different **Aztec** possibilities can be realized. As an example, consider the following situation. A user needs to solve the linear systems in the order shown below:

$$A_1 x = b, A_2 y = x, \text{ and } A_1 z = y.$$

The first and second systems are solved with *options[*AZ_pre_calc*]* set to AZ_calc. However, the name (i.e. *data_org[*AZ_name*]*) is changed between these two solves. In this way, scaling and preconditioning information computed from the first solve is not overwritten during the second solve. By then setting *options[*AZ_pre_calc*]* to AZ_reuse and *data_org[*AZ_name*]* to the name used during the first solve, the third system is solved reusing the scaling information (to scale the right hand side, initial guess, and rescale

the final solution[3]) and the preconditioning factorizations (e.g. ILU) used during the first solve. While in this example the same matrix system is solved for the first and third solve, this is not necessary. In particular, preconditioners can be reused from previous nonlinear iterates even though the linear system being solved are changing. Of course, many times information from previous linear solves is not reused. In this case the user must explicitly free the space associated with the matrix or this information will remain allocated for the duration of the program. Space is cleared by invoking AZ_clear(*data_org*[AZ_name]).

6.3. Important Constants. **Aztec** uses a number of constants which are defined in the file `az_aztec_defs.h`. Most users can ignore these constants. However, there may be situations where they should be changed. Below is a list of these constants with a brief description:

AZ_MAX_NEIGHBORS	Maximum number of processors with which information can be exchanged during matrix-vector products.
AZ_MSG_TYPE AZ_NUM_MSGS	All message types used inside **Aztec** lie between AZ_MSG_TYPE and AZ_MSG_TYPE + AZ_NUM_MSGS - 1.
AZ_MAX_BUFFER_SIZE	Maximum message information that can be sent by any processor at any given time before receiving. This is used to subdivide large messages to avoid buffer overflows.
AZ_MAX_MEMORY_SIZE	Maximum available memory. Used primarily for the LU-factorizations where a large amount of memory is first allocated and then unused portions are freed after factorization.
AZ_TEST_ELE	Internal algorithm parameter that can effect the speed of the AZ_find_procs_for_externs calculation. Reduce AZ_TEST_ELE if communication buffers are exceeded during this calculation.

6.4. AZ_transform Subtasks. The function AZ_transform described in Section 4 is actually made up of 5 subtasks. In most cases the user need not be concerned with the individual tasks. However, there might arise situations where additional information is available such that some of the subtasks can be omitted. In this case, it is possible for the user to edit the code for AZ_transform located in the file `az_tools.c` to suit the application. In this section we briefly describe the five subroutines which make up the transformation function. More detailed descriptions are given in [5]. Prototypes for these subroutines as well as for AZ_transform are given in Section 7.

AZ_transform begins by identifying the *external* set needed by each processor. Here, each column entry must correspond to either an element updated by this processor or

[3] The matrix does not need to be rescaled as the scaling during the first solve overwrites the original matrix.

an *external* component. The function AZ_find_local_indices checks each column entry. If a column is in *update*, its number is replaced by the appropriate index into *update* (i.e. *update[new column index]* = old column index). If a column number is not found in *update*, it is stored in the *external* list and the column number is replaced by an index into *external* (i.e. *external[new column index - N_update]* = old column index).

AZ_find_procs_for_externs queries the other processors to determine which processors update each of its *external* components. The array *extern_proc* is set such that *extern_proc[i]* indicates which processor updates *external[i]*.

AZ_order_ele reorders the *external* components such that elements updated by the same processor are contiguous. This new ordering is given by *extern_index* where *extern_index[i]* indicates the local numbering of *external[i]*. Additionally, *update* components are reordered so the *internal* components precede the *border* components. This new ordering is given by *update_index* where *update_index[i]* indicates the local numbering of *update[i]*.

AZ_set_message_info initializes *data_org* (see Section 6.1) This is done by computing the number of neighbors, making a list of the neighbors, computing the number of values to be sent and received with each neighbor and computing the list of elements which will be sent to other processors during communication steps.

Finally, AZ_reorder_matrix permutes and reorders the matrix nonzeros so that its entries correspond to the newly reordered vector elements.

7. Aztec Functions . In this section we describe the **Aztec** functions available to the user. Certain variables appear many times in the parameter lists of these frequently used functions. In the interest of brevity we describe these variables at the beginning of this section and then proceed with the individual function descriptions.

Frequently Used Aztec Parameters

data_org	Array describing the matrix format (Section 6.1). Allocated and set AZ_set_message_info and AZ_transform.
extern_index	*extern_index[i]* gives the local numbering of global element *external[i]*. Allocated and set by AZ_order_ele and AZ_transform.
extern_proc	*extern_proc[i]* is updating processor of *external[i]*. Allocated and set by AZ_find_procs_for_externs.
external	Sorted list (global indices) of external elements on this node. Allocated and set by AZ_find_local_indices and AZ_transform .
N_external	Number of *external* components. Set by AZ_find_procs_for_externs and AZ_transform.
N_update	Number of *update* components assigned to this processor. Set by AZ_read_update.
options, params	Arrays describing AZ_solve options (Section 2).

proc_config[AZ_node]	Node i.d. of this processor.
proc_config[AZ_N_proc]	Total number of processors used in current simulation. Allocated and set by AZ_processor_info.
update_index	*update_index[i]* gives the local numbering of global element *update[i]*. Allocated and set by AZ_order_ele and AZ_transform.
update	Sorted list of elements (global indices) to be updated on this processor. Allocated and set by AZ_read_update.
val, bindx, bpntr, cpntr, indx, rpntr	Arrays used to store matrix. For MSR matrices *bpntr, cpntr, indx, rpntr* are ignored (Section 3).

Prototype

void AZ_broadcast(char **ptr*, int *length*, int **proc_config*, int *action*)

Description

Used to concatenate a buffer of information and to broadcast this information from processor 0 to the other processors. The four possible actions are

- *action* == AZ_PACK
  - proc_config[AZ_node] == 0: store *ptr* in the internal buffer.
  - proc_config[AZ_node] $\neq$ 0: read from the internal buffer to *ptr*. If the internal buffer is empty, first receive the broadcast information.
- *action* == AZ_SEND
  - proc_config[AZ_node] == 0: broadcast the internal buffer (filled by AZ_broadcast) and then clear it.
  - proc_config[AZ_node] $\neq$ 0: clear internal buffer.

Sample Usage:

The following code fragment broadcasts the information in 'a' and 'b'.

```
if (proc_config[AZ_node] == 0) {
   a = 1;
   b = 2;
}
AZ_broadcast(&a, sizeof(int), proc_config, AZ_PACK);
AZ_broadcast(&b, sizeof(int), proc_config, AZ_PACK);
AZ_broadcast(NULL ,       0 , proc_config, AZ_SEND);
```

NOTE: There can be no other communication calls between the **AZ_PACK** and **AZ_SEND** calls to AZ_broadcast.

Parameters

ptr	On input, data string of size *length*. Information is either stored to or retrieved from *ptr* as described above.
length	On input, length of *ptr* to be broadcast/received.
action	On input, determines AZ_broadcast behavior.

Prototype

```
int AZ_check_input(int *data_org, int *options, double *params, int *proc_config)
```

Description

Perform checks for iterative solver library. This is to be called by the user of the solver library to check the values in *data_org*, *options*, *params*, and *proc_config*. If all the values are valid AZ_check_input returns 0, otherwise it returns an error code which can be deciphered using AZ_print_error.

Prototype

```
void AZ_check_msr(int *bindx, int N_update, int N_external, int option,
                  int *proc_config)
```

Description

Check that the number of nonzero off-diagonals in each row and that the column indices are nonnegative and not too large (see *option*).

Parameters

option

AZ_LOCAL	On input, indicates matrix uses local indices. The number of nonzeros in a row and the largest column index must not exceed the total number of elements on this processor.
AZ_GLOBAL	On input, indicates matrix uses global indices. The number of nonzeros in a row and the largest column index must not exceed the total number of elements in the simulation.

Prototype

```
void AZ_check_vbr(int N_update, int N_external, int option, int *bindx,
                  int *bpntr, int *cpntr, int *rpntr, int *proc_config )
```

Description

Check VBR matrix for the following:

- number of columns within each block column is nonnegative.
- *rpntr[i]* == *cpntr[i]* for $i \leq$ *N_update*.
- number of nonzero blocks in each block row is nonnegative and not too large.
- block column indices are nonnegative and not too large.

Parameters

option

AZ_LOCAL	On input, indicates matrix uses local indices. The number of block nonzeros in a row and the largest block column index must not exceed the total number of blocks columns on this processor.
AZ_GLOBAL	On input, indicates matrix uses global indices. The number of block nonzeros in a row and the largest block column index must not exceed the total number of blocks rows in the simulation.

Prototype

```
int AZ_defaults(double *options, int *params )
```

Description

Set *options* and *params* so that the default options are chosen.

Parameters

options	On output, set to the default options.
params	On output, set to the default parameters.

Prototype

```
void AZ_exchange_bdry(double *x, int *data_org)
```

Description

Locally exchange the components of the vector x so that the *external* components of x are updated.

Parameters

x	On input, vector defined on this processor. On output, *external* components of x are updated via communication.

Prototype

```
int AZ_find_index(int key, int *list, int length )
```

Description

Returns the index, i, in *list* (assumed to be sorted) which matches the key (i.e. *list[i]* == *key*). If *key* is not found AZ_find_index returns -1. See also AZ_quick_find.

Parameters

key	On input, element to be search for in list.
list	On input, sorted list to be searched.
length	On input, length of list.

Prototype

```
void AZ_find_local_indices(int N_update, int *bindx, int *update,
                           int **external, int *N_external, int mat_type,
                           int *bpntr)
```

Description

Given the global column indices for a matrix and a list of elements updated on this processor, compute the *external* set and change the global column indices to local column indices. Specifically,

- allocate *external*, compute and store the external components in *external*.
- renumber column indices so that column entry k is renumbered as j where either *update[j]* == k or *external[j-N_update]* == k .

Called by AZ_transform.

Parameters

mat_type	On input, indicates whether matrix format is MSR (= AZ_MSR_MATRIX) or VBR (= AZ_VBR_MATRIX).
external	On output, allocated and set to sorted list of the external elements.
bindx	On input, contains global column numbers of MSR or VBR matrix (Section 3). On output, contains local column numbers as described above.

Prototype

```
void AZ_find_procs_for_externs(int N_update, int *update, int *external,
                               int N_external, int *proc_config, int **extern_proc)
```

Description

Determine which processors are responsible for updating each external element.
Called by AZ_transform.

Parameters

extern_proc	On output, *extern_proc[i]* contains the node number of the processor which updates *external[i]*.

Prototype

void AZ_free_memory(int *name*)

Description

Free **Aztec** memory associated with matrices with *data_org*[AZ_name] = *name*. This is primarily scaling and preconditioning information that has been computed on earlier calls to AZ_solve.

Parameters

name	On output, all preconditioning and scaling information is freed for matrices which have *data_org*[AZ_name] = *name*.

Prototype

double AZ_gavg_double(double *value*, int **proc_config*)

Description

Return the average of the numbers in *value* on all processors.

Parameters

value	On input, *value* contains a double precision number.

Prototype

double AZ_gdot(int N, double *r, double *z, int **proc_config*)

Description

Return the dot product of r and z with unit stride. This routine calls the BLAS routine ddot to do the local vector dot product and then uses the global summation routine AZ_gsum_double to obtain the required global result.

Parameters

N	On input, length of r and z on this processor.
r, z	On input, vectors distributed over all the processors.

Prototype

```
double AZ_gmax_double(double value, int *proc_config )
```

Description

Return the maximum of the numbers in *value* on all processors.

Parameters

value	On input, *value* contains a double precision number.

Prototype

```
int AZ_gmax_int(int value, int *proc_config )
```

Description

Return the maximum of the numbers in *value* on all processors.

Parameters

value	On input, *value* contains an integer.

Prototype

```
double AZ_gmax_matrix_norm(double *val, int *indx, int *bindx, int *rpntr, int *cpntr,
                           int *bpntr, int *proc_config, int *data_org)
```

Description

Returns the maximum matrix norm $\|A\|_\infty$ for the distributed matrix encoded in *val, indx, bindx, rpntr, cpntr, bpntr* (Section 3).

Prototype

```
double AZ_gmax_vec(int N, double *vec, int *proc_config )
```

Description

Return the maximum of all the numbers located in *vec[i]* ($i < N$) on all processors.

Parameters

vec	On input, *vec* contains a list of numbers.
N	On input, length of *vec*.

Prototype

```
double AZ_gmin_double(double value, int *proc_config )
```

Description

Return the minimum of the numbers in *value* on all processors.

Parameters

value	On input, *value* contains a double precision number.

Prototype

```
int AZ_gmin_int(int value, int *proc_config )
```

Description

Return the minimum of the numbers in *value* on all processors.

Parameters

value	On input, *value* contains an integer.

Prototype

```
double AZ_gsum_double(double value, int *proc_config )
```

Description

Return the sum of the numbers in *value* on all processors.

Parameters

value	On input, *value* contains a double precision number.

Prototype

```
int AZ_gsum_int(int value, int *proc_config )
```

Description

Return the sum of the integers in *value* on all processors.

Parameters

value	On input, *value* contains an integer.

Prototype

```
void AZ_gsum_vec_int(int *values, int *wkspace, int length, int *proc_config )
```

Description

values[i] is set to the sum of the input numbers in *values[i]* on all processors ($i <$ *length*).

Parameters

values	On input, *values* contains a list of integers. On output, *values[i]* contains the sum of the input *values[i]* on all the processors.
wkspace	On input, workspace array of size *length*.
length	On input, length of *values* and *wkspace*.

Prototype

```
double AZ_gvector_norm(int n, int p, double *x, int *proc_config)
```

Description

Returns the p norm of the vector x distributed over the processors:

$$\|x\|_p = (x[0]^p + x[1]^p + \cdots + x[N-1]^p)^{1/p}$$

where N is the total number of elements in x over all processors.
NOTE: For the $\|\cdot\|_\infty$ norm, set $p = -1$.

Parameters

n	On input, number of *update* components of x on this processor.
p	On input, order of the norm to perform, i.e., $\|x\|_p$.
x	On input, vector whose norm will be computed.

Prototype

```
void AZ_init_quick_find(int *list, int length, int *shift, int *bins )
```

Description

shift and *bins* are set so that they can be used with AZ_quick_find. On output, shift satisfies

$$\frac{range}{2^{shift-1}} > \left\lfloor \frac{length}{4} \right\rfloor \quad \text{and} \quad \frac{range}{2^{shift}} \leq \left\lfloor \frac{length}{4} \right\rfloor$$

where *range = list[length - 1] - list[0]*. The array *bins* must be of size $2 + length/4$ and is set so that

$$bins[k] \leq list[j] < bins[k+1]$$

where $k = (list[j] - list[0])/2^{shift}$.
This routine is used in conjunction with AZ_quick_find. The idea is to use *bins* to get a good initial guess as to the location of *value* in *list*.

Parameters

list	On input, sorted *list*.
length	On input, length of *list*.
shift	On output, *shift* is set as described in above.
bins	On input, array of size $2 + length/4$. On output, *bins* is set as described above.

Prototype

```
void AZ_matvec_mult(double *val, int *indx, int *bindx, int *rpntr, int *cpntr,
                    int *bpntr, double *b, double *c, int exchange_flag,
                    int *data_org )
```

Description

Perform the matrix-vector multiply

$$c \leftarrow Ab$$

where the matrix A is encoded in *val, indx, bindx, rpntr, cpntr, bpntr* (Section 3).

Parameters

b	On input, distributed vector to use in multiplication.

c	On output, result of matrix-vector multiplication.
exchange_flag	On input, dictates whether communication needs to occur. If *exchange_flag* == 1, communication occurs. If *exchange_flag* == 0, no communication occurs.

Prototype

```
void AZ_msr2vbr(double *val, int *indx, int *rpntr, int *cpntr, int *bpntr, int *bindx,
                int *bindx2, double *val2, int total_blk_rows, int total_blk_cols,
                int blk_space, int nz_space, int blk_type)
```

Description

Convert the DMSR matrix defined in $(val2, bindx2)$ to a DVBR matrix defined in $(val, indx, rpntr, cpntr, bpntr, bindx)$.

Parameters

val2, bindx2	On input, DMSR arrays holding the matrix to be converted.
cpntr	On input, $cpntr[i]$ is the block size of the i^{th} block in the resulting DVBR matrix. Columns 0 to $cpntr[0]-1$ form the first block column, columns $cpntr[0]$ to $cpntr[0] + cpntr[1]-1$ form the second block column, etc. On output, *cpntr* corresponds to the resulting DVBR matrix.
val, indx, rpntr, bpntr, bindx	On output, DVBR arrays of converted DMSR matrix.
total_blk_rows	On input, number of block rows in resulting local VBR matrix.
total_blk_cols	On input, number of block columns in resulting local VBR matrix.
blk_space	On input, length allocated for *bindx* and *indx*.
nz_space	On input, length allocated for *val*.
blk_type	On input, if *blk_type* > 0, indicates that all block rows (and columns) have the same size given by *blk_type*. If *blk_type* < 0, the block rows have different sizes.

Prototype

```
void AZ_order_ele(int *update_index, int *extern_index, int *N_internal,
                  int *N_border, int N_update, int *bpntr, int *bindx,
                  int *extern_proc, int N_external, int option, int mat_type)
```

Description

Find orderings for *update* and *external*. *external* are ordered so that elements updated by the same processor are contiguous. If *option* == AZ_ALL, *update* are ordered so that the *internal* components have the lowest numbers followed by the *border* components. Otherwise, the order of *update* is unchanged. The ordering information is placed in *update_index* and *extern_index* (Section 4). Called by AZ_transform.

Parameters

N_internal	On output, number of *internal* components on processor.
N_border	On output, number of *border* components on processor.
update_index	On output, *update_index[i]* indicates the local index (or order) of *update[i]*.
extern_index	On output, *extern_index[i]* indicates the new local index (or order) of *external[i]*.
option	On input, indicates whether to reorder *update*.
AZ_ALL	Order *update* and *external*.
AZ_EXTERNS	Order only external elements.
mat_type	On input, indicates whether matrix format is MSR (= AZ_MSR_MATRIX) or VBR (= AZ_VBR_MATRIX).

Prototype

```
void AZ_print_error(int error_code)
```

Description

Prints out an error message corresponding to *error_code*. Typically, *error_code* is generated by AZ_check_input.

Parameters

error_code	On input, error code generated by AZ_check_input.

Prototype

```
void AZ_processor_info(int *proc_config)
```

Description

proc_config[AZ_node] is set to the node name of this processor. *proc_config[AZ_N_proc]* is set to the number of processors used in simulation.

Prototype

```
int AZ_quick_find(int key, int *list, int length, int shift, int *bins )
```

Description

Return the index, i, in *list* (assumed to be sorted) which matches the key (i.e. *list[i]* = *key*). If *key* is not found AZ_quick_find returns -1.
NOTE: This version is faster than AZ_find but requires *bins* to be set and stored using AZ_init_quick_find.

Parameters

key	On input, element to search for in *list*.
list	On input, sorted list to be searched.
length	On input, length of list.
shift	On input, used for initial guess (computed by previous AZ_init_quick_find call).

bins — On input, computed by AZ_init_quick_find for initial guess. *bins* is set so that $list[bins[k]] \leq key < list[bins[k+1]]$ where $k = (key - list[0])/2^{shift}$.

Prototype

```
void AZ_read_msr_matrix(int *update, double **val, int **bindx, int N_update,
                        int *proc_config )
```

Description

Read the file `.data` and create a matrix in the MSR format. Processor 0 reads the input file. If the new row to be added resides in processor 0's *update*, it is added to processor 0's matrix. Otherwise, processor 0 determines which processor has requested this row and sends it to this processor for its local matrix.
The form of the input file is as follows:

```
num_rows
col_num1  entry1  col_num2  entry2
col_num3  entry3 -1
col_num4  entry4  col_num5  entry5
col_num6  entry6 -1
```

This input corresponds to two rows: 0 and 1. Row 0 contains `entry1` in column `col_num1`, `entry2` in column `col_num2` and `entry3` in column `col_num3`. Row 1 contains `entry4` in column `col_num4`, `entry5` in column `col_num5` and `entry6` in column `col_num6`.
NOTE: row and column numbers must start from 0.
NOTE: AZ_read_msr_matrix() is inefficient for large matrices.

Parameters

val, bindx — On output, these two arrays are allocated and filled with the MSR representation corresponding to the file `.data`.

Prototype

```
void AZ_read_update(int *N_update, int **update, int *proc_config,
                    int N, int chunk, int input_option )
```

Description

This routine initializes *update* to the global indices updated by this processor and initializes *N_update* to the total number of elements to be updated.

Parameters

N_update — On output, number of elements updated by processor.

update — On output, *update* is allocated and contains a list of elements updated by this processor in ascending order.

chunk — Number of indices within a group. For example, $chunk == 2 \Rightarrow chunk_0 = \{0, 1\}$, and $chunk_1 = \{2, 3\}$.

N — Total number of chunks in the vector.

input_option

AZ_LINEAR — Processor 0 is assigned the first $\left\lfloor \frac{N+P-1}{P} \right\rfloor$ chunks, processor 1 is assigned the next $\left\lfloor \frac{N+P-2}{P} \right\rfloor$ chunks, etc. where $P =$ proc_config[AZ_N_proc]).

AZ_BOX — The processor system is viewed as a $p_2 \times p_1 \times p_0$ where $p_i = 2^{\lfloor (k+i)/3 \rfloor}$ (so proc_config[AZ_N_proc] must equal 2^k). The chunks are viewed as an $n \times n \times n$ cube where n is divisible by each p_i. Chunks are distributed into uniform boxes such that each processor has the same number of chunks.

AZ_FILE — Read the proc_config[AZ_N_proc] lists contained in the file update. Each list contains a set of global indices preceeded by the of number of indices in this set. List 0 is sent to processor proc_config[AZ_N_proc] - 1, list 1 is sent to processor proc_config[AZ_N_proc] - 2, etc. Note: A graph partitioning package named **Chaco** [2] produces files in this format.

Prototype

```
void AZ_reorder_matrix(int N_update, int *bindx, double *val, int *update_index,
                       int *extern_index, int *indx, int *rpntr, int *bpntr,
                       int N_external, int *cpntr, int option, int mat_type)
```

Description

Reorder the matrix so that it corresponds to the new ordering given by *update_index* and *extern_index*. Specifically, global matrix entry $(update[i], update[j])$ which was stored as local matrix entry (i, j) is stored as $(update\_index[i], update\_index[j])$ on output. Likewise, global matrix entry $(update[i], external[k])$ which was stored as local matrix entry $(i, k + N\_update)$ is stored locally as (update_index[i], extern_index[k]) on output. Called by AZ_transform.
IMPORTANT: This routine assumes that *update_index* contains two sequences of numbers that are ordered but intertwined. For example,

update_index:	4	5	0	6	1	2	3	7
sequence 1:			0		1	2	3	
sequence 2:	4	5		6				7

Parameters

option	On input, indicates whether to reorder update elements.
AZ_ALL	All the rows and columns are renumbered.
AZ_EXTERNS	Only columns corresponding to external elements are renumbered.
mat_type	On input, indicates matrix format.
AZ_MSR_MATRIX	DMSR matrix format.
AZ_VBR_MATRIX	DVBR matrix format.
bindx, val, indx, rpntr, bpntr, cpntr	On input, matrix ordered as described above. On output, matrix reordered using *update_index* and *extern_index* as described above.

Prototype

```
void AZ_set_message_info(int N_external, int *extern_index, int N_update,
                         int *external, int *extern_proc, int *update,
                         int *update_index, int *proc_config, int *cpntr,
                         int **data_org, int mat_type)
```

Description

Initialize *data_org* so that local communications can occur to support matrix vector products. This includes:

- determine neighbors with which we send or receive.
- determine the total number of elements that we send and allocate *data_org*.
- initialize *data_org* as described in Section 6.1.
  Note: *data_org*[AZ_name] is set to a number (starting from 1) that is incremented each time AZ_set_message_info is called.

Called by AZ_transform.
NOTE: Implicitly the neighbors are numbered using the ordering of the external elements (which have been previously ordered such that elements updated by the same processor are contiguous).

Parameters

data_org	On output, *data_org* is allocated and completely initialized as described in Section 6.1.
mat_type	On input, indicates matrix format.
AZ_MSR_MATRIX	DMSR matrix.
AZ_VBR_MATRIX	DVBR matrix.

Prototype

```
void AZ_solve(double *x, double *b, int *options, double *params, int *indx,
              int *bindx, int *rpntr, int *cpntr, int *bpntr, double *val,
              int *data_org, double *status, int *proc_config)
```

Description

Solve the system of equations $Ax = b$ via an iterative method where the matrix A is encoded in *indx, bindx, rpntr, cpntr, bpntr* and *val* (see Section 3 and Section 2).

Parameters

x	On input x contains the initial guess. On output x contains the solution to linear system.
b	Right hand side of linear system.
options, params	Options and parameters used during the solution process (Section 2).
status	On output, status of iterative solver (Section 2).

Prototype

```
void AZ_sort(int *list1, int N, int *list2, double *list3 )
```

Description

Sort the elements in *list1*. Additionally, move the elements in *list2* and *list3* so that they correspond with the moves done to *list*. NOTE: If *list2* == NULL, *list2* is not manipulated. If *list3* == NULL, *list3* is not manipulated.

Parameters

list1	On input, values to be sorted. On output, sorted values (i.e. *list1[i]* $\leq$ *list1[i+1]*)
N	On input, length of lists to be sorted.
list2	On input, a list associated with *list1*. On output, if *list1[k]* on input is now stored in *list1[j]* on output, *list2[k]* on input is also stored as *list2[j]* on output.
list3	On input, a list associated with *list1*. On output, if *list1[k]* on input is now stored in *list1[j]* on output, *list3[k]* on input is also stored as *list3[j]* on output. Note: if *list3* == NULL on input, it is unchanged on output.

Prototype

```
void AZ_transform(int *proc_config, int **external, int *bindx,
                  double *val, int *update, int **update_index,
                  int **extern_index, int **data_org, int N_update,
                  int *indx, int *bpntr, int *rpntr, int **cpntr, int mat_type)
```

Description

Convert the global matrix description to a distributed local matrix format (see Section 2 and Section 6.4).

Parameters

external	On output, allocated and set to components that must be communicated during the matrix vector multiply.
bindx, val, index, bpntr, rpntr	On input, matrix arrays (MSR or VBR) corresponding to global format. On output, matrix arrays (DMSR or DVBR) corresponding to local format. See Section 2.
update_index	On output, allocated and set such that *update_index*[*i*] is the local numbering corresponding to *update*[*i*].
extern_index	On output, allocated and set such that *extern_index*[*i*] is the local numbering corresponding to *external*[*i*].
data_org	On output, allocated and set to data layout information, see Section 6.1.
cpntr	On output, allocated and set for VBR matrices to the column pointer array.
mat_type	On input, matrix format: either AZ_VBR_MATRIX or AZ_MSR_MATRIX.

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