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SIAM Applied Linear Algebra  
October 2015

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## Outline

We consider solving the linear system of equations,

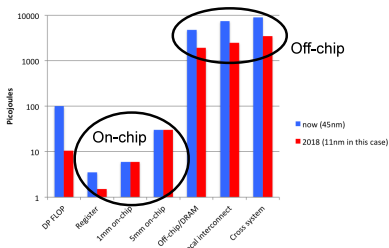
$$Ax = b,$$

where  $A$  is large and non-symmetric.

- ▶ Many applications: e.g., scientific/engineering applications when solving PDEs
- ▶ Communication-avoiding Krylov method:
  - GMRES for solving large-scale problems
- ▶ Communication-Avoiding Preconditioners for CA methods
  - A domain decomposition framework for CA preconditioning
- ▶ Hybrid CPU/GPU cluster implementation

# Communication-Avoiding Methods

- ▶ **Communication:**
  - Moving data between levels of memory
  - Moving data between processors in a network
- ▶ **Communication-Avoiding:** Reduce Communication (messages, volume)
  - Not Communication hiding
- ▶ Improves Time to solution and Reduces energy consumption
- ▶ More important in future architectures



(Image Courtesy: John Shalf, LBL)

## Communication-Avoiding Iterative Methods

- ▶ Originally proposed 30 years ago for Conjugate Gradient (J. van Rosendale, 1983).
- ▶ Chronopoulos and Gear - “s-step iterative methods” (1989)
- ▶ R. Leland - The effectiveness of these methods (1989)
- ▶ Walker - Implementation of the GMRES method using Householder transformations (1988)
- ▶ E. de Sturler and H. A. van der Vorst - GMRES and CG, basis vectors (2005)
- ▶ M. Hoemmen (2010) - TSQR, “Communication-Avoiding” methods
- ▶ Two main problems:
  - “Good” basis vectors (works for practical ‘s’)
  - Lack of preconditioners (This talk)



## Preconditioners for Communication-Avoiding Iterative Methods

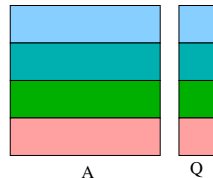
- ▶ Preconditioners that are like SpMV or that add no communication
  - Polynomial preconditioning, Sparse approximate inverse
  - CA-ILU(0) (L. Grigori and S. Moufawad, 2013)
  - Deflation based preconditioning (E. Carson 2014)
- ▶ Preconditioners that use low-rank like structures might be an option
  - Need changes to how the matrix is stored and no known evaluation with s-step methods
- ▶ Other related methods
  - s-step GMRES as bottom solver for multigrid (IPDPS 14)
  - Communication hiding pipelined Krylov methods do not have the preconditioning problem (P. Ghysels et al., 2013)
  - Hierarchical Krylov Methods [L. McInnes et al.]

## Restarted GMRES with GPUs

```

1 Generate Krylov Basis on GPUs:  $O(m \cdot nnz(A) + m^2 n)$  flops
  for  $j = 1, 2, \dots, m$  do
    Sparse Matrix-Vector Multiply (SpMV (+ Precond)):
       $\mathbf{q}_{j+1} := A\mathbf{q}_j$ 
    Orthonormalization (Orth):
       $\mathbf{q}_{j+1} := \mathbf{q}_{j+1} - Q_{1:j}Q_{1:j}^T \mathbf{q}_{j+1}$ 
  end for
2 Solve Projected Subsystem on CPUs:  $O(m^2)$  flops
  small structured least-square problem
  → restart with “best” initial vector  $\mathbf{q}_1$  in  $Q_{1:m}$ 

```



- ▶ generating basis vectors dominates computational cost.
  - ▶ distribute  $A$  and  $Q$  in a 1D block row among GPUs.
  - ▶ redundantly solve least-squares by each process.
- ▶ both *SpMV* and *Orth* require “expensive” communication:
  - ▶ point-to-point/neighborhood for *SpMV* (inter-GPU).
  - ▶ global all-reduces in *Orth* (inter-GPU).
  - ▶ data movements through local memory hierarchy (intra-GPU).

## Communication-Avoiding Implementation of $s$ -step GMRES

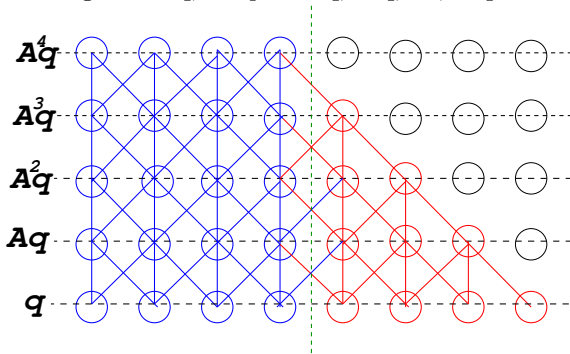
```
1. Generate Krylov Basis:
  for  $j = 1, 1 + s, \dots, m$  do
    Matrix Powers Kernel (MPK):
       $\mathbf{q}_{k+1} := A\mathbf{q}_k$ , for  $k = j, \dots, j + s - 1$ 

    Block Orthogonalization (BOrth):
      orthogonalize  $Q_{j+1:j+s}$  against  $Q_{1:j}$ 
    Tall-skinny QR (TSQR):
      orthogonalize  $Q_{j+1:j+s}$ 
      compute  $H_{j:j+s-1, j+1:j+s}$ 
  end for
2. Solve Projected Subsystem on CPUs:  $\sim O(m^2)$  flops.
   structured small least-square problem
    $\rightarrow$  restart with “best” initial vector  $\mathbf{q}_1$  in  $Q_{1:m}$ .
```

- ▶ replace *SpMV* and *Ortho* with *MPK* and *BOrth*+*TSQR*.
- ▶ reduce comm by generating  $s$  vectors “at once”  
(e.g., replace BLAS-2 with BLAS-3).

## Matrix Powers Kernel for a tridiagonal matrix

For a given starting vector  $\mathbf{q}$ , compute  $A\mathbf{q}, A^2\mathbf{q}, \dots, A^s\mathbf{q}$  (e.g.,  $s = 4$ ):

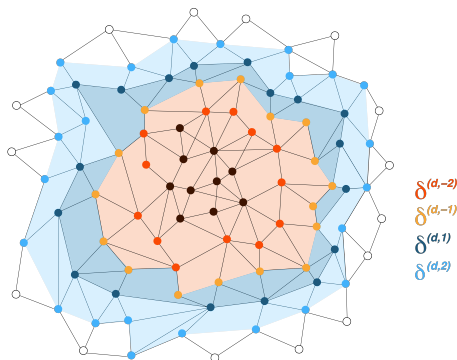


1. communicate required nonlocal elements for  $s$ -step between GPUs
2. apply  $s$  *SpMV*s with extra computation on shrinking ghost
  - local submatrix is expanded with  $s$ -level ghost

→ reduce inter-GPU latency by  $s$  (with redundant computation).

## Matrix Powers Kernel for a general matrix ( $s = 2$ ):

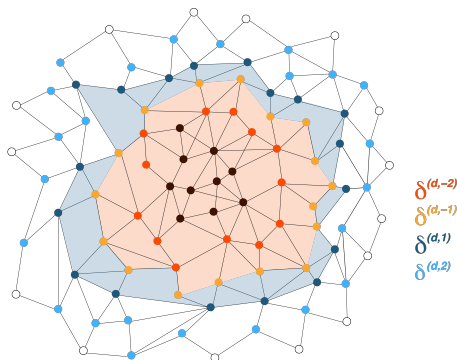
In adjacency graph of  $A$ ,



- at 1st step of *MPK*,  
we perform *SpMV* with *local* and 2-level *ghost* elements of  $\mathbf{q}_1$

## Matrix Powers Kernel for a general matrix ( $s = 2$ ):

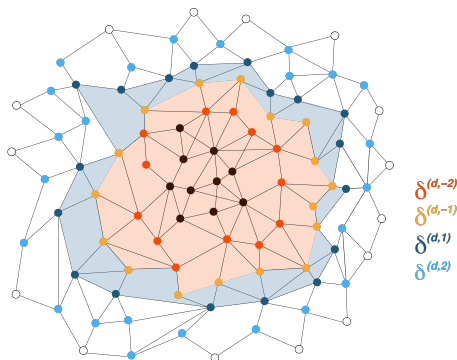
In adjacency graph of  $A$ ,



- ▶ at 1st step of *MPK*,  
we perform *SpMV* with **local** and 2-level **ghost** elements of  $\mathbf{q}_1$   
→ compute **local** and 1-level **ghost** elements of  $\mathbf{q}_2$

## Matrix Powers Kernel for a general matrix ( $s = 2$ ):

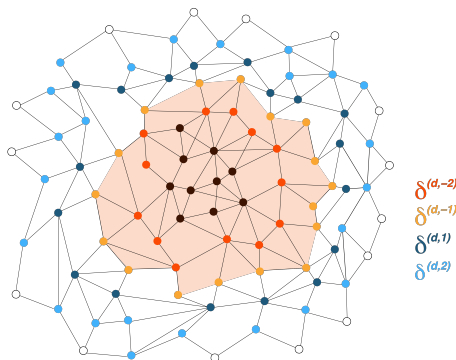
In adjacency graph of  $A$ ,



- ▶ at 2nd step of *MPK*, we perform *SpMV* with local and 1-level ghost elements of  $\mathbf{q}_2$

## Matrix Powers Kernel for a general matrix ( $s = 2$ ):

In adjacency graph of  $A$ ,



- at 2nd step of *MPK*,  
we perform *SpMV* with **local** and 1-level **ghost** elements  
→ compute **local** elements of  $\mathbf{q}_3$



## Our Matrix Powers Kernel Implementation with multiple GPUs

Initialize *MPK*:

set up communication pattern.

expand *local* submatrix with *ghost* elements, etc.

CA-GMRES with GPUs.

1. Generate Krylov Basis:

for  $j = 1, 1 + s, \dots, m$  do

*MPK*:

*Inter-GPU Communication*: each MPI process

1. CPU  $\leftarrow$  GPUs using CUDA

2. CPUs  $\longleftrightarrow$  CPUs using MPI

3. CPU  $\rightarrow$  GPUs using CUDA

*GPU Kernel*:

for  $k = 1, 2, \dots, s$  do

*SpMV* with *local* and  $k$ -level *ghost* elements

end for

*BOrth* and *TSQR*.

end for

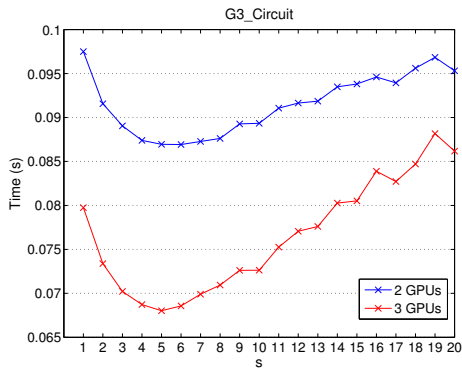
2. Solve projected system.

- currently optimized only for inter-GPU communication, and not for intra-GPU communication

## Matrix Powers Kernel Performance on a node

Our *MPK* requires overheads, but reduces inter-GPU latency:

- ▶ additional memory to store “ghost” elements
- ▶ addition computation for *SpMV* with “ghost” elements
- ▶ potentially, increasing total inter-GPU communication volume.



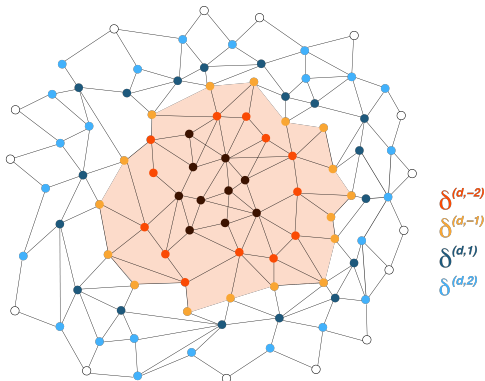
## Integrating preconditioner into *MPK*

Apply *Preco* followed by *SpMV* at each step of *MPK*

```
for  $k = j, j + 1, \dots, j + s - 1$  do  
  Preco:  $\mathbf{q}_{k+1} := M^{-1}\mathbf{q}_k$   
  SpMV:  $\mathbf{q}_{k+1} := A\mathbf{q}_{k+1}$   
end for
```

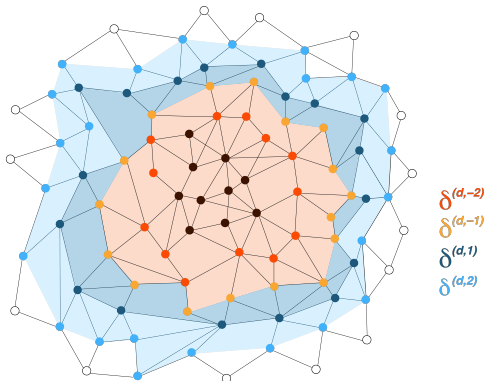
- ▶ focus on right-preconditioning, generating  $\kappa(AM^{-1}, \mathbf{q}_1)$ 
  - can be easily extended to left-preconditioning
- ▶ not increase inter-GPU comm from what is already needed by *MPK*

## Challenge: block Jacobi preconditioner increases communication



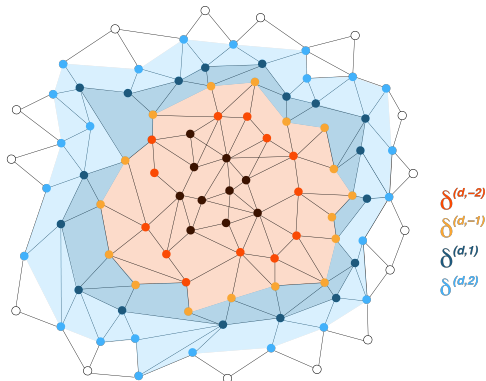
- ▶ each GPU *Precon* local elements of  $\mathbf{q}_1$ , solving its local sub-problem.
- ▶ *SpMV* requires “preconditioned”  $s$ -level ghost elements of  $\mathbf{q}_1$   
→ additional communication

## Challenge: block Jacobi preconditioner increases communication



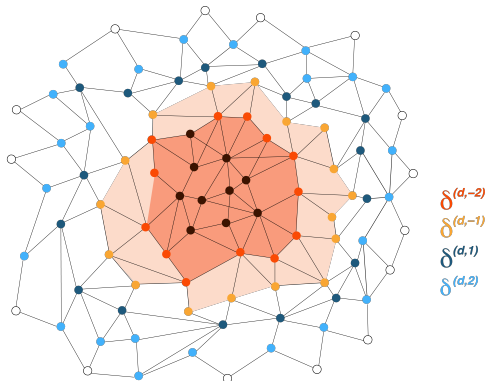
- ▶ each GPU *Precon* local elements of  $\mathbf{q}_1$ , solving its local sub-problem.
- ▶ *SpMV* requires “preconditioned”  $s$ -level ghost elements of  $\mathbf{q}_1$   
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## Challenge: block Jacobi preconditioner increases communication



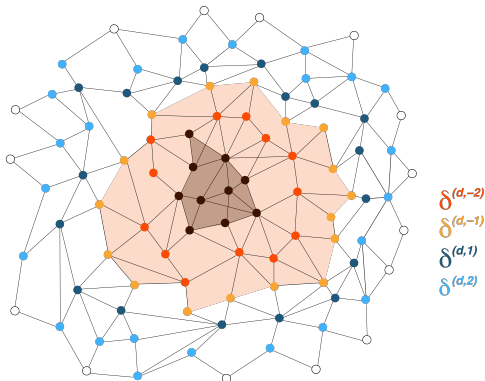
- **Solution 1:** consider  $2 \times s$  levels of **ghost** (*Preco* then *SpMV*)
- **Solution 2:** consider what we can do without additional comm

## Domain Decomposition Preconditioner for CA-Krylov



- ▶ for 1st  $SpMV$ , neighboring GPUs require elements on 1-level **underlap**
  - **local** elements reachable from other subdomains by one edge

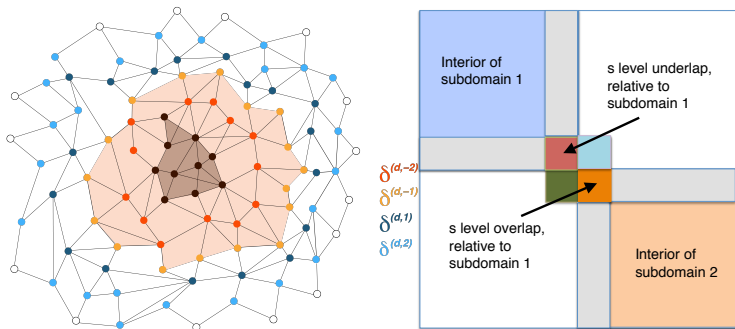
## Domain Decomposition Preconditioner for CA-Krylov



- for 2nd *SpMV*, neighboring GPUs require elements on 2-level **underlap**
  - **local** elements reachable from other subdomains by two edges



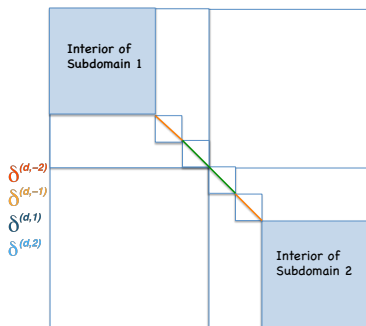
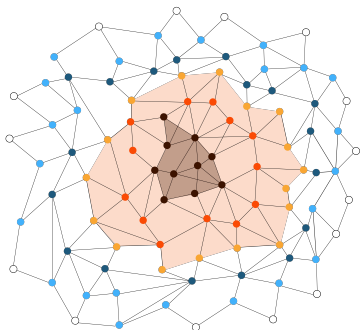
# Domain Decomposition Preconditioner for CA-Krylov



In order to “localize” effects of preconditioner,

- ▶ form “interior” by removing s-level “underlap”
- ▶ apply “local” preconditioner on “interior” and “underlap/ghost,” separately
  - ILU( $k$  or  $\tau$ ), SAI( $k$ ), Jacobi, GaussSeidel, etc. on “interior”

# Domain Decomposition Preconditioner for CA-Krylov

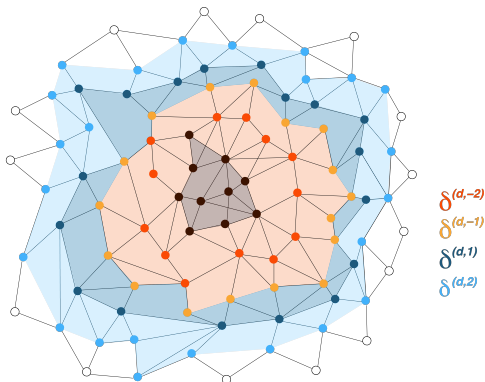


In order to “localize” effects of preconditioner,

- ▶ form “interior” by removing  $s$ -level “underlap”
- ▶ apply “local” preconditioner on “interior” and “underlap/ghost,” separately
  - ILU( $k$  or  $\tau$ ), SAI( $k$ ), Jacobi, GaussSeidel, etc. on “interior”
  - diagonal Jacobi on “underlap” and “ghost”

# Domain Decomposition Preconditioner for CA-Krylov

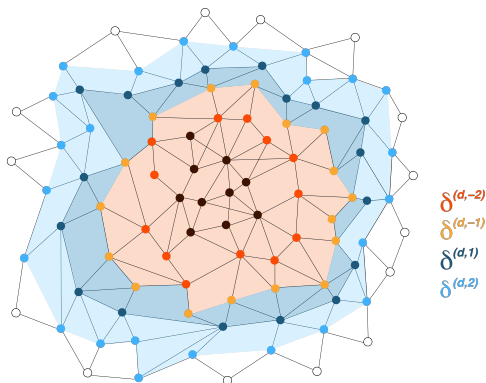
For *Precon* at 1st step of *MPK*,



- local preconditioning on interior and 2-level underlap/ghost of  $\mathbf{q}_1$ 
  - ILU( $k$  or  $\tau$ ), SAI( $k$ ), Jacobi, GaussSeidel, etc. on interior
  - diagonal Jacobi on underlap and 2-level ghost

# Domain Decomposition Preconditioner for CA-Krylov

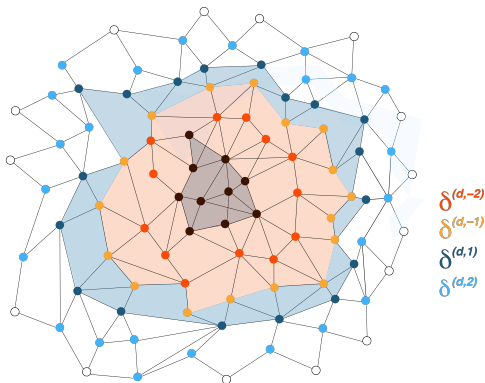
For  $SpMV$  at 1st step of  $MPK$ ,



►  $SpMV$  with local subdomain and 2-level ghost of  $\mathbf{q}_1$

## Domain Decomposition Preconditioner for CA-Krylov

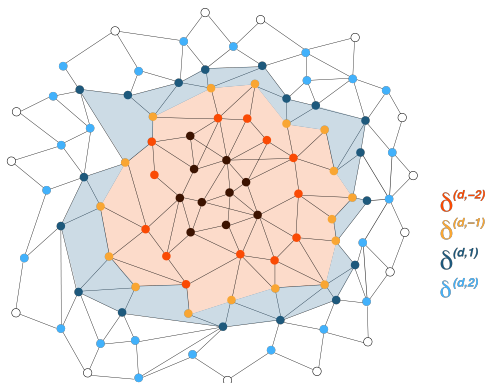
For *Precon* at 2nd step of *MPK*,



- local preconditioning on interior and 1-level underlap/ghost of  $\mathbf{q}_2$ 
  - ILU( $k$  or  $\tau$ ), SAI( $k$ ), Jacobi, GaussSeidel, etc. on interior
  - diagonal Jacobi on underlap and 1-level ghost

## Domain Decomposition Preconditioner for CA-Krylov

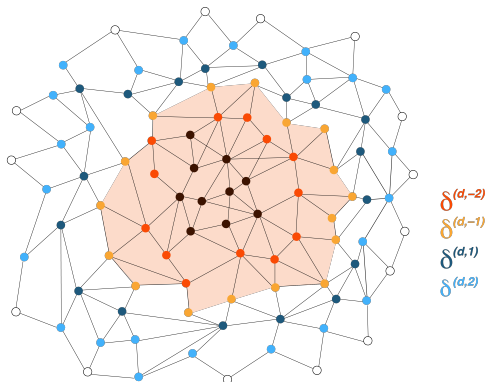
For  $SpMV$  at 2nd step of  $MPK$ ,



- $SpMV$  with local subdomain and 1-level ghost of  $q_2$

## Domain Decomposition Preconditioner for CA-Krylov

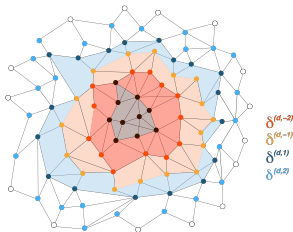
For  $SpMV$  at 2nd step of  $MPK$ ,



- apply  $SpMV$  with local subdomain and 1-level ghost  
→ compute local elements of  $\mathbf{q}_3$

# Domain Decomposition Preconditioner for CA-Krylov

## Summary:



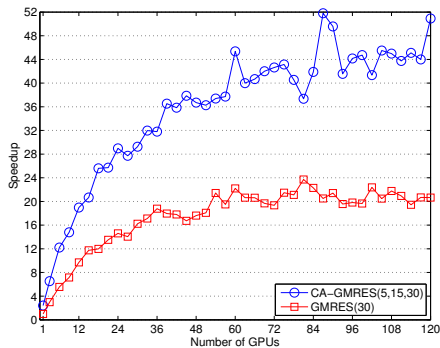
- ▶ no increase in inter-GPU communication
- ▶ any local preconditioner/solver on interior
  - ILU( $k$  or  $\tau$ ), SAI( $k$ ), Jacobi, Gauss-Seidel, etc.
- ▶ preconditioner on **underlap**/**ghost**
  - diagonal Jacobi: **interior** preconditioner propagates only within **subdomain**
  - extension in current work



## Experimental Setup

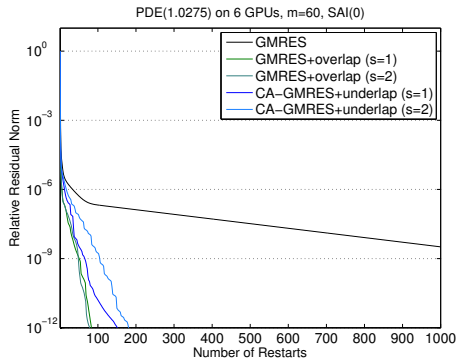
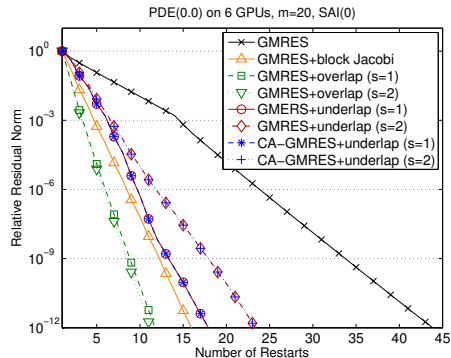
- ▶ graph partitioning (e.g., METIS) for load balance and small communication
- ▶ local matrix reordering (e.g., METIS, RCM) for performance (e.g., nested dissection for triangular solves on GPU)
- ▶ matrix equilibration for numerical stability
- ▶ Newton basis to enhance *MPK* stability
- ▶ *Precon* computed on CPU (e.g., ITSOL, ParaSail), and copied and apply it on GPU (e.g., trsv/spmv of CuSPARSE)
- ▶ Keeneland at Georgia Tech  
each node has  $2 \times 6$  Intel Xeon + 3 NVIDIA M2090.
- ▶ Test matrix:  $\text{PDE}(\alpha)$ :  $n \approx 10^6$ , symmetric but can be indefinite
  - larger  $\alpha$  makes it more ill-conditioned
  - $\alpha > 1$  makes it indefinite

## CA-GMRES Performance (speedups vs. GMRES on one GPU)



- ▶ obtained speedups of up to 2.5
  - more details in IPDPS/SC'14 papers.

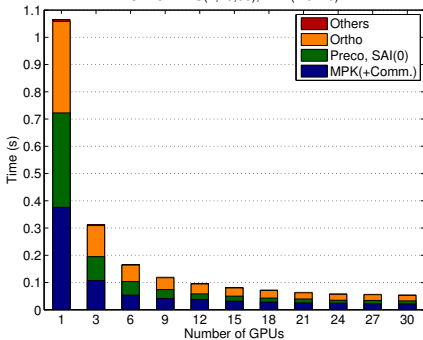
# Convergence Results



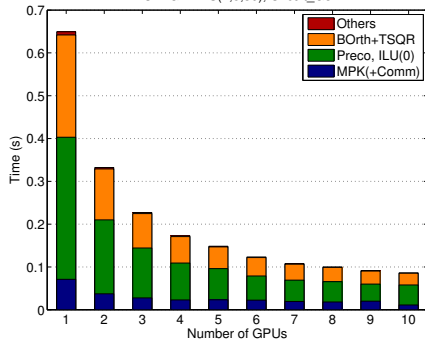
- DD preconditioner improves the convergence
  - faster convergence with a larger overlap
  - slower convergence with a larger underlap

# Restart Cycle Time Breakdown

CA-GMRES(1,10,60), PDE(1.0275)

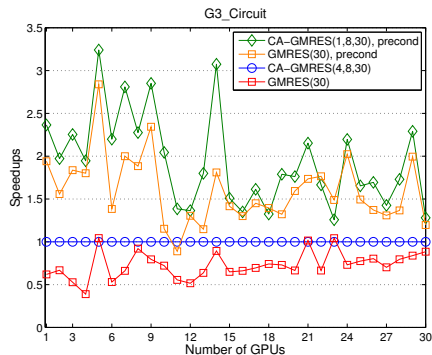
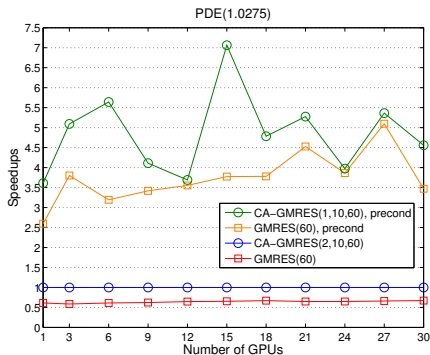


CA-GMRES(1,8,30), Circuit\_G3



- ▶ SAI(0) is used for PDE(1.0275)
- ▶ ILU(0) is required for Circuit\_G3

## Time to Solution Speedups vs. CA-GMRES



- ▶ speedups of up to  $7.5\times$  over CA-GMRES without preconditioner
- ▶ speedups of up to  $1.7\times$  over GMRES with preconditioner
  - ▶ Our *MPK* is not optimized on a GPU
  - ▶ On GPUs, *Ortho* performs great, and *SpMV/Preco* can dominate easily.

## Summary

- ▶ proposed domain decomposition preconditioners for CA-Krylov
  - ▶ do not increase inter-process communication
  - ▶ can use any solver on interior problem
- ▶ presented results of a block Jacobi like implementation
  - ▶ diagonal Jacobi on underlap/ghost
  - ▶ potential to improve convergence/performance
    - over precond GMRES or standard CA-GMRES

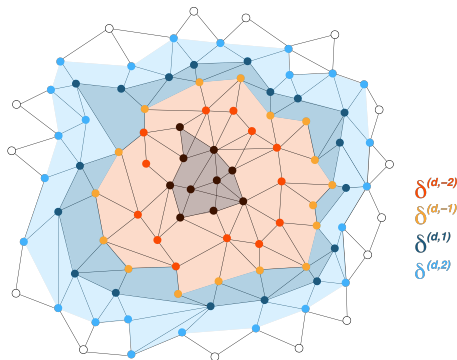
## Future work

- ▶ improving performance
  - ▶ utilizing CPU, partitioning, etc.
- ▶ underlap/ghost preconditioning
- ▶ more extensions (e.g., “flexible” preconditioner)

Thank you!!

## Domain Decomposition Preconditioner for CA-Krylov

For **SpMV** at 1st step of **MPK**,

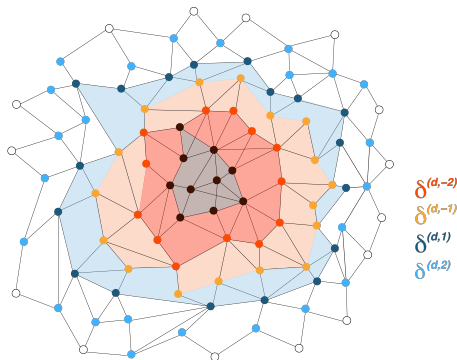


- perform **SpMV** with local subdomain and 2nd-level ghost



## Domain Decomposition Preconditioner for CA-Krylov

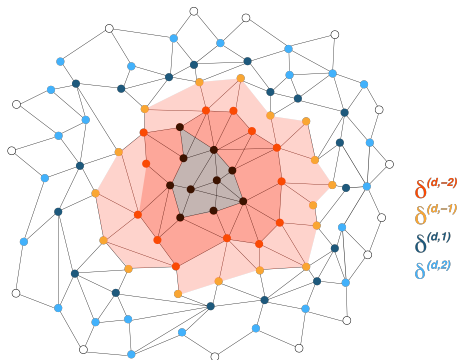
After SpMV at 1st step of MPK,



- effects of interior precondition propagates into 2nd-level underlap

## Domain Decomposition Preconditioner for CA-Krylov

After SpMV at 2nd step of MPK,



- effects of interior preconditioner pages into 1st-level ghost

## Matrix Powers Kernel Performance on a node

