HPX Applications and Performance Adaptation

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SAND2015-8999C

The HPX runtime system is a critical component of the DOE XPRESS (eXascale PRogramming Environment and System Software) project and other projects world-wide. We are exploring a set of innovations in execution models, programming models and methods, runtime and operating system software, dynamic and adaptive scheduling and resource management algorithms and mechanisms, and instrumentation and introspection techniques to achieve unprecedented efficiency, scalability, and programmability in the context of billion-way parallelism. A number of applications have been implemented to drive system development and quantitative evaluation of the HPX system implementation details and operational efficiencies and scalabilities.

Applications and Characteristic Behavior

- **LULESH:** (Livermore Unstructured Lagrangian Explicit Shock Hydrodynamics) For details see Deep Dive to right.
- **Mini-ghost:** A miniapp for exploring boundary exchange strategies using stencil computations in scientific parallel computing. Implemented by decomposing the spatial domain, inducing a "halo exchange" of process-owned boundary data.
- **N-Body Code:** An event driven constraint based execution model using the Barnes-Hut algorithm where the particles are grouped by a hierarchy of cube structures using a recursive algorithm. It uses an adaptive octree data structure to compute center of mass and force on each of the cubes with resultant O(N logN) computational complexity making use of LibGeoDecomp an auto-parallelizing library.
- **PIC:** 3D particle-in-cell (PIC) codes, such as GTC developed for studying turbulent transport in magnetic confinement fusion plasmas, which models interaction between fields and particles by solving the 5D gyro-averaged kinetic equation coupled to the Poisson equation, and PICSAR a miniapp with key functionalities of PIC accelerator codes including Maxwell solver using an arbitrary order finite-difference scheme (staggered/centered), a particle pusher using the Boris algorithm, and an energy conserving field gathering routine is energy conserving with high order particle shape factors.
- **miniTri:** A newly developed triangle enumeration-based data analytics miniapp. miniTri mimics the computation requirements of an important set of data science applications, not well represented by traditional graph search benchmarks such as Graph500. An asynchronous HPX-based approach enables our linear algebra-based implementation of miniTri to be significantly more memory efficient, allowing us to process much larger graphs
- **CMA:** (Climate Mini-App) For details see Deep Dive to right.
- **Kernels:** Studying various computational kernels, such as matrix transpose, and fast multipole algorithms, to explore features of HPX and compare to other approaches.

LULESH (Deep Dive)

LULESH--The Sedov blast wave problem in three dimensions is spherically-symmetric and the code solves the problem in a parallelepiped region. In the figure, symmetric boundary conditions are imposed on the colored faces such that the normal components of the velocities are always zero; free boundary conditions are imposed on the remaining boundaries.

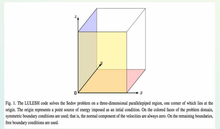


Fig. 1: The LULESH code solves the Sedov problem on a three-dimensional parallelepiped region, one corner of which lies at the origin. The origin represents a point source of energy imposed as an initial condition. On the colored faces of the problem domain, symmetric boundary conditions are used; that is, the normal component of the velocities are always zero. On the remaining boundaries, free boundary conditions are used.

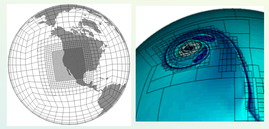
The LULESH algorithm is implemented as a hexahedral mesh-based code with two centerings. Element centering stores thermodynamic variables such as energy and pressure. Nodal centering stores kinematics values such as positions and velocities. The simulation is run via time integration using a Lagrange leapfrog algorithm. There are three main computational phases within each time step: advance node quantities, advance element quantities, and calculate time constraints. There are three communication patterns, each regular, static, and uniform: face adjacent, 26 neighbor, and 13 neighbor communications, illustrated below:

[1] LULESH is available from: https://codesign.llnl.gov/lulesh.php

[2] Karlin I, Bhatele A, Keasler J, Chamberlain BL, Cohen J, DeVito Z, et al. Exploring Traditional and Emerging Parallel Programming Models using a Proxy Application. In: Proc. of the 27-th IEEE International Parallel and Distributed Processing Symposium (IPDPS); 2013.

CMA (Deep Dive)

The climate mini app (CMA) models the performance profile of an atmospheric "dynamic core" (dycore) for non-hydrostatic flows. The codes use a conservative finite-volume discretization on an adaptively-refined cubed-sphere grid. An implicit-explicit (IMEX) time integrator combines a vertical implicit operator (which is FLOP-bound) with a horizontal explicit operator (which is bandwidth-bound).

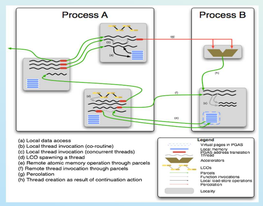


Figures: Left image shows an example of an adaptively refined cubed-sphere grid used in climate codes. Right image is vorticity dynamics for a climate test problem with AMR.

The mini app is implemented using the Chombo adaptive mesh refinement (AMR) framework, and has both an MPI+OMP and HPX backend. The mini-app is being used to explore performance on multi-core architectures (e.g. Xeon Phi) and to explore the benefits of using HPX for finite-volume AMR codes to combat dynamic load imbalance.

Background Info

ParalleX Execution Model

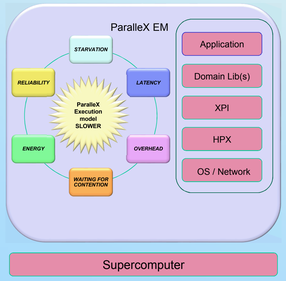


- **Lightweight multi-threading**
  - Divides work into smaller tasks
  - Increases concurrency
- **Message-driven computation**
  - Move work to data
  - Keeps work local, stops blocking
- **Constraint-based synchronization**
  - Declarative criteria for work
  - Event driven
  - Eliminates global barriers
- **Data-directed execution**
  - Merger of flow control and data structure
- **Shared name space**
  - Global address space
  - Simplifies random gathers

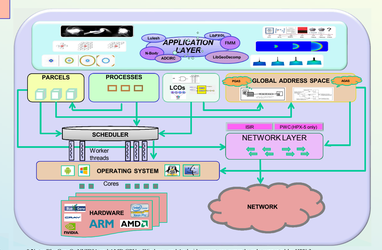
HPX-5 is the High Performance ParalleX runtime library from Indiana University. The HPX-5 interface and C99 library implementation is guided by the ParalleX execution model (http://hpx.crest.iu.edu).
HPX-3 is the C++11/14 implementation of ParalleX execution model from Louisiana State University (http://stellar-group.org/libraries/hpx/).

High Performance ParalleX (HPX)

* The HPX runtime system reifies the ParalleX execution model to support large-scale irregular applications:
  - Localities
  - Active Global Address Space (AGAS)
  - ParalleX Processes
  - Complexes (ParalleX Threads and Thread Management)
  - Parcel Transport and Parcel Management
  - Local Control Objects (LCOs)
* Sits between the application and OS
* Portable interface: C++11/14 (*HPX-3 only*), XPI
  - Comprehensive suite of parallel C++ algorithms (*HPX-3 only*)
* Automatic distributed garbage collection in AGAS (*HPX-3 only*)
* Flexible set of execution and scheduling policies
* Performance counter framework (*HPX-3 only*)



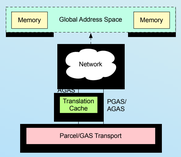
HPX Architecture



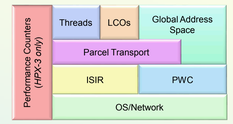
* **Note:** BlueGene/Q, NVIDIA and AMD GPUs, Windows, and Android support are currently only supported by HPX-3

Memory models and Transport

HPX supports three memory models: Symmetric Multi-Processing (SMP, no global memory), a Partitioned Global Address Space (PGAS), an Active Global Address Space (AGAS).



HPX avoids the use of locks and/or barriers in parallel computation through the use of LCOs, which are lightweight synchronization objects used by threads as a control mechanism. Reads and writes on LCOs are globally atomic and require no other synchronization mechanism.

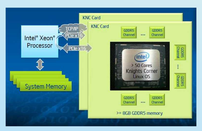


* Optimized transports built on top of Photon (*HPX-5 Only*) and other communication libraries
* Two-sided Isend/Irecv transport (ISIR)
Pre-posts irecvs to reduce probe overhead
* One-sided Put-With-Command/Completion (PWC)
Local/remote notifications for RDMA operations
RDMA communication optimizations using Photon: turn large puts into gets, buffer coalescing

- Thomas Sterling, Daniel Kogler, Matthew Anderson, and Maciej Brodowicz. SLOWER: A performance model for Exascale computing. Supercomputing Frontiers and Innovations, 1:42–57, September 2014.
- Hartmut Kaiser, Thomas Heller, Bryce Adelstein-Lelbach, Adrian Serio, Dietmar Fey, HPX – A Task Based Programming Model in a Global Address Space, PGAS 2014: The 8th International Conference on Partitioned Global Address Space Programming Models (2014).

Results Showing Benefits of HPX

NERSC's Edison, a Cray XC30 using the Aries interconnect and Intel Xeon processors with a peak performance of more than 2 petaflops.

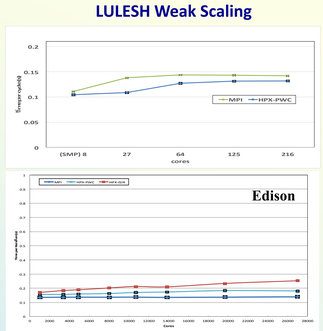


NERSC's Babbage machine uses the Intel Xeon Phi™ coprocessor (codenamed "Knights Corner"), which combines many Intel CPU cores onto a single chip. Knights Corner is available in multiple configurations, delivering up to 61 cores, 244 threads, and 1.2 teraFLOPS of performance.

GTCX Communication Reduction

Reduction of communication in GTCX (GTC with HPX added) compared to original GTC (upper image) is shown.

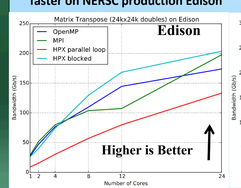
LULESH Weak Scaling



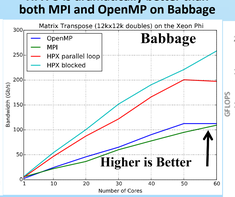
Top: HPX-5 weak-scaling LULESH performance on 256 core cluster. **Bottom**: HPX-5 weak scaling LULESH performance on Edison up to 14000 cores. ISIR and PWC are HPX-5 network back-ends. Lower values are better and we demonstrate developing 27k+ core scaling.

Matrix Transpose Kernel

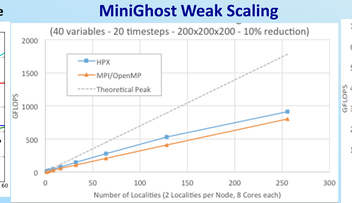
HPX-3 blocked implementation faster on NERSC production Edison



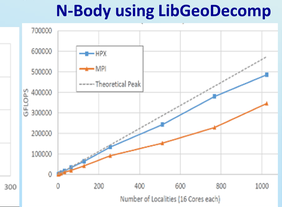
HPX-3 is dramatically better than both MPI and OpenMP on Babbage



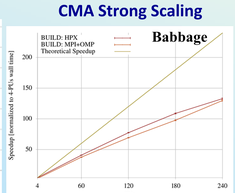
MiniGhost Weak Scaling



N-Body using LibGeoDecomp



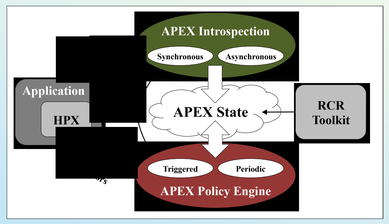
CMA Strong Scaling

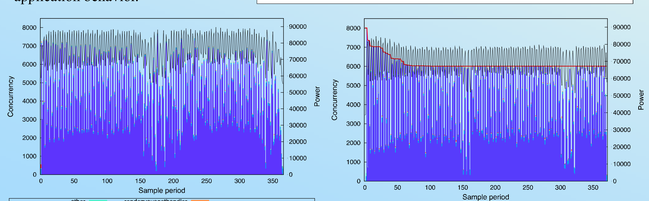


Performance Adaptation, Legacy Applications, and Summary

APEX : Performance Adaptation

HPX runtime implementations are integrated with APEX (Autonomic Performance Environment for Exascale), a feedback/control library for performance measurement and runtime adaptation. *APEX Introspection* observes the application, runtime, OS and hardware to maintain the *APEX state*, while the Policy Engine enforces policy rules to adapt, constrain or otherwise modify application behavior.

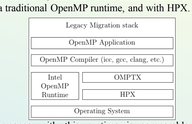




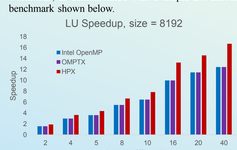
Figures: Concurrency views of LULESH (HPX-5), as observed and adapted by APEX. The left image is an unmodified execution, while the right image is a runtime power-capped (220W per-node) execution, with equal execution times. (8000 subdomains, 643 elements per subdomain on 8016 cores of Edison, 334 nodes, 24 cores per node). Concurrency throttling by APEX resulted in 12.3% energy savings with no performance change.

Legacy Application support

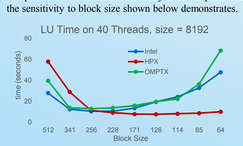
Legacy OpenMP Application support is implemented using an HPX implementation of the intel OpenMP runtime. As illustrated below, this allows the same binaries to be executed with a traditional OpenMP runtime, and with HPX.



Performance with this runtime is comparable with most workloads, as demonstrated with a simple LU decomposition benchmark shown below.



While the existing OpenMP applications have performance differences, the performance of an application written and compiled with HPX has substantially different performance, as the sensitivity to block size shown below demonstrates.



Summary

Exascale programming models and runtime systems are at a critical juncture in development.

Systems based on light-weight tasks and data dependence are an excellent method for extracting parallelism and achieving performance.

HPX is emerging as an important new path with support from US Department of Energy, the National Science Foundation, the Bavarian Research Foundation, and the European Horizon 2020 Programme.

Application performance of HPX codes on the very recent architectures including the current and prototypical next-generation Cray-Intel machines is very good.

For some of the applications the performance using HPX is significantly better than using standard MPI + OpenMP implementations.

Performance adaptation using APEX provides significant energy savings with no performance change.

We have shown that legacy application using OpenMP can run under the HPX runtime system effectively.

Funding by DOE Office of Science through grants: DE-SC0008714, DE-SC0008809, DE-AC04-94AL85000, DE-SC0008596, DE-SC0008638, and DE-AC02-05CH11231, by National Science Foundation through grants: CNS-1117470, AST-1240655, CCF-1160602, IIS-1447831, and ACI-1339782, and by Friedrich-Alexander-University Erlangen-Nuremberg through grant: H2020-EU.1.2.2. 671603.