

# Secure Distributed Set Membership through Secret Sharing

## Sandia National Laboratories

T. M. Kroeger, D. J. Zage and C. A. Phillips  
Livermore, California & Albuquerque, New Mexico

## University of New Mexico

J. Saia  
Albuquerque, New Mexico

## Tufts University

T. R. Benson  
Medford, Massachusetts

## Problem

### Data Security and Availability Conflict

Data security and availability for operational use are frequently seen as fundamentally opposing forces.

#### Encryption is fragile:

Tools like homomorphic encryption are a start but most encryption algorithms are at best assumed to be secure, and often in reality just delayed release.

#### Secret Sharing Provides Provably Secure Systems

Archives that distribute data with secret sharing can provide information theoretic data protections and a resilience to:

1. malicious insiders,
2. compromised systems, and
3. untrusted components.

### We are developing ways to functionally use secret shares without reassembly.

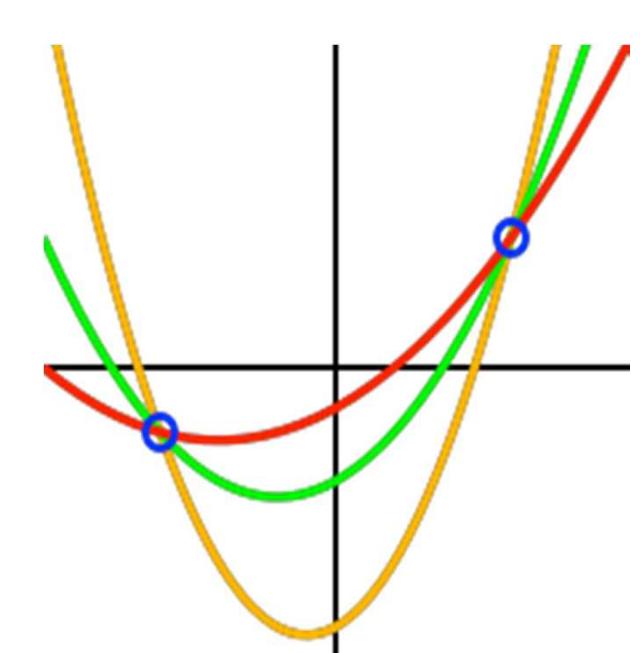
## Approach

### Secure Foundation: Shamir Secret Splitting

Shamir's Secret Sharing (SSS) provides a data protection that goes far beyond just splitting the data into parts. It is **information theoretically secure** and uses points on a polynomial curve to securely encode sensitive data.

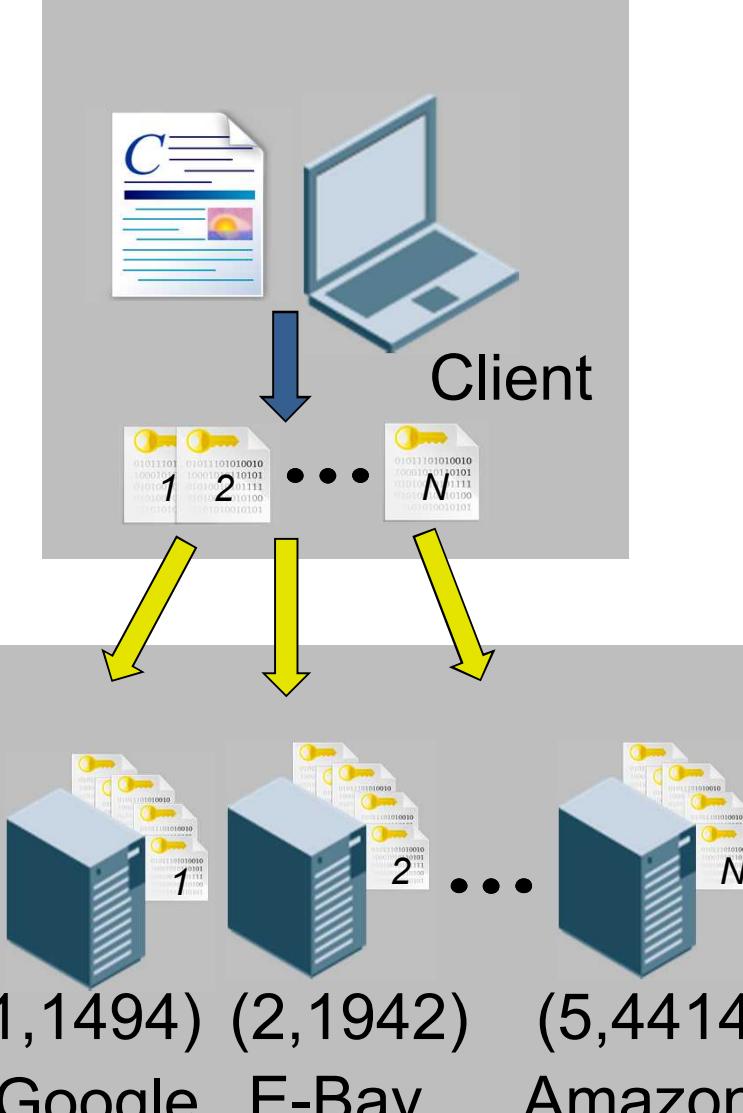
#### Example:

Any 3 of 5 Secret  $S = 1234$



An infinite number of polynomials of degree 2 exist through 2 points.

#### Distributed Secure Archive



Any three points enables a user to solve for  $S$ . With one or two points you know nothing more than when you had none;  $S$  is one of infinite possible Y-intercepts.

## Results

### Serial Interpolation Filter (SIF)

By using Lagrange Interpolation serially across the archive, we can query our data **without** exposing the data.

Lagrange Interpolation allows for the recovery of the original function via the stored points:

$$f(x) = \sum_{i=1}^n L_i(x)f(i)$$

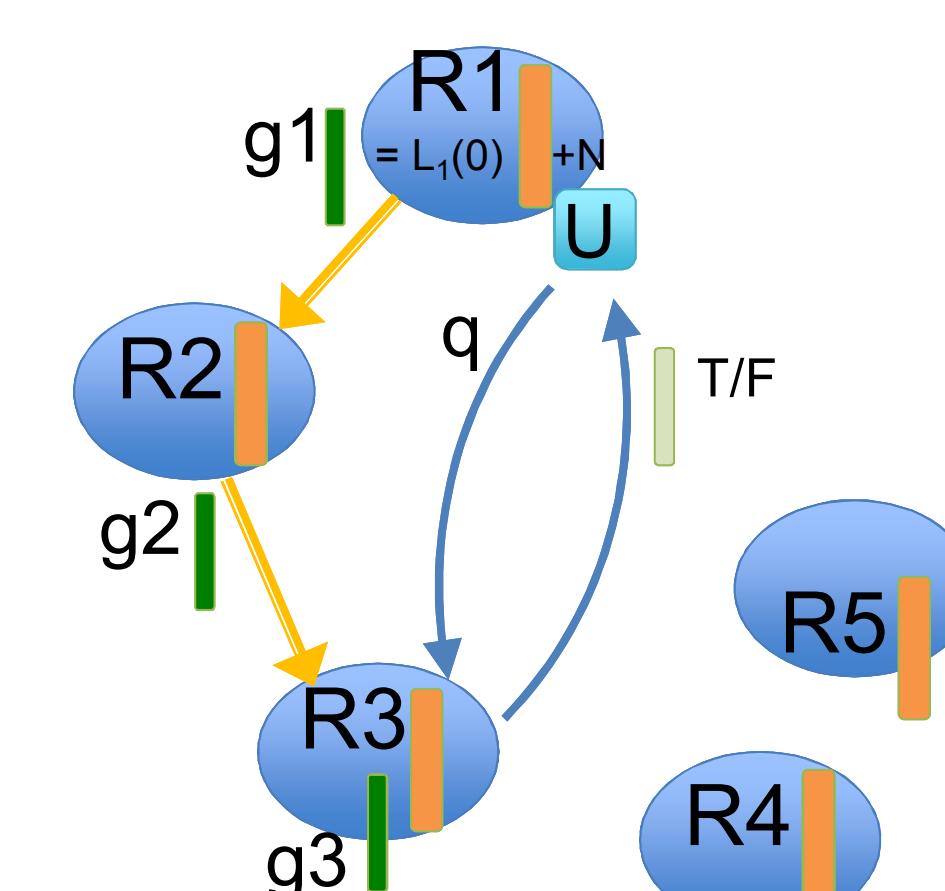
where  $L_i(x)$  is the Lagrange basis polynomial and  $f(i)$  is the function evaluated at  $i$ . Reconstructed,  $f(0) = S$ .

#### Example:

5 companies are willing to share their list of sensitive security data (e.g., list of known bad address). They encode each entry with a 2<sup>nd</sup> order polynomial. Now we can use a SIF to determine if a given address is in the list.

User  $U$ , who works at company  $R1$ , does the following to test if an address  $Z$  is on the list.

- 1)  $U$  creates random private random number called a nonce  $N$ .
- 2)  $U$  &  $R1$  calculate  $g1 = L_1(x)f(1) + N$  and send  $g1$  to  $R2$
- 3)  $R2$  calculates  $g2 = L_2(x)f(2) + g1$  sends  $g2$  to  $R3$
- 4)  $R3$  calculates  $g3 = L_3(x)f(3) + g2$
- 5)  $U$  sends  $q = Z + N$  to  $R3$
- 6)  $R3$  For all entries where  $g3 == q$  send TRUE to  $U$  else send FALSE to  $U$



## Significance

With data breaches plaguing both government and corporate America, the need for better data security architectures is painfully obvious.

#### Our approach:

- provides information theoretic protections,
- avoids difficult key management, and
- enables interacting with the data without reassembly, preserving the protections of SSS.