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QUANTUM INFORMATION PROCESSING

An Overview

Uzoma Onunkwo

Advanced Information and Network Systems Engineering Department
Sandia National Laboratories
Albuquerque, NM

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- ① DEMANDS OF WIRELESS COMMUNICATION
- ② INFORMATION PROCESSING
 - Digital Information
 - Quantum Information
- ③ QUANTUM ALGORITHMS
- ④ CRYPTOGRAPHY
 - Confidentiality
 - Authentication
 - Signature
- ⑤ QUANTUM CRYPTOGRAPHY
 - *Quantumness*: Relevant Laws
 - Achieving Perfect Secrecy
 - Quantum Key Distribution (QKD) Purpose
 - BB84 - A Quantum Cryptography Protocol
- ⑥ FINAL REMARKS
- ⑦ REFERENCES



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- ① Higher effective data rates: 1G (*circa 1981*) → 2G (*circa 1991*) → 3G (*circa 2001*) → 4G (*circa 2012*) → 5G (*maybe 2020*).



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- 2 Higher capacity in terms of number of simultaneous users.



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- ② Higher capacity in terms of number of simultaneous users.
- ③ Secure communications for privacy.



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COMPUTATION 101

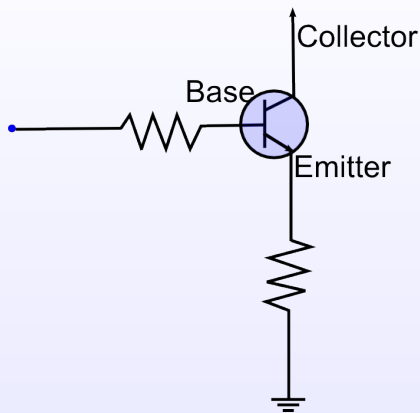
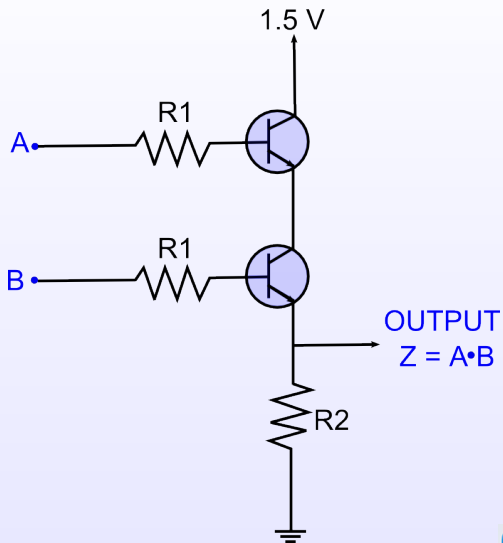


FIGURE 1: Transistor circuit: *increasing voltage at base, exponentially increases current flow from collector to emitter; the converse is also true...we have a switch!*



COMPUTATION 101



COMPUTATION 101

BASIC UNIVERSAL GATES (generally, irreversible)



AND gate			OR gate			NOT gate	
A	B	Q	A	B	Q	A	Q
0	0	0	0	0	0	0	1
0	1	0	0	1	1	1	0
1	0	0	1	0	1		
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1-BIT ADDER

input (two 1-bits)		output (2-bit)		decimal value
x	y	b_1	b_2	
0	0	0	0	0
0	1	0	1	1
1	0	0	1	1
1	1	1	0	2



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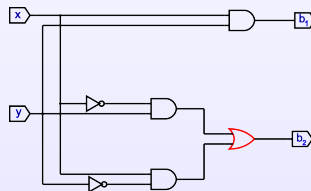
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QUANTUM COMPUTING 101: *A compare and contrast approach*

	Classical	Quantum
Basis of information	Bits: 0 or 1 2-level system	Qubits: $ \psi\rangle = \begin{pmatrix} a \\ b \end{pmatrix}$ infinitely possible vectors or states, e.g., $ 0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$, $ \phi\rangle = \begin{pmatrix} 0.8 \\ 0.6 \end{pmatrix}$
Gates	AND, OR, NOT generally irreversible	Bit-flip = $\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ Idle rotation = $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$ reversible operations
Can transport (qu)bits	Yes	Yes
Can clone (copy) (qu)bits	Yes	No
Effect of measurement	Nothing	Destroys original qubit



POTENTIAL REALIZATIONS OF QUBITS

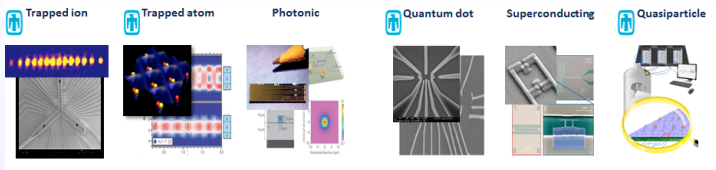


FIGURE 2: Source: Technical Overview presentation on AQUARIUS Grand Challenge

- Physical realization of qubits; *think of vacuum tubes and transistors for bits*
 - Quantum dots (*non-mobile*)
 - Ion traps (*mobile*)
 - Neutral atom laser (*non-mobile?*)
 - Photons (*mobile*)
 - Superconducting flux (*non-mobile*)

MODELS OF QUANTUM COMPUTING

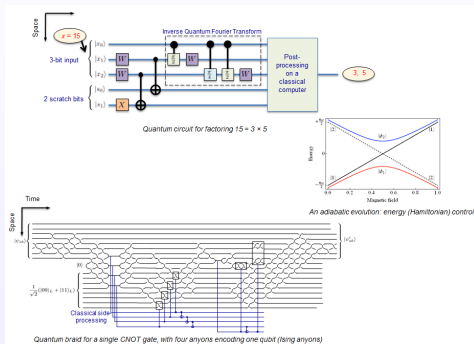


FIGURE 3: Source: Technical Overview presentation on AQUARIUS Grand Challenge

- Models of quantum computing
 - Circuit or network model
 - Adiabatic quantum computing model
 - One-way (Cluster state) computing model



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QUANTUM ALGORITHMS

- *Grover's Algorithm*: Existence in an Unsorted Database.
- *Shor's Algorithm*: Factorization of Composite Numbers.
 - Extremely hard **classically**. Best classical solution takes exponential time in the number of bits to run.
 - Relatively easy to solve **quantum-wise** using *Shor's method*.

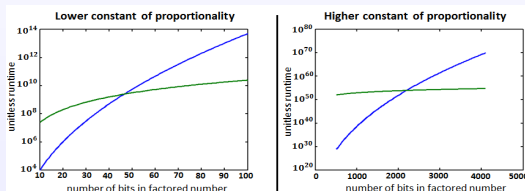


FIGURE 4:

- *Hallgren's Algorithm*: Solution to Pell's equation ($x^2 - dy^2 = 1$, $\{x, y, d\} \in \mathbb{Z}^+$ and d is a non-square positive integer).



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INFORMATION CONFIDENTIALITY

Purpose is to obfuscate information.

- Universally used:
 - ① bank transactions,
 - ② investment account management,
 - ③ sensitive e-mails, ...
- Desired property:
 - ① *impossible* to decipher (decrypt) by an unwanted entity,
 - ② should stand the test of time, and
 - ③ be easy enough to be *implementable* today.
- Decent solutions: block cipher vs. stream cipher, secret-key vs. public-key, ...



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Can encrypt as

cmbicmbi jt nz qbttxpse.

Is this good?



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Is this good? No, it is **horrible!** Susceptible to *Chosen-Plaintext Attack*



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AUTHENTICATION: *They are who we thought they were!*



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- Validates the source of transmitted data.
- This is implemented over a classical channel, even in *quantum crptography*.
- Three main techniques exist:
 - ① Use of *message authentication code* (MAC) - amenable to secret-key ciphers;
 - ② Use of public key infrastructure (PKI) - amenable to public-key ciphers;
 - ③ Use of *universal families of hash functions* - amenable to quantum cryptography.



SIGNATURE: THIS IS JAMES BOND AND I APPROVE THIS MESSAGE

- It is okay to eavesdrop, but it is not okay to forge my signature.
- Achievable with a form of encryption, *private-key encryption*.
- **Alice** signs a message. **Bob** receives the message and decrypts with *Alice's public-key*; this confirms Alice's signature.
- **Eve** can listen in the middle and even extract the message, but she cannot forge the signature



PERFECT SECRECY IN COMMUNICATIONS

- Can only be achieved if key size is greater than or equal to message size.
- Call $P = \{p : p \text{ is in the set of all plaintext}\}$ and $C = \{c : c \text{ is in the set of corresponding ciphertext}\}$, then if

$$\Pr[P = p | C = c] = \Pr[P = p]$$

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- Alice and Bob are legit users, while Eve is the eavesdropper. Eve should be able to listen/view cipher message all she wants, but not decipher it.



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- Perfect secrecy of information **cannot** be achieved with key lengths less than message length.
 - Today's ciphers use key lengths much smaller than message lengths!
 - AES-256 uses a key length of 256 bits, message block size of 128 bits, but the same key for a long message.
 - RSA (cipher for SSH) uses keys of size 1024, 2048, and 4096 bits, but much longer messages.



SO, AREN'T CURRENT CRYPTOGRAPHIC TECHNIQUES SUFFICIENT?



SO, AREN'T CURRENT CRYPTOGRAPHIC TECHNIQUES SUFFICIENT?

- Current classical cryptographic systems are mostly *cat-and-mouse* solutions...propose solutions based on current attacks.
- Strong presumptions made on solvability of currently hard problems, like *modulo-factoring* for the famous RSA-cipher system.



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RELEVANT LAWS OF QUANTUM PHYSICS

- 1 It is impossible to clone a qubit, the equivalent of classical bits. This limitation is mathematically proven and does NOT rely on technological limitation.



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- ① It is impossible to clone a qubit, the equivalent of classical bits. This limitation is mathematically proven and does NOT rely on technological limitation.
- ② Observing the state of a qubit destroys its previous state; hence, **we cannot distinguish between non-orthogonal states**:
 - Classically, we only have 1 and 0's...we can distinguish these.
 - Quantum-wise, we have $|1\rangle, |0\rangle, |+\rangle, |-\rangle, \dots$ in fact infinite representative qubits!



CLONING QUBITS IS IMPOSSIBLE: *A No-go Theorem*

- One can make multiple copies of an *a priori* known qubit, but...
- One cannot make a gadget that makes copy of any arbitrary input qubit

PROOF.

If such a gadget existed, then it should do the following: $|\psi\rangle \mapsto |\psi\rangle \otimes |\psi\rangle \equiv |\psi\psi\rangle$. Let such a gadget be represented by the linear operator, \mathcal{C} . The requirement for linearity of the operator comes from Schrödinger's equation. Then,

$$\mathcal{C} |\psi\rangle = |\psi\psi\rangle$$

By the same token, we expect

$$\mathcal{C} |0\rangle = |00\rangle \quad (1)$$

$$\mathcal{C} |1\rangle = |11\rangle \quad (2)$$

$$\mathcal{C} (|0\rangle + |1\rangle) = (|00\rangle + |11\rangle) \quad (3)$$

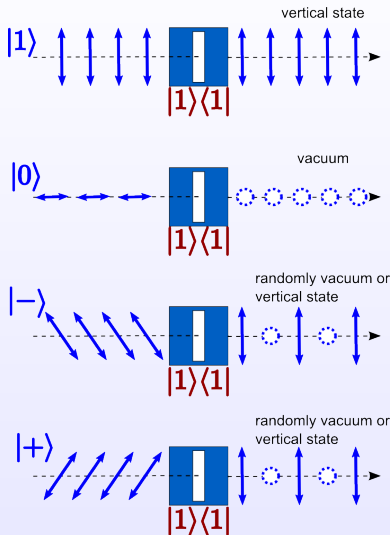
The last line comes from the linearity property of the operator \mathcal{C} , which contradicts the expected result of

$$\begin{aligned} \mathcal{C} (|0\rangle + |1\rangle) &= (|0\rangle + |1\rangle) \otimes (|0\rangle + |1\rangle) \\ &\neq (|00\rangle + |11\rangle) \end{aligned}$$

Thus, no such gadget can exist; we call this the *No-cloning Theorem*. □



MORE DETAILS ON QUBIT OBSERVATION



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FIGURE 5: photon qubits being measured by polarizer

PERFECT SECRECY REVISITED

- One-time pad with key length at least as long as message:
assume plaintext is 0110 and random key is 1010, then
ciphertext can be 1100. Can't decipher plaintext from this!



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- Cannot reuse cipher key or Eve can decipher message by performing XOR on consecutive message.
- Requires true random keys. Pseudo-random numbers prevalent on today's computers is not sufficient; they are deterministic unfortunately. TRNG generators exist:
 - ① Hotbits - rely on radioactive decay, but slow rates < 1Kbits/sec.
 - ② Protego's SG100 EVO-USB - rely on resistive circuit elements and yields about 16 Mbits/sec per module.
 - ③ Quantum measurement of qubits.



QUANTUM CRYPTOGRAPHY IS REALLY QUANTUM KEY DISTRIBUTION

- We are not encrypting *qubits*, we are encrypting *bits* in plaintext.
- One-time pad needs constantly changing keys equal to message length or more.
- QKD achieves safe secret sharing using (a) quantum channel and (b) classically authenticated channel.



A QUANTUM CRYPTOGRAPHY PROTOCOL - BB84

- Alice will send messages to Bob. Eve is the eavesdropper.
- Alice and Bob authenticate to each other, using classical channel but maybe with universal hash function technique.
- Alice chooses four *non-orthogonal* qubit states.
- Bob measures these states, with some errors.
- Alice and Bob negotiate error and decide whether to discard or to correct and use distilled key.



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- Bob measures these states, with some errors.
- Alice and Bob negotiate error and decide whether to discard or to correct and use distilled key.
- If Eve observes the channel, she will likely cause errors that both Alice and Bob will detect.
- In addition, Eve cannot store the qubits for future use in deciphering messages (*No-cloning principle*).



BB84 - PICTORAL DEPICTION

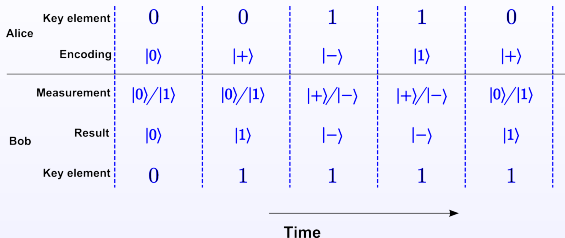
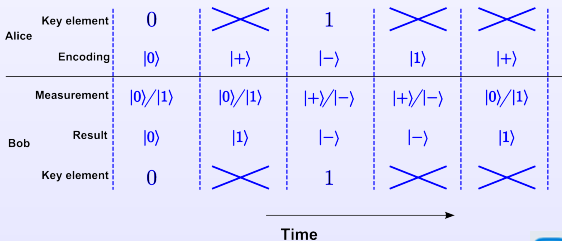


FIGURE 6:



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FIGURE 7: Redrawn from G.V. Assche's textbook [2].

FINAL REMARKS

- Quantum science shows promises for alternative way of information processing.
- QKD aides secret sharing of cryptographic keys based on information-theoretic limits not computational constraint. When combined with a one-time padding cipher, one can attain perfect secrecy.
- Commercial companies providing quantum key distribution (QKD) systems are:
 - ① **id Quantique** in Geneva, Switzerland,
 - ② **MagiQ Technologies** in New York, USA, and
 - ③ **QuintessenceLabs** in Australia.



QUESTIONS

Thank you for being here!
Q & A



- [1] M. A. Nielsen and I. L. Chuang, *Quantum Computation and Quantum Information*, Cambridge University Press, 2000.
- [2] G. V. Assche, *Quantum cryptography and secret-key distillation*. Cambridge University Press, 2006.
- [3] U. M. Maurer, "Authentication theory and hypothesis testing," *Information Theory, IEEE Transactions on* vol. 46, no. 4, pp. 1350 - 1356, 2000.

