

War Stories : Graph Algorithms in GPUs

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FASTMath Meeting



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Increasingly Complex Heterogeneous Future;

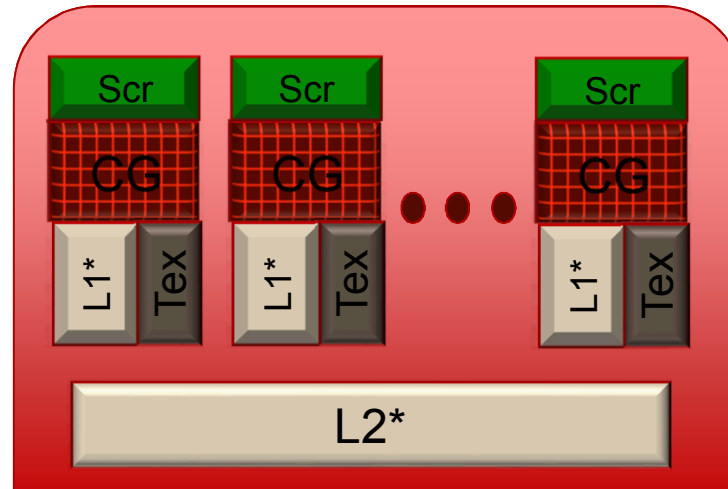
¿ Future Proof Performance Portable Code?

Memory Spaces

- Bulk non-volatile (Flash?)
- Standard DDR (DDR4)
- Fast memory (HBM/HMC)
- (Segmented) scratch-pad on die

Execution Spaces

- Throughput cores (GPU)
- Latency optimized cores (CPU)
- Processing in memory

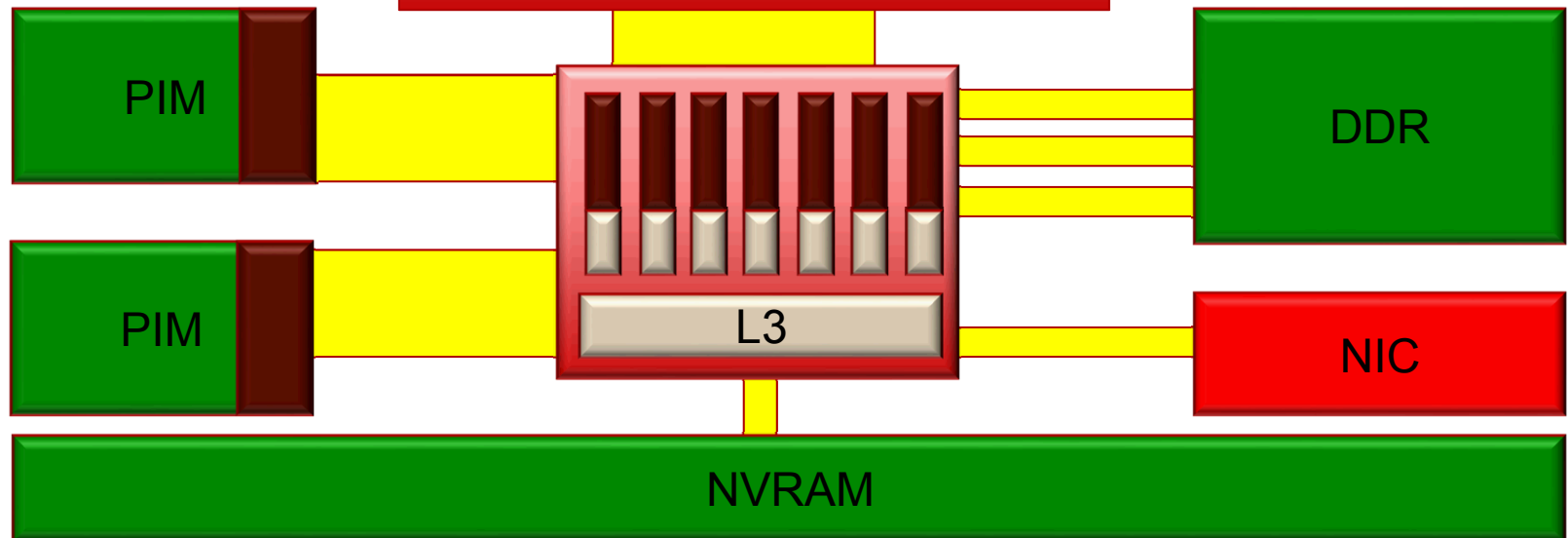


Special Hardware

- Non caching loads
- Read only cache
- Atomics

Programming models

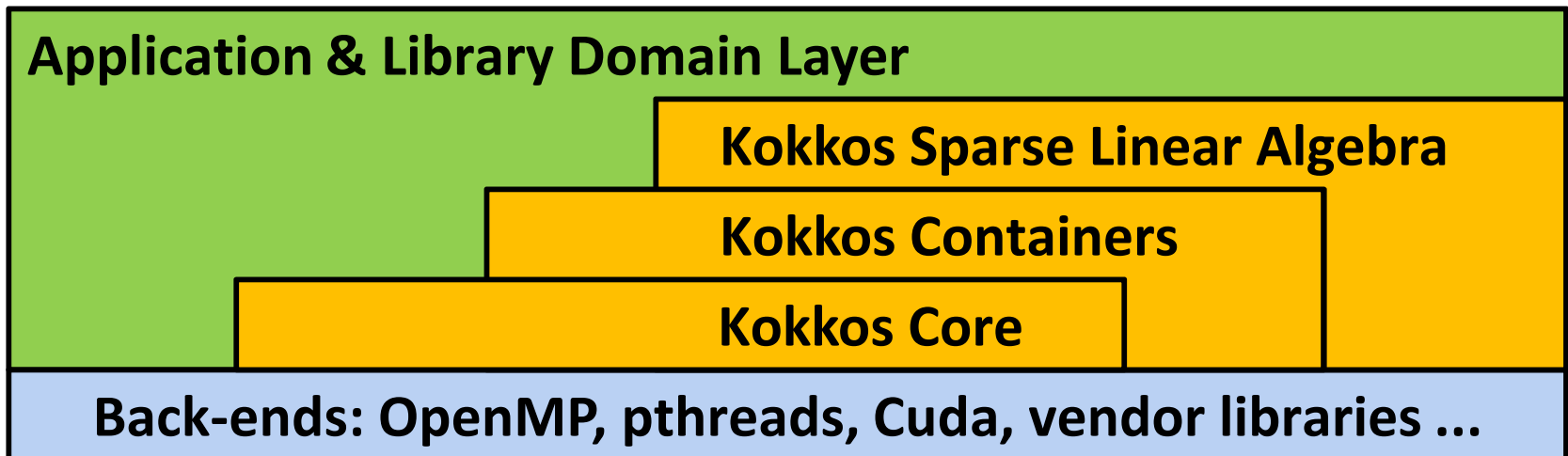
- GPU: CUDA-ish
- CPU: OpenMP
- PIM: ??



- **What is Kokkos (Slides from Kokkos Developers: Carter Edwards, Christian Trott, Dan Sunderland)**
 - Layered collection of C++ libraries
 - Thread parallel programming model that managed data access patterns
- Graph Algorithms with OpenMP
- Graph Algorithms with Kokkos
- Conclusion

Kokkos: A Layered Collection of Libraries

- Standard C++, Not a language extension
 - In *spirit* of Intel's TBB, NVIDIA's Thrust & CUSP, MS C++AMP, ...
 - *Not* a language extension: OpenMP, OpenACC, OpenCL, CUDA
- Uses C++ template meta-programming
 - Currently rely upon C++1998 standard (everywhere except IBM's xLC)
 - Prefer to require C++2011 for lambda syntax
 - Need CUDA with C++2011 language compliance



Kokkos' Layered Libraries

■ Core

- Multidimensional arrays and subarrays in memory spaces
- `parallel_for`, `parallel_reduce`, `parallel_scan` on execution spaces
- Atomic operations: compare-and-swap, add, bitwise-or, bitwise-and

■ Containers

- *UnorderedMap* – fast lookup and thread scalable insert / delete
- *Vector* – subset of `std::vector` functionality to ease porting
- *Compress Row Storage (CRS) graph*
- *Host mirrored & synchronized device resident arrays*

■ Sparse Linear Algebra

- *Sparse matrices and linear algebra operations*
- *Wrappers for vendors' libraries*
- *Portability layer for Trilinos manycore solvers*

Performance Portability Challenge: Require Device-Dependent Memory Access Patterns

- CPUs (and Xeon Phi)
 - Core-data affinity: consistent NUMA access (first touch)
 - Hyperthreads' cooperative use of L1 cache
 - Alignment for cache-lines and vector units
- GPUs
 - Thread-data affinity: coalesced access with cache-line alignment
 - Temporal locality and special hardware (texture cache)
- ¿ “Array of Structures” vs. “Structure of Arrays” ?
 - This is, and has been, the *wrong* question

Right question: Abstractions for Performance Portability ?

Kokkos Core: Fundamental Abstractions

- Devices have Execution Space and Memory Spaces
 - Execution spaces: Subset of CPU cores, GPU, ...
 - Memory spaces: host memory, host pinned memory, GPU global memory, GPU shared memory, GPU UVM memory, ...
 - Dispatch *computation* to *execution space* accessing data in *memory spaces*
- Multidimensional Arrays, *with a twist*
 - Map multi-index (i,j,k,...) \leftrightarrow memory location *in a memory space*
 - Map is derived from an array *layout*
 - Choose layout for device-specific memory access pattern
 - Make layout changes transparent to the user code;
 - IF the user code honors the simple API: $a(i,j,k,...)$

Separates user's index space from memory layout

Kokkos Core: Multidimensional Array Layout and Access Attributes

- Override device's default array layout

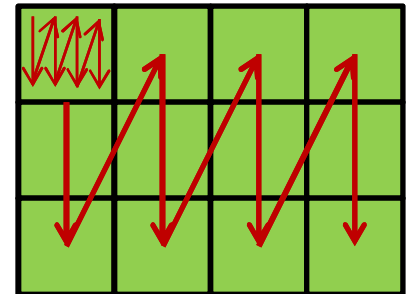
```
class View<double**[3][8], Layout , Device> a("a",N,M);
```

- E.g., force row-major or column-major
- Multi-index access is unchanged in user code
- *Layout* is an extension point for blocking, tiling, etc.

- Example: Tiled layout

```
class View<double** , TileLeft<8,8> , Device> b("b",N,M);
```

- Layout changes are transparent to user code
- IF the user code honors the `a(i,j,k,...)` API



- Data access attributes – user's intent

```
class View<const double**[3][8], Device, RandomRead> x = a ;
```

- Constant + RandomRead + GPU → read through GPU texture cache
- Transparent to user code

‘NW’ units of data parallel work

- `parallel_for(NW , functor)`
 - Call `functor(iw)` with $iw \in [0, NW)$ and $\#thread \leq NW$
- `parallel_reduce(NW , functor)`
 - Call `functor(iw , value)` which contributes to reduction ‘value’
 - Inter-thread reduction via `functor.init(value)` & `functor.join(value,input)`
 - Kokkos manages inter-thread reduction algorithms and scratch space
- `parallel_scan(NW , functor)`
 - Call `functor(iw , value , final_flag)` multiple times (possibly)
 - if `final_flag == true` then ‘value’ is the prefix sum for ‘iw’
 - Inter-thread reduction via `functor.init(value)` & `functor.join(value,input)`
 - Kokkos manages inter-thread reduction algorithms and scratch space

Kokkos Core: Dispatch Data Parallel Functors

League of Thread Teams (grid of thread blocks)

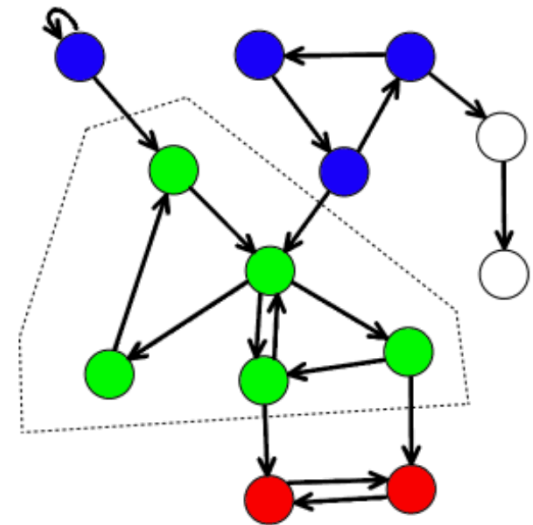
- `parallel_for({ #teams , #threads/team } , functor)`
 - Call functor(*teaminfo*)
 - *teaminfo* = { #teams, team-id, #threads/team, thread-in-team-id }
- `parallel_reduce({ #teams , #threads/team } , functor)`
 - Call functor(*teaminfo* , value)
- `parallel_scan({ #teams , #threads/team } , functor)`
 - Call functor(*teaminfo* , value , final_flag)
- A Thread Team has
 - Concurrent execution with intra-team collectives (barrier, reduce, scan)
 - Team-shared scratch memory
 - Exclusive use of CPU and Xeon Phi cores while executing

Outline

- What is Kokkos
- **Graph Algorithms with OpenMP**
- Graph Algorithms with Kokkos
- Conclusion

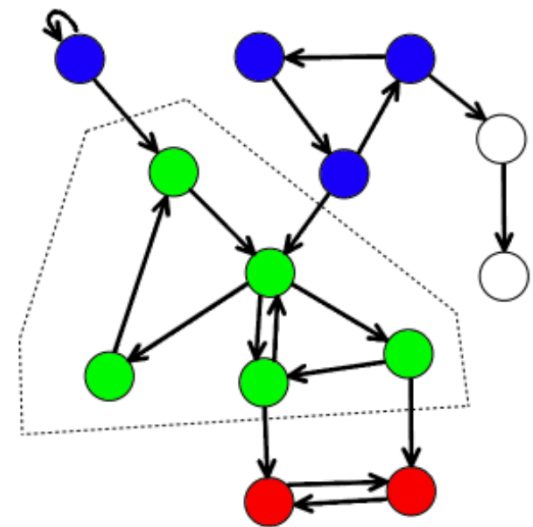
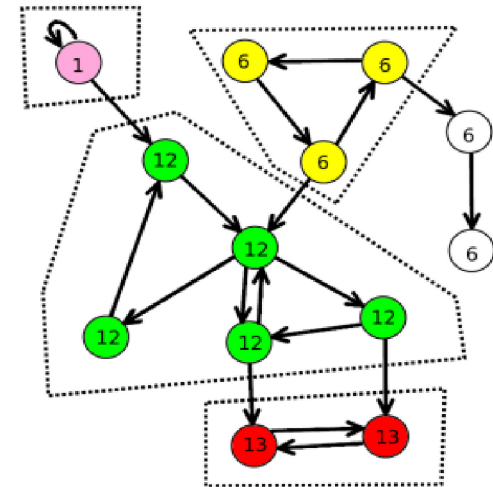
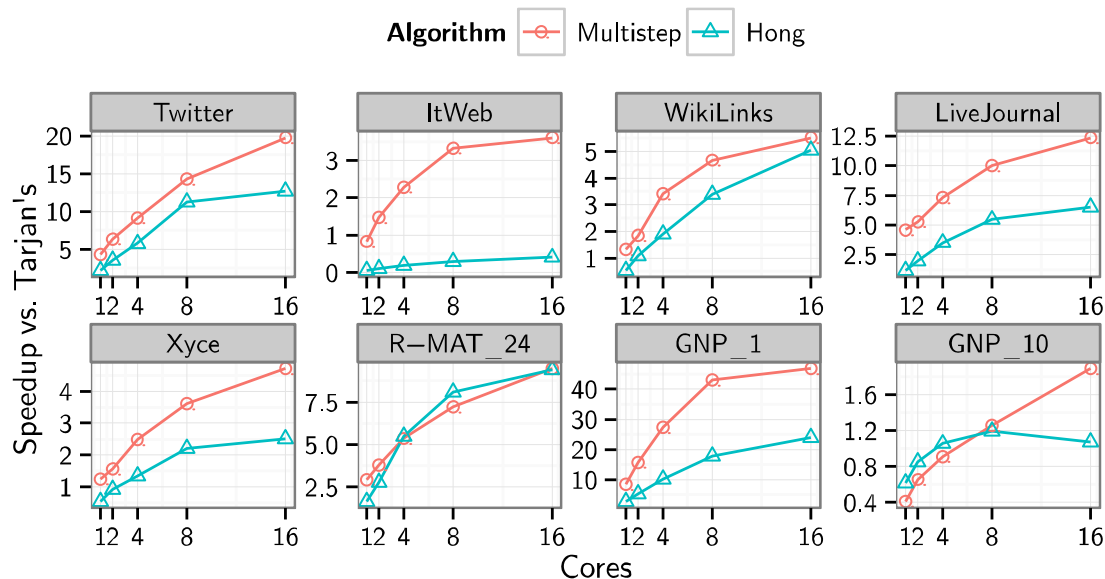
Computing Strongly Connected Components

- **Problem:** Given a graph find all the strongly connected components in the graph
- Multistep method:
 - Multithreaded with OpenMP
 - Optimized for the best CPU performance, state-of-the-art code.
 - Scales to millions of vertices and billions of edges
 - Data-Parallel code, minimal synchronization
 - Good as a baseline for porting to Kokkos
 - Scales to 16-32 threads
- FASTMath session in PP14 and IPDPS 14.



Multistep Method with OpenMP

- Different Steps of the Algorithm uses different types of parallelism
 - Per-vertex for-loop
 - Level synchronous BFS



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Thread Parallel vs Thread Scalable

- **Common construct in OpenMP programming**

- Allocate threadlocal data, do parallel work
- Non-Starter in the GPUs for large arrays

```
#pragma omp  
{  
    // allocate an array  
    parallel for do work  
}
```

- **Count, Allocate, Fill, paradigms**

- Non-Starter for graph algorithms

- **Need Algorithms that use tiny threadlocal data and synchronize with global memory**

- Tiny == 16 – 32 edges
- Expensive, too many synchronizations

- **Need Algorithms that use threadteams and shared memory (scratch space) between a team of threads.**

Thread Teams in GPUs

- Multiprocessor (up to about 15/GPU)
 - Multiple groups of stream processors (12×16)
 - Warps of threads all execute SIMT on single group of stream processors (32 threads/warp, two cycles per instruction)
 - Irregular computation (high degree verts, if/else, etc.) can result in most threads in warp doing NOOPs
- Kokkos Model:
 - Thread team - multiple warps on same multiprocessor, but all still SIMT for GPU
 - Thread league - multiple thread teams
 - Work statically partitioned to teams before parallel code is called

Challenges in ThreadScalable codes

- **Goal:** Fast Kokkos-based ThreadScalable algorithm for CPU/GPU/Phi
- Challenges:
 - No persistent thread-local storage
 - Minimize serial portions for GPU/Phi
 - Mitigate effect of high degree vertices, irregular graphs
 - Mitigate algorithmic differences of various architectures
- Solutions:
 - Very small static thread-owned arrays
 - No more Tarjan's, minimize possible per-thread work
 - Implement new algorithmic tweaks, for loadbalancing, to current methods
 - HOPE this doesn't happen !!!!

Handling Imbalance in GPU threads

- **Chunking:**

- Transform vertex queue into edge queue
- Each thread can explore only a few edges and chunks the rest of the edges for later stages

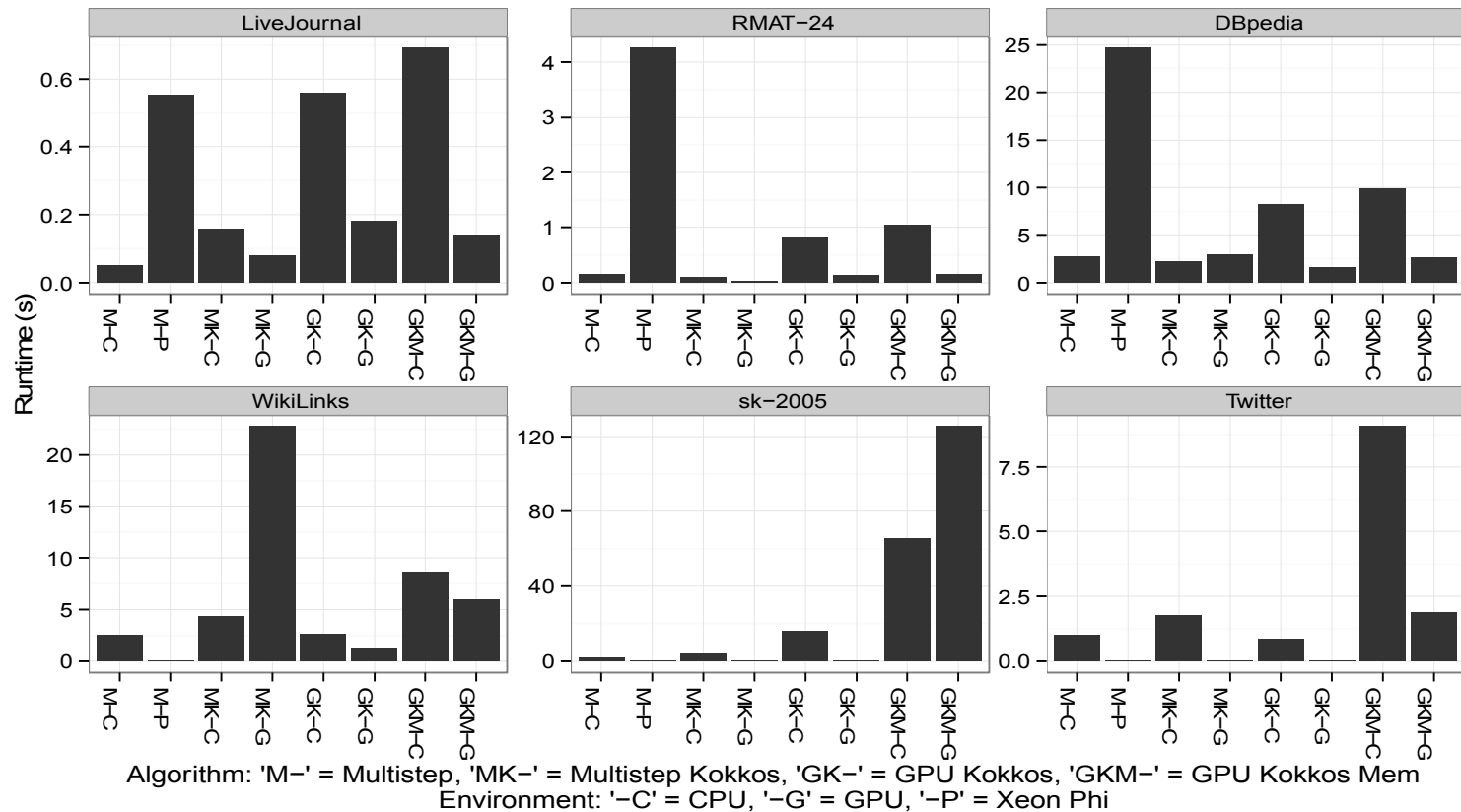
- **Delayed Exploration of High Degree Vertices:**

- When a single thread in a team encounters a high degree vertex, its exploration is delayed
- The vertex is placed in shared memory queue, only accessible by thread team (just a template parameter in Kokkos)
- Once team finishes original work (minus large degree vertices), the team works to explore all delayed vertices via inner loop parallelism
- On CPU, size of thread teams is usually 1, so this algorithm would default back to standard approach on that architecture

Results - Algorithms

- **Multistep (M):** Simple trimming, Dir. Opt BFS, Coloring until less than 100k vertices remain with single thread exploration on backward stage, Serial Tarjan algorithm.
- **Multistep in Kokkos (MK):** Simple trimming, Dir. Opt BFS with small thread owned queues, coloring with fully parallel forward and backward, no Tarjan
- **GPU in Kokkos (GK):** Simple trimming, Dir. Opt BFS with chunking, coloring with delayed exploration
- **GPU min Memory in Kokkos (GKM):** Only utilize out edges, simple trimming, Forward BFS with chunking and fully bottom-up BFS on backward stage, forward coloring with delayed exploration and fully bottom-up reverse search

Results with Kokkos versions



- Multistep (M) is the fastest algorithm in CPU
- **GK** is the fastest algorithm in the GPU
- Phi is (a lot) slower

Conclusion

- **Kokkos provides the path-forward for refactoring codes to different architectures**
 - Handles data layout
 - Portable, ThreadScalable performance
- **Algorithmic Challenges:**
 - Still different algorithms perform better in different architectures
 - Hard to see a single refactor for algorithms in different architectures
- **Kokkos programming is C++ programming not CUDA programming**