

Macroscale simulation of million core systems with the Structural Simulation Toolkit

JOWOG-34

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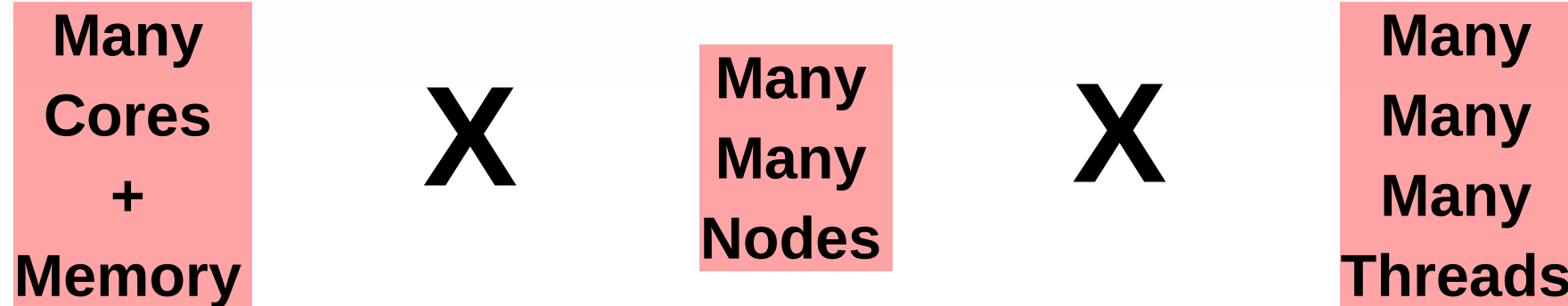
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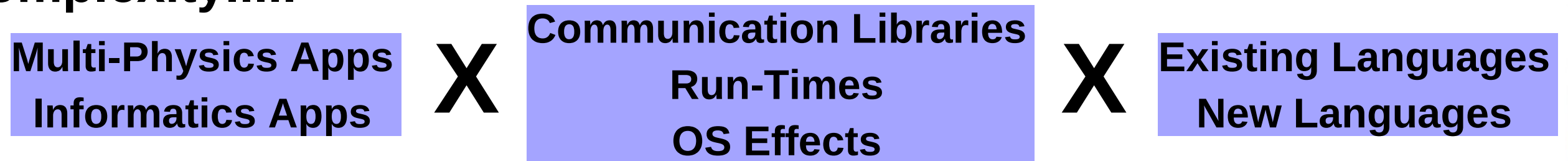
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Designing, developing for, and procuring large-scale system is a multi-faceted problem

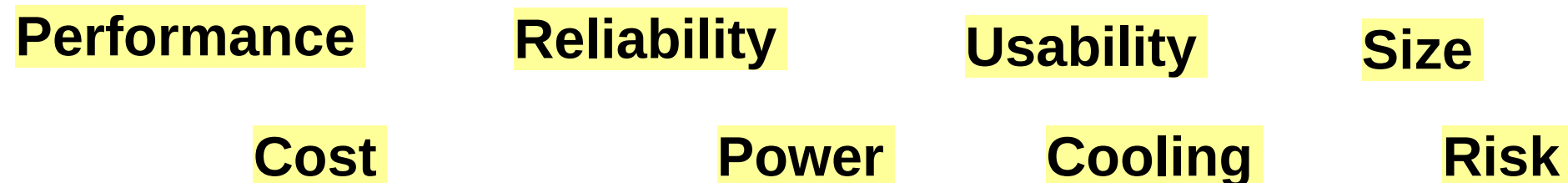
Scale.....



Complexity.....



Constraints.....



Architecture simulation provides a way to explore this large problem space

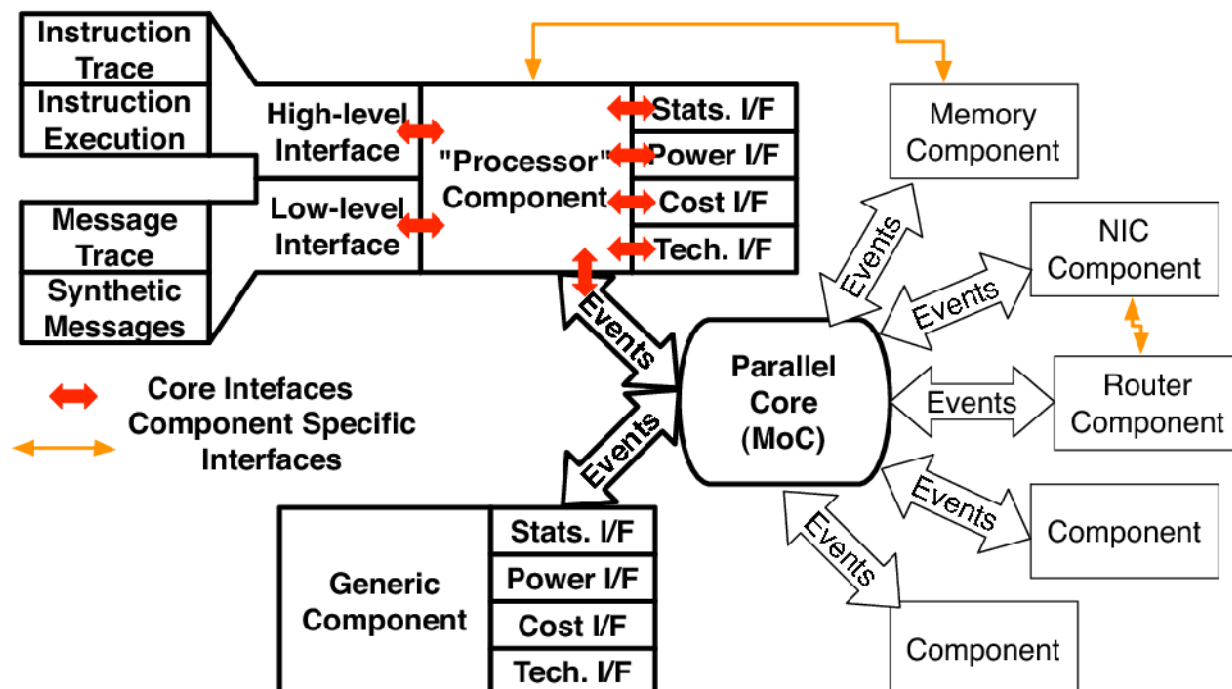
Institute for Advanced Architectures provides locus for community simulation efforts

SST Simulator Goals

- Become **the** standard architectural simulator for the HPC community
- Be able to evaluate future systems on DOE workloads
- Use supercomputers to design supercomputers and applications

Technical Approach

- Multiscale (Cycle-accurate to analytic)
- Parallel (design HPC with HPC)
- Whole system (integrated models)
- SST has core and rich component set



Consortium

- Utilize SST core to interoperate
- “Best of Breed” simulation suite
- Combine Lab, academic, & industry



Sun Microsystems



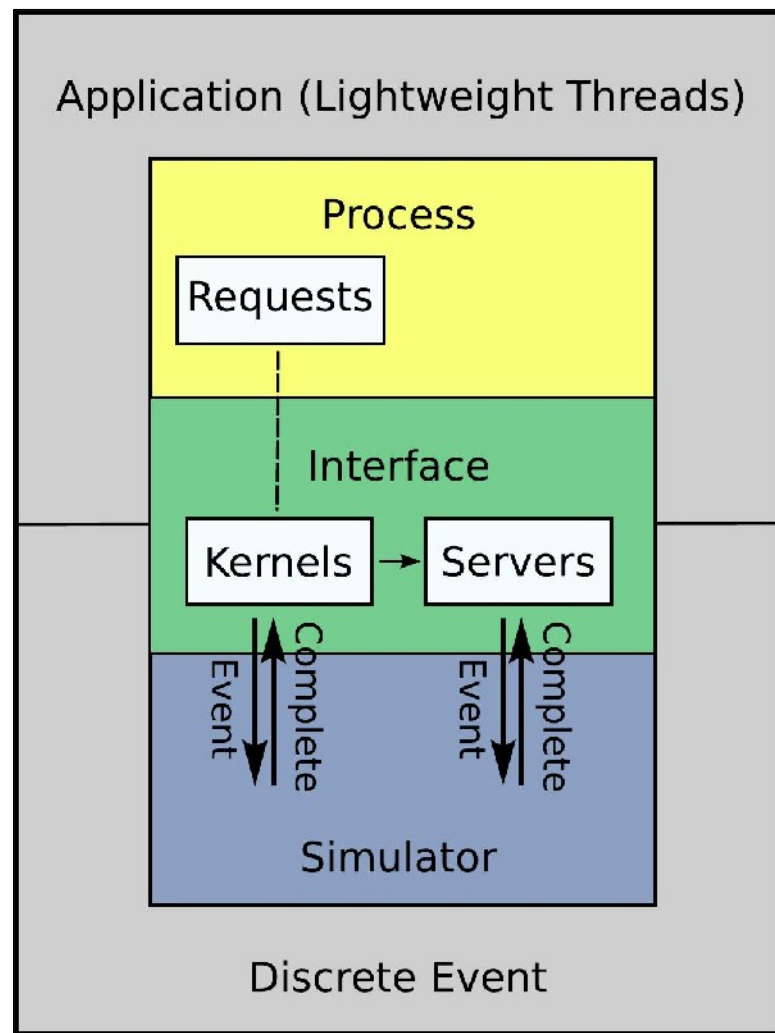
Project goals for Structural Simulation Toolkit macroscale (SST/macro) components

- Take into account the coupling of communication and computation in applications using
 - Trace files that record an actual application run
 - Skeleton applications that mimic application behaviour
- Run very large simulations (complex applications at large scales)
 - 1,000,000's of cores simulated on a single processor
 - Parallelism used for parameter studies & multiscale
- Learn from currently available simulators such as BigSim and the Fujitsu simulator and provide new capabilities beyond that now available
- Allow investigation of effects of
 - Topology, process placement, and interference between jobs
 - Changes in the routing algorithm
 - Changes in the network latency and bandwidth
 - Having many cores share the same network interface
 - Modifications to the MPI layer
 - Modifications to the application
 - Incorporation of more detailed models w/multiscale simulation

SST/macro design

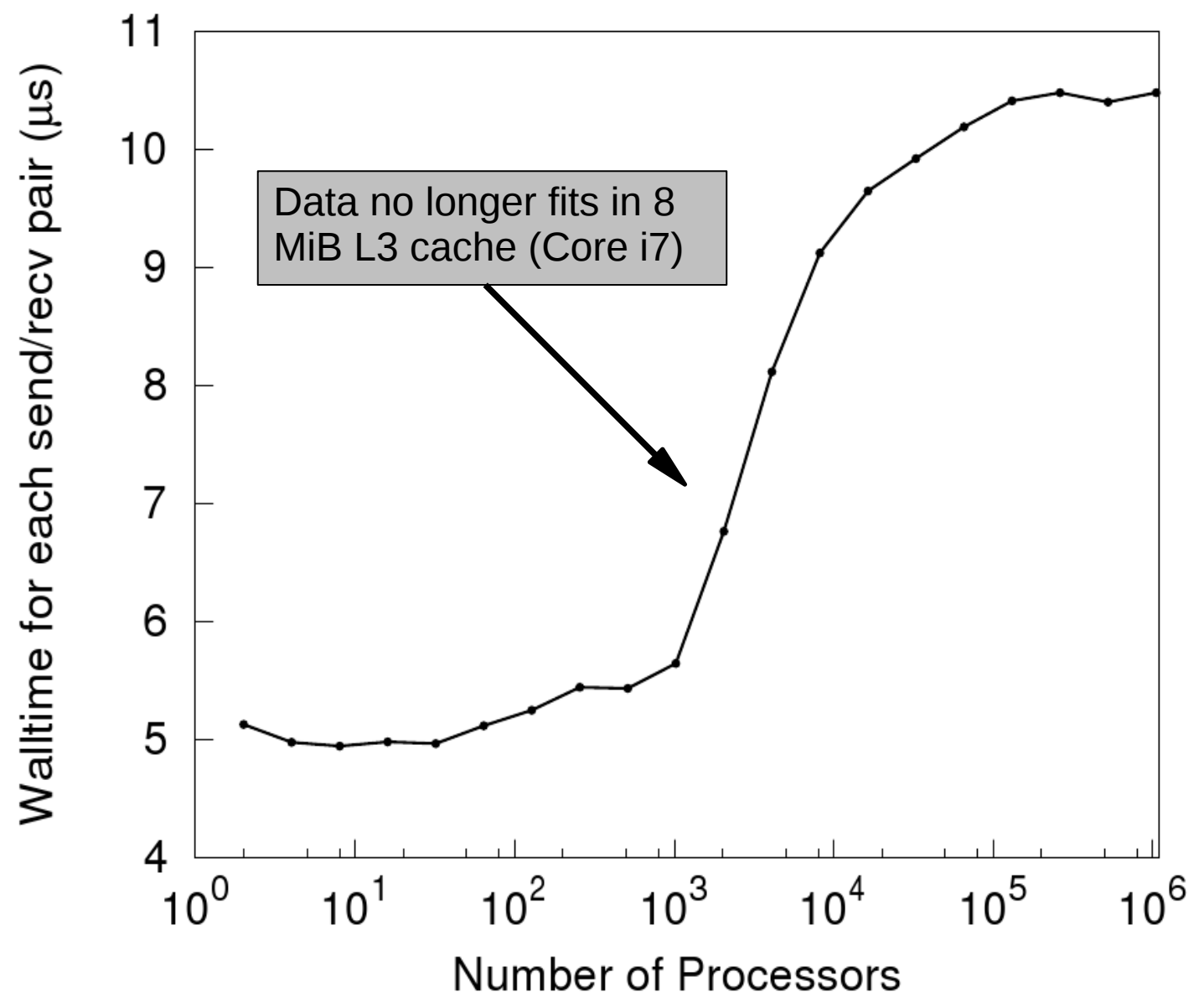
Generic event interface:

permits integration into the hybrid multi-scale SST simulator framework:



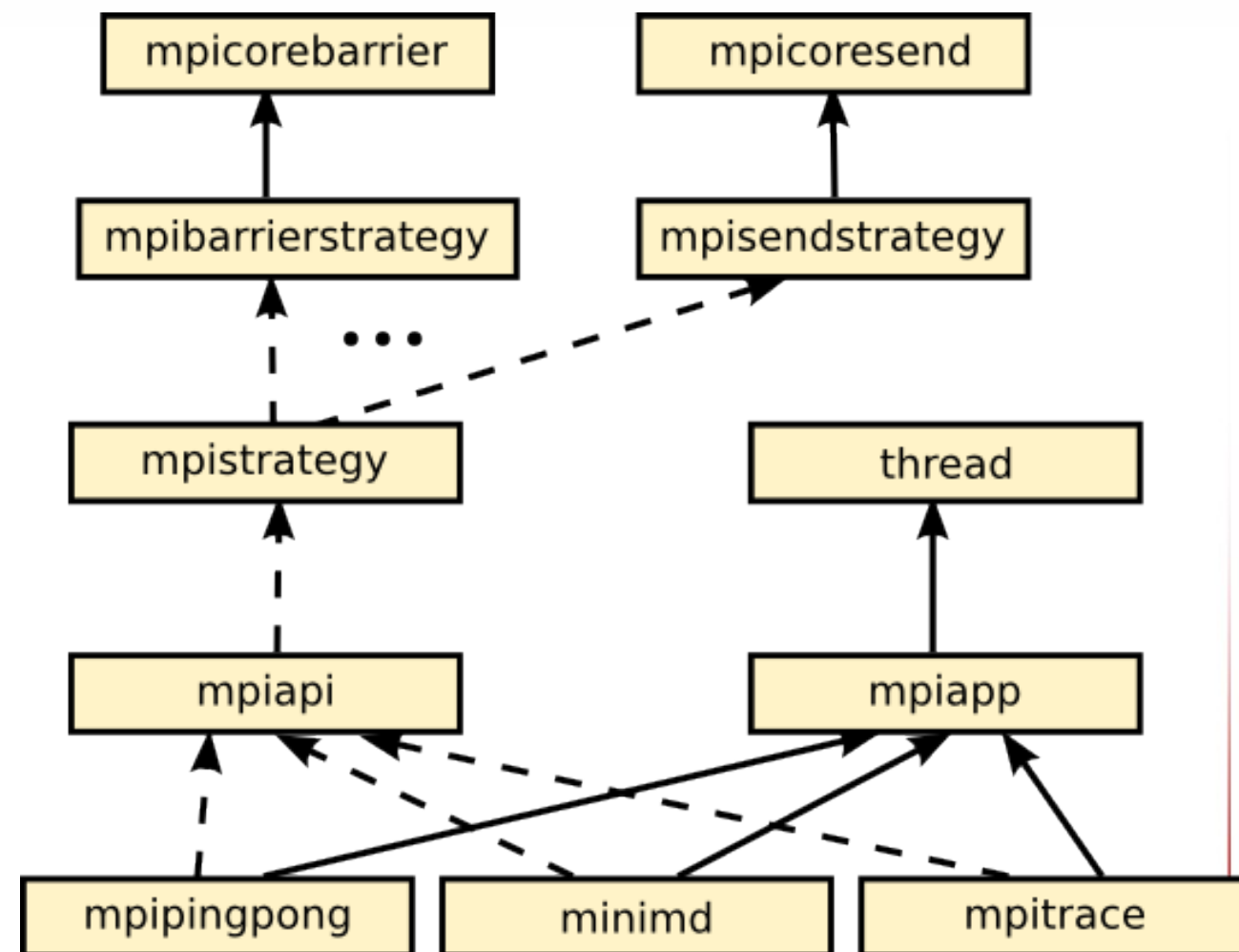
Extremely lightweight events:

Measured real time to perform MPI ping pong round trips (simulator ran on a single processor):



SST/macro has a flexible MPI model

- Traces & skeleton apps record MPI calls—but that can entail a lot
 - MPI model can generate all of the messages that a real MPI library generates ...
 - ... or a simplified model with less overhead can be used
- Allows details of the MPI implementation to be a part of the design space
- Not just restricted to MPI2
 - Other programming models can be easily added
 - Implemented immediate mode collectives (MPI3 proposal)



Examples of SST usage and impact

Obtain traces for applications and compact applications. Use SST for parameter studies.

Write skeleton applications for extreme scale studies.

Simulate new architectural feature such as extended memory semantics and transactional memory.

COMPLEXITY

Develop acceptance tests and estimate performance before machine is built.

Understand performance and issues for machines several years from deployment.

Allow co-design of advance architectures and applications many years before deployment.

SST/macro customers

SNL: studying Zia benchmarks suite and IAA-related mini applications.

ORNL: examining performance of MADNESS.

PNNL: understanding performance of one-sided programming models.

LLNL: mapping of applications into Sequoia network.

Methods for driving the SST/macro simulator

Trace Driven

- Open Trace Format
 - Tools exist to generate trace
 - Visualizers exist for trace
 - Data not complete
- dumpi trace format
 - Custom SST/macro format
 - Records full MPI signature
 - size vectors
 - MPI_Request info
 - Well behaved when application is not
 - Skips irrelevant but resource intensive info (like MPI_Iprobe)

Skeleton Application

- Programmer writes program that behaves like application
 - Skip heavy computation
- Permits extreme-scale runs

```
void mpipingpong::run() {
    this->mpi_->init();
    mpicomm world = this->mpi_->comm_world();
    mpitype type = mpitype::mpi_double;
    int rank = world.rank().id;
    int size = world.size().id;
    if(! ((size % 2) && (rank+1 >= size))) {
        mpiid peer(rank ^ 1);
        mpiapi::const_mpistatus_t stat;
        for(int half_cycle = 0;
            half_cycle < 2*iterations_; ++half_cycle) {
            if((half_cycle + rank) & 1)
                mpi_->send(count_, type, peer,
                           mpitag(0), world);
            else
                mpi_->recv(count_, type, peer,
                           mpitag(0), world, stat);
        }
    }
    mpi_->finalize();
}
```

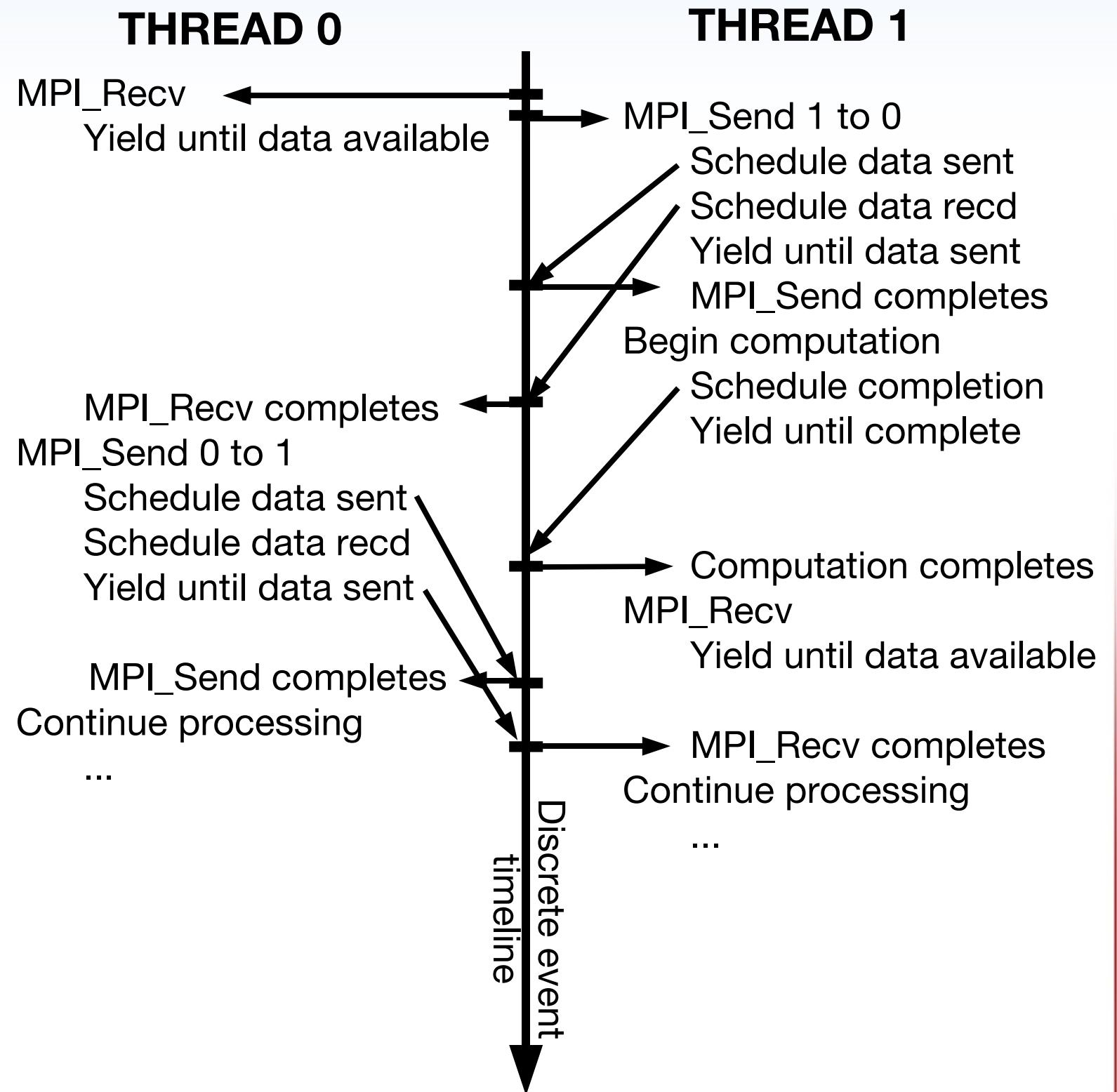

A more sophisticated skeleton app: miniMD

- Simulator runs a skeletonized molecular dynamics application
- Computation time is derived from measurements with various input parameters
 - The calls to estimate computation time is shown in **blue** type
- MiniMD control logic remains mostly intact

```
void minimd::integrate::run(shared_ptr<atom> atm, shared_ptr<force> frc,
                           shared_ptr<neighbor> nbr, shared_ptr<comm> cmm,
                           shared_ptr<thermo> thm, shared_ptr<timer> tmr)
{
    mpiid rank = mpi_->comm_world().rank();
    for(int n = 0; n < this->ntimes; ++n) {
        env_->compute(this->interpolator_->get("integrate::run", 0));
        if((n+1) % nbr->every) {
            cmm->communicate(atm);
        }
        else {
            cmm->exchange(atm);
            cmm->borders(atm);
            nbr->build(atm);
        }
        frc->compute(atm, nbr);
        env_->compute(this->interpolator_->get("integrate::run", 1));
        if(thm->nstat)
            thm->compute(n+1, atm, nbr, frc);
    }
}
```

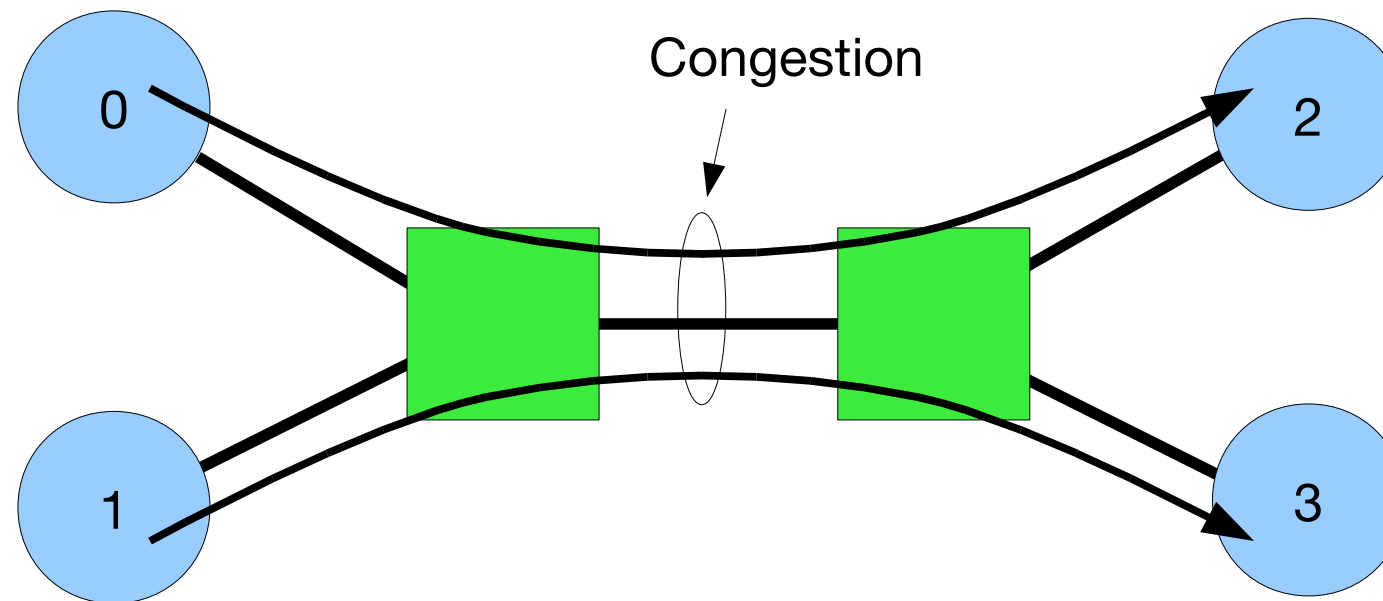
Translating an application into discrete events

- MPI ping pong with computation: Simulation time increases going down, events are bars on the time axis, thread 0 executes code on left and process 1 on right
- Each thread inserts events into the discrete event queue until it yields
- Events can cause a process to resume execution
- Single MPI calls can result in multiple events
- Scheduling of data sent/received events depends on network traffic



The circuit network component: a simple model that includes network congestion effects

Example: Two pairs of nodes try to use the same network link simultaneously:

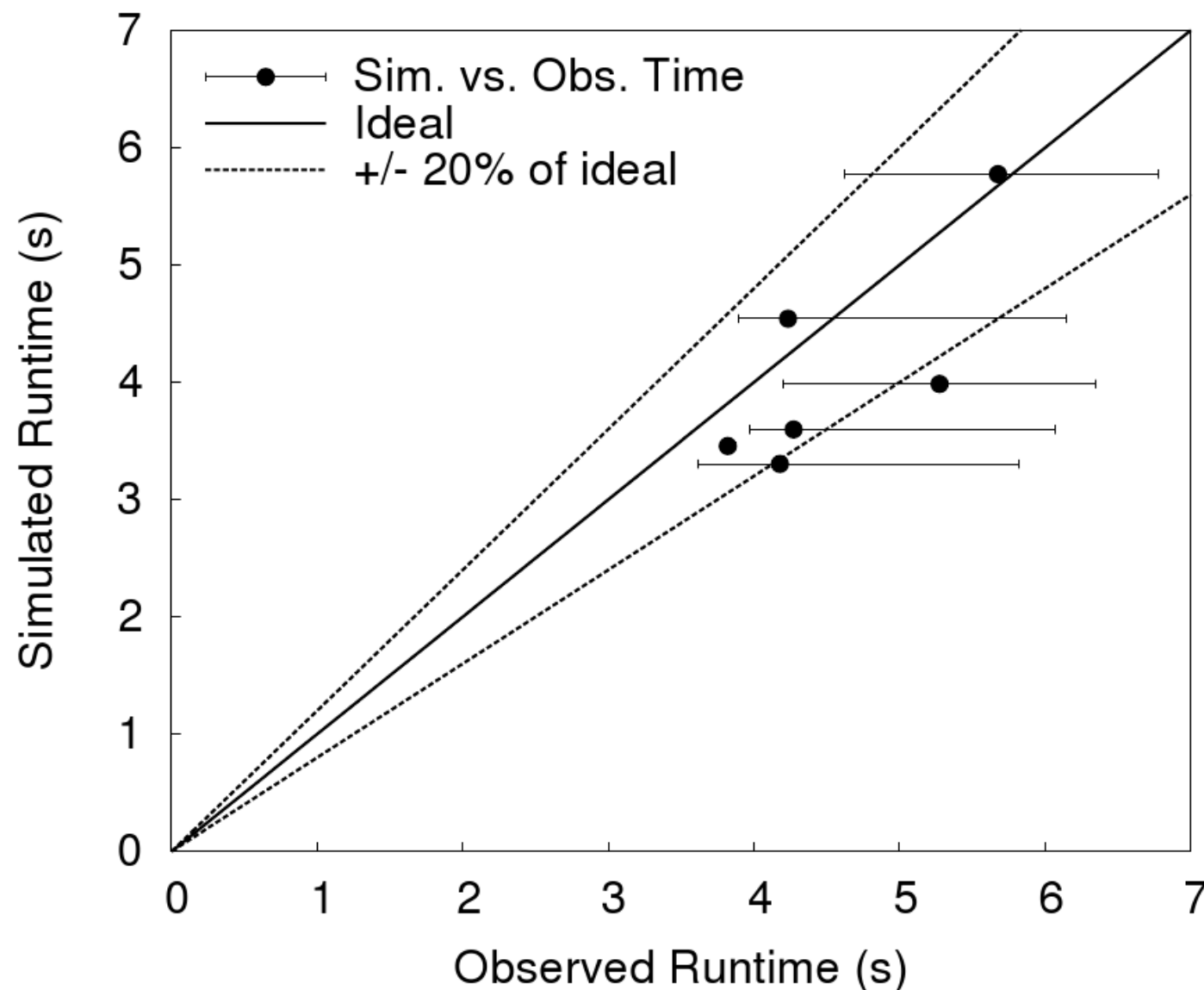


Circuit model handles this case as follows:

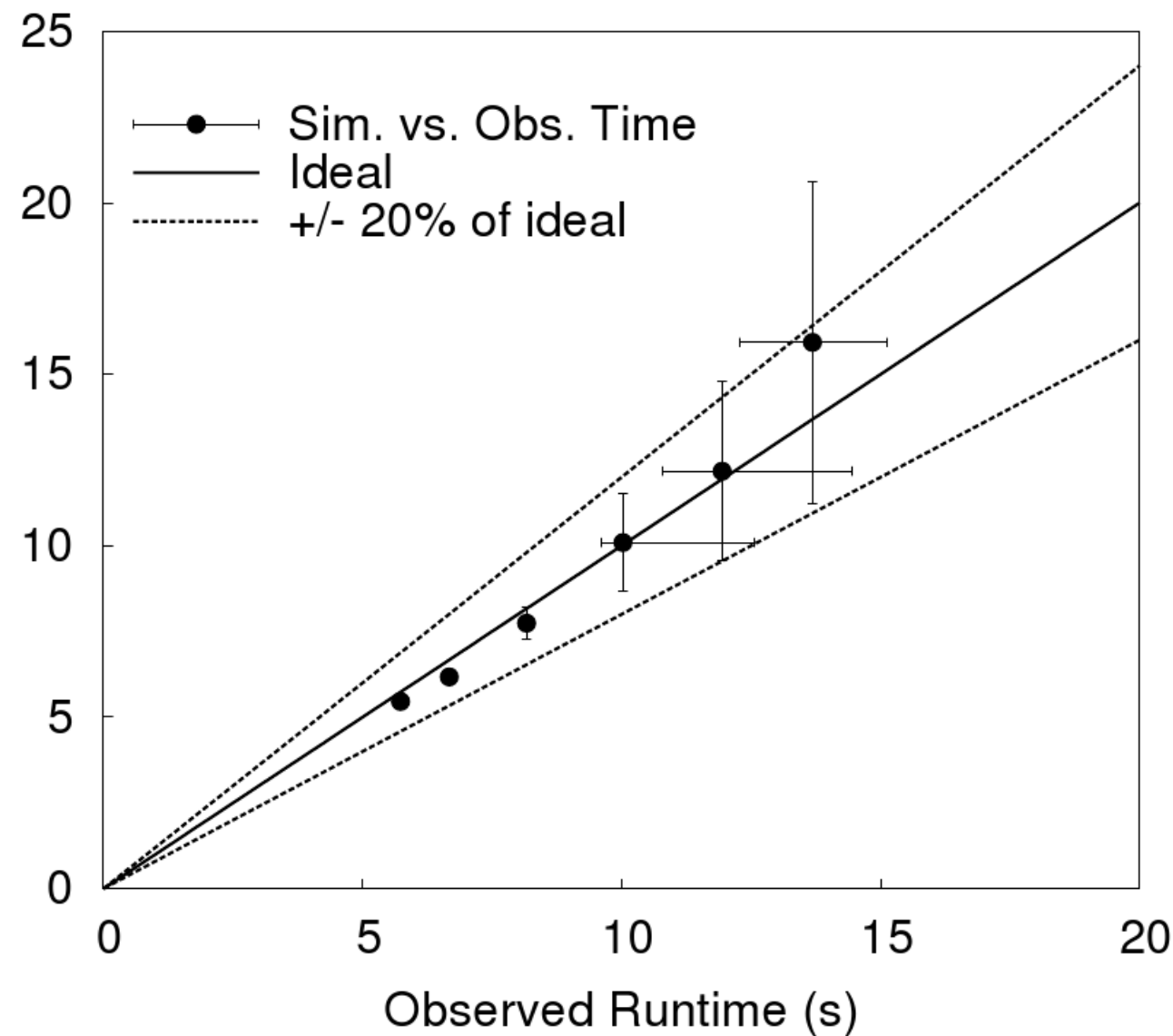
- $T = 0s$: Traffic begins from node 1 to node 3; duration is 3s.
- $T = 1s$: Attempt to begin traffic from node 0 to 2; duration is 1s. Attempt fails so attempt is rescheduled at $T = 3s$.
- $T = 3s$: Traffic from node 1 to node 3 completes.
- $T = 3s$: Traffic from node 0 to 2 begins; duration is 1s.
- $T = 4s$: Traffic from node 0 to 2 completes.

Validation of simulator

- Used AMG2006: part of NNSA ASC/Sequoia acceptance tests
- Collected traces on the Thunderbird machine and played back through simulator



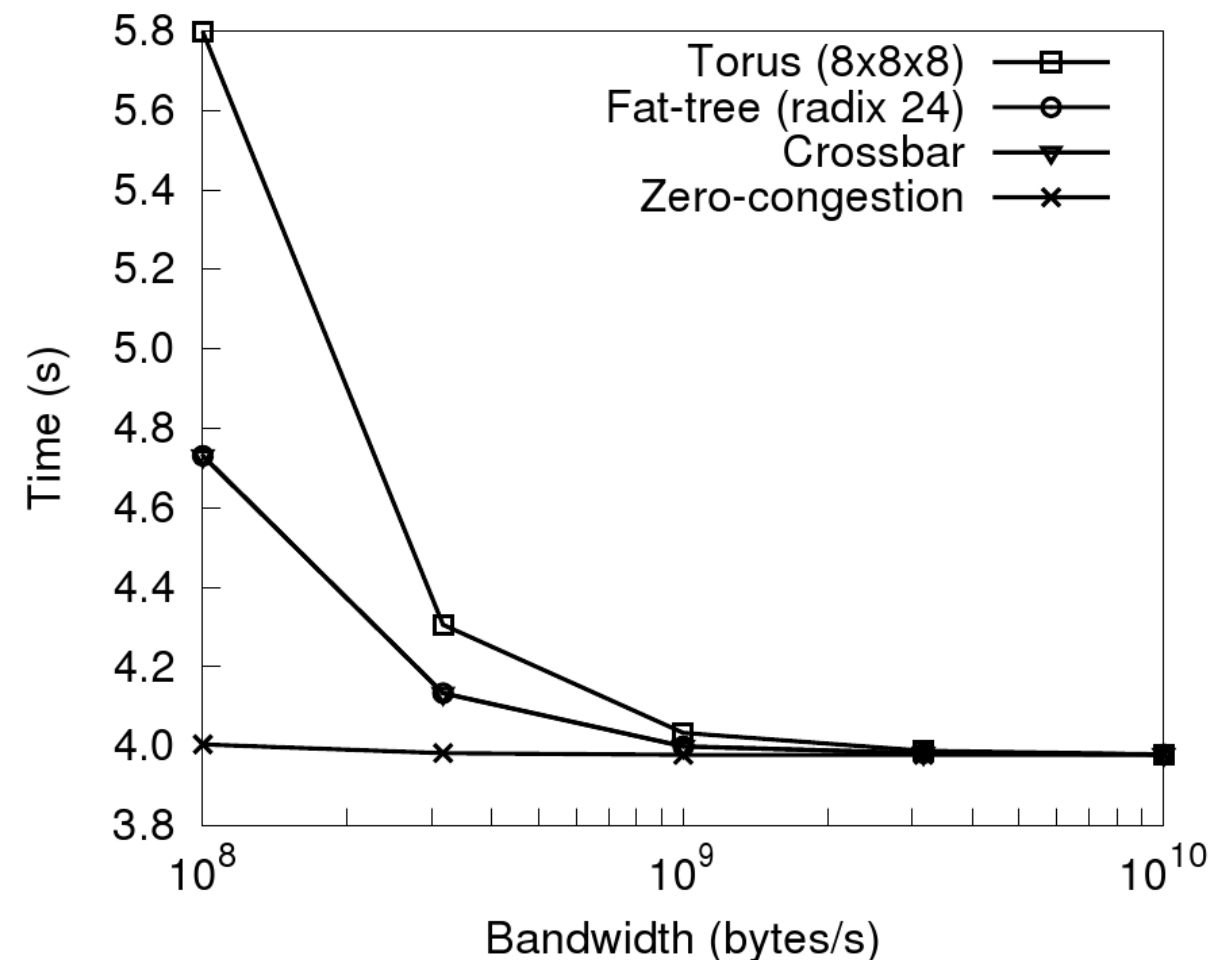
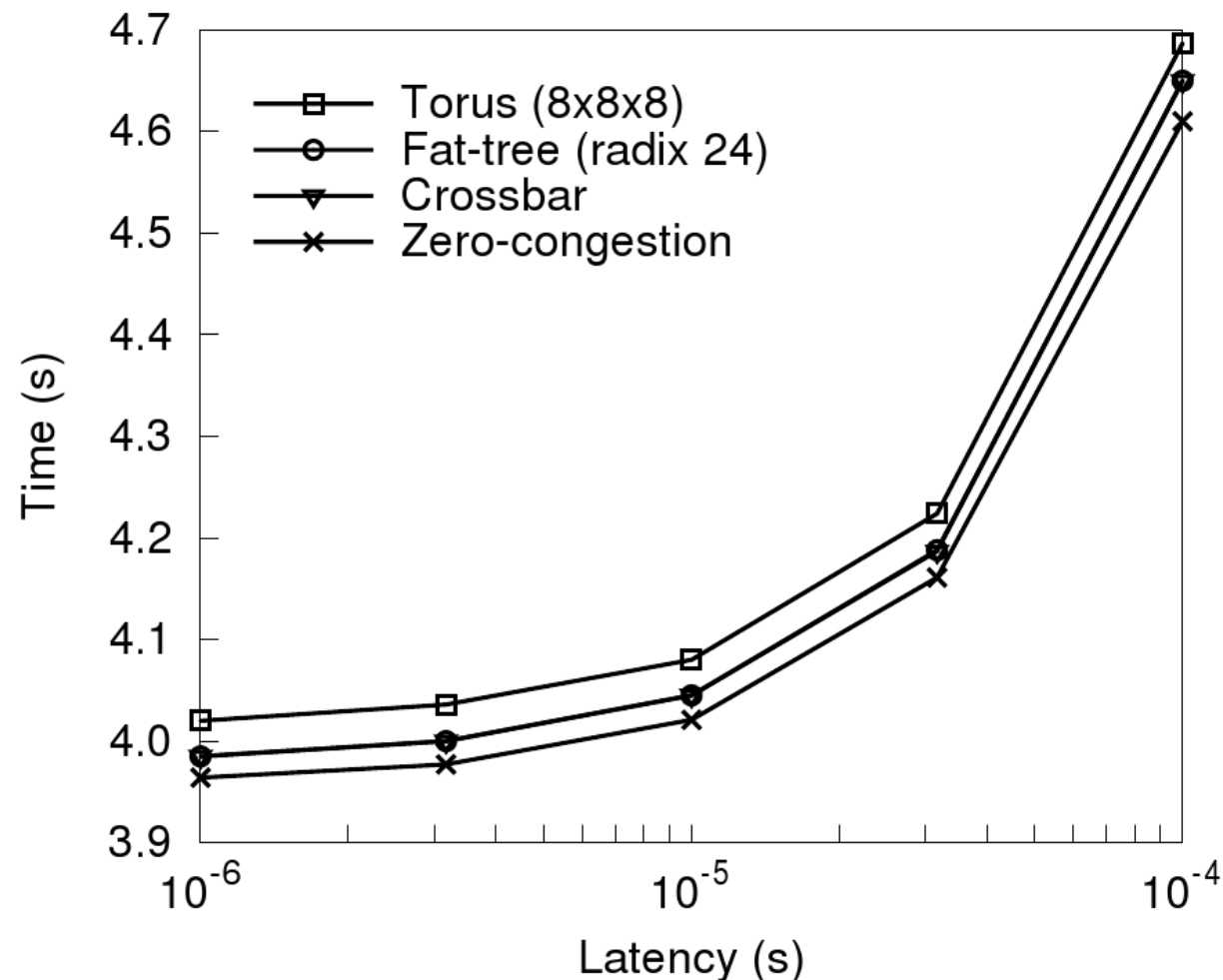
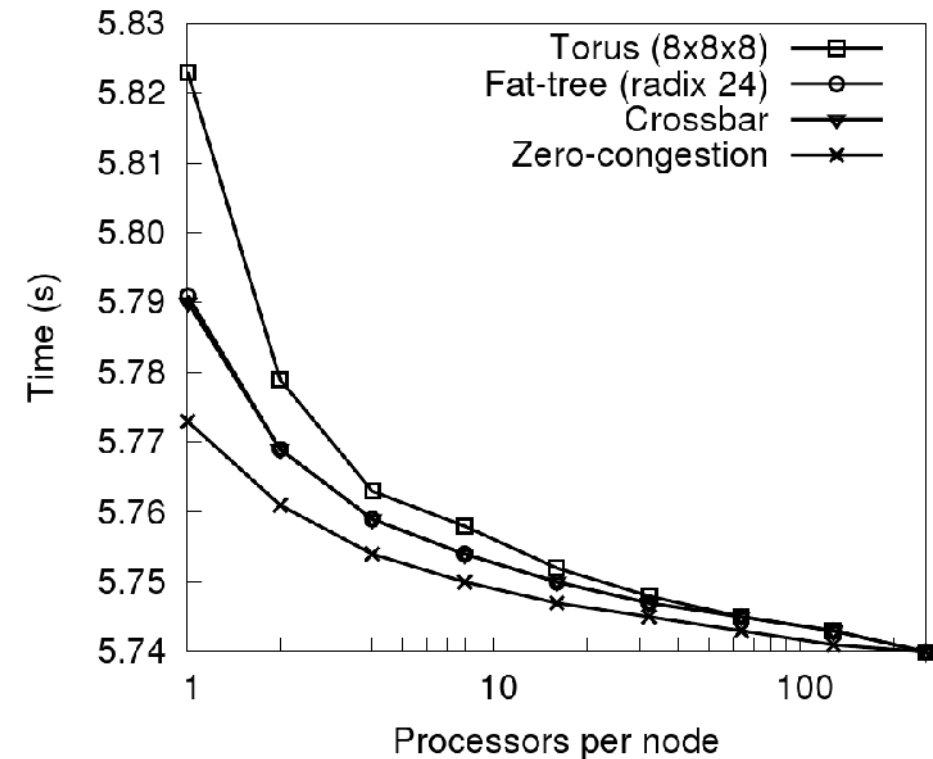
“Narrow” decomposition



“Fat” decomposition

Sensitivity of AMG2006 to architectural parameters

- Examined simulated time to solution as bandwidth, latency, and processors per node are varied for several topologies.



Current Work

- Extended validation studies
 - Using Red Storm Qualification system
 - Can control process placement
 - Results are highly consistent from run to run
- Advanced routing
 - Dispersive routing
 - Adaptive routing
- Trace file format (dumpi)
 - Finalizing full MPI support
- User interface and visualization
 - Developing a GUI to simplify problem setup
 - Working with others to visualize network congestion, etc.
- Processor models
 - Developing more sophisticated but inexpensive processor models
- Integration into the SST/Core interface