

Detecting Insider Activity Using Enhanced Directory Virtualization

October 8, 2010

**William R. Claycomb
Member of Technical Staff
Sandia National Laboratories**

**Dongwan Shin
Associate Professor
New Mexico Tech**



Agenda

- **Background on Insider Threats**
- **Background on Directory Services**
 - **Directory Virtualization**
- **Enhancing Directory Virtualization**
- **Implementation Examples**
- **Primary Contributions**
- **Future Work**



Insider Threats

- **“Insider Threat Study: Illicit Cyber Activity in the Government Sector”, a study conducted by U.S. Secret Service and CERT (2008) found:**
 - **Most of the insiders had authorized access at the time of their malicious activities**
 - **Access control gaps facilitated most of the insider incidents, including:**
 - **The ability of an insider to use technical methods to override access controls without detection**
 - **System vulnerabilities that allowed technical insiders to use their specialized skills to override access controls without detection**





Agenda

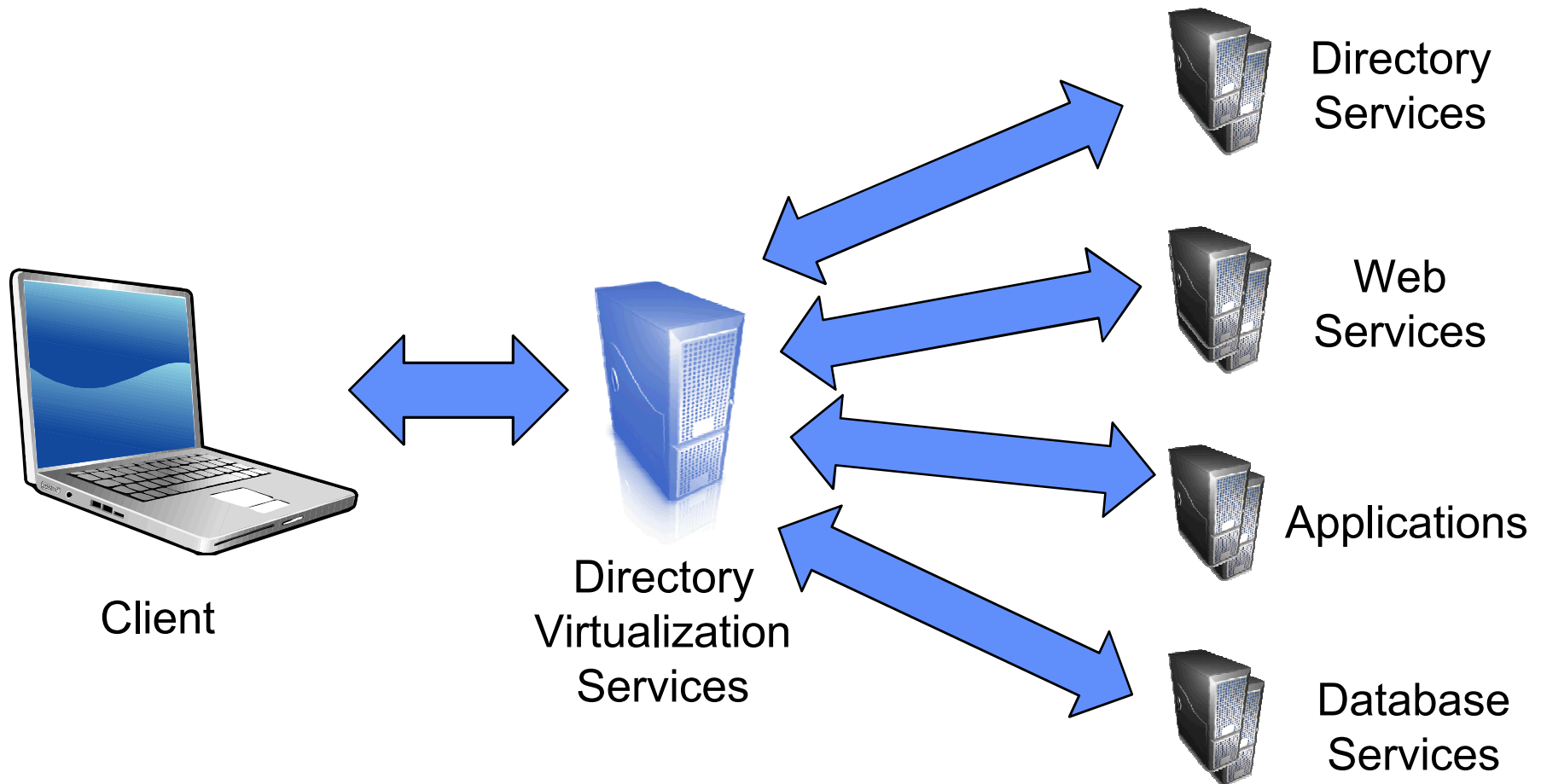
- Background on Insider Threats
- **Background on Directory Services**
 - **Directory Virtualization**
- Enhancing Directory Virtualization
- Implementation Examples
- Primary Contributions
- Future Work



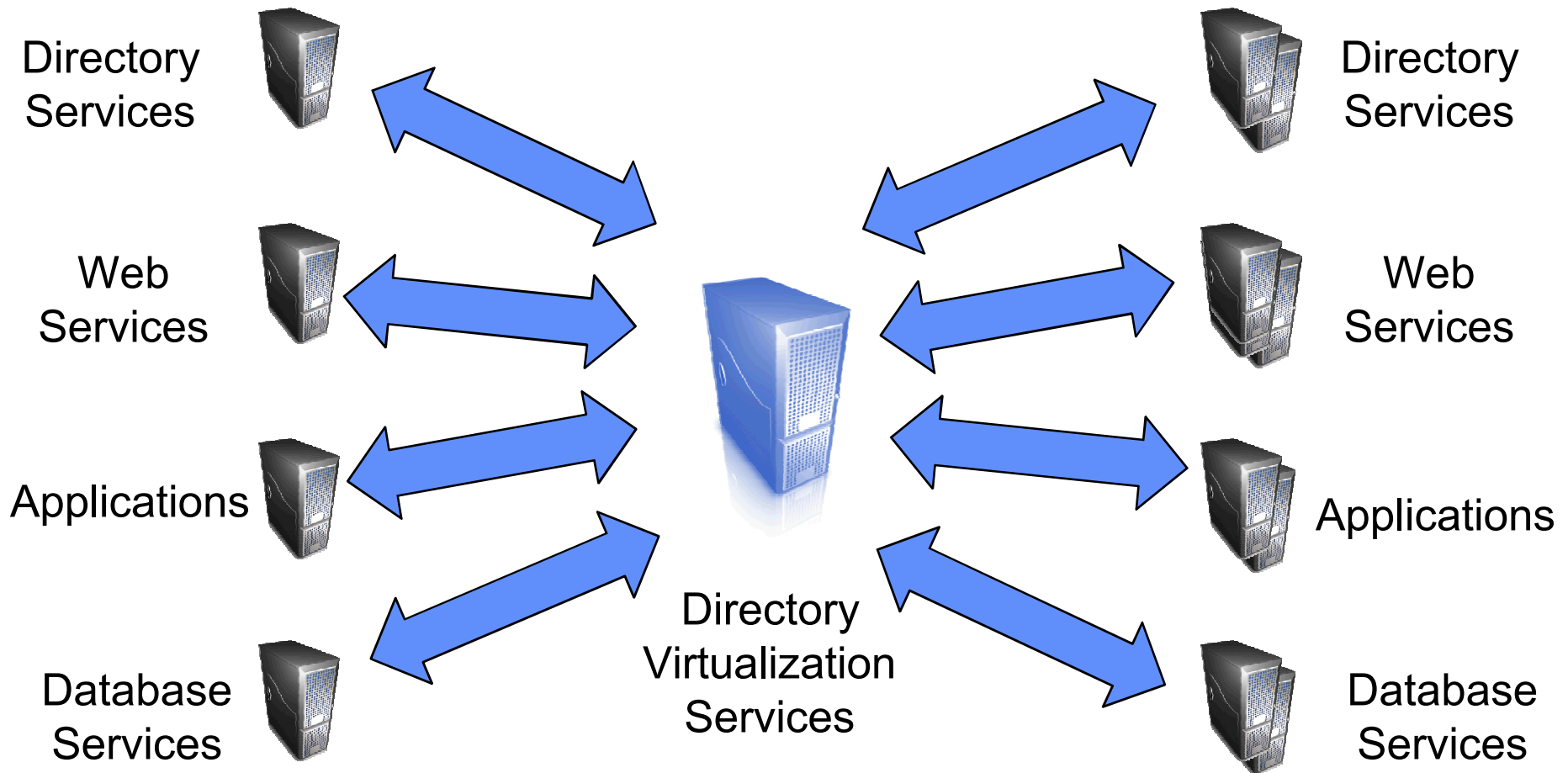
Directory Services

- **Localized data store containing information about objects (Users, Computers, Contacts, etc.)**
- **Provide information to applications**
 - Authentication and access control
 - Contact information
 - Group membership
- **Popular Directory Services Implementations**
 - Windows Server *Active Directory*
 - IBM *Tivoli*
 - *OpenLDAP*
 - Fedora *Directory Server*
 - Sun JAVA System *Directory Server*

Directory Virtualization



Directory Virtualization and Synchronization





Attacks on Directory Services

- **Create a new user**
 - Allows insider to impersonate another account, rather than using her own
- **Change group membership**
 - Modify ACL to gain unauthorized access to protected resources
- **Reset user password**
 - Allows insider to masquerade as an existing user
 - Used to circumvent processes that monitor for unauthorized account creation





Agenda

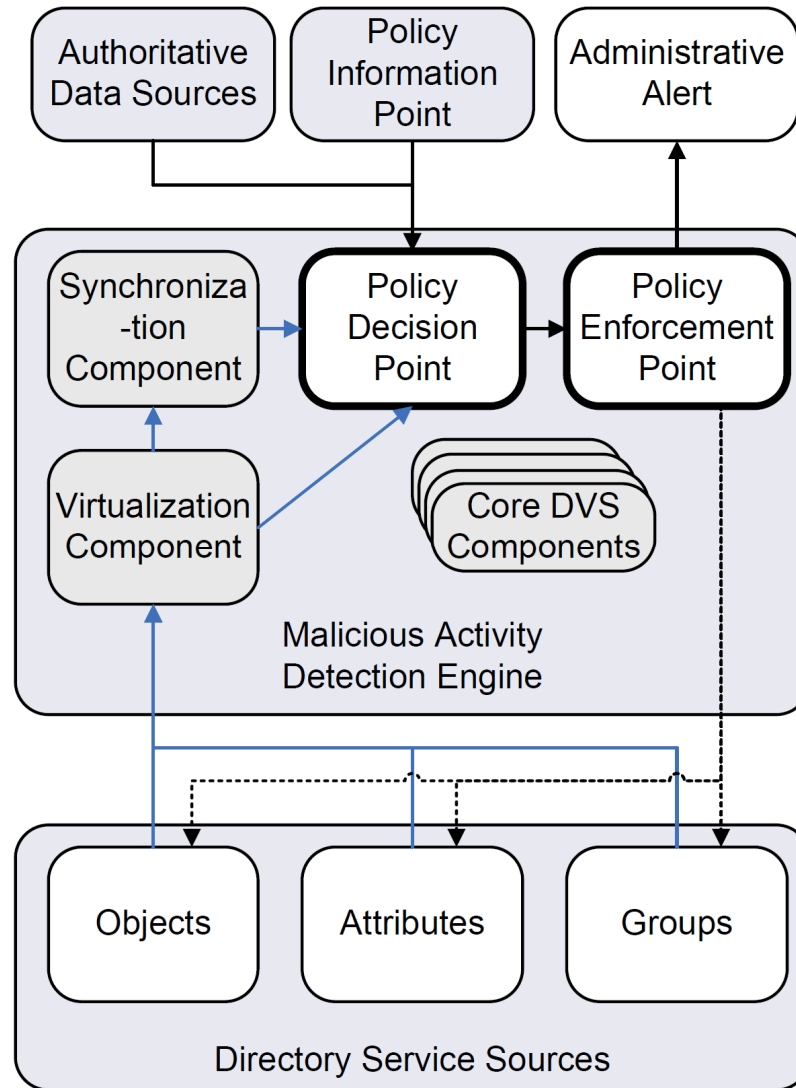
- Background on Insider Threats
- Background on Directory Services
 - Directory Virtualization
- **Enhancing Directory Virtualization**
- Implementation Examples
- Primary Contributions
- Future Work



Enhancing Directory Virtualization

- **Malicious Activity Detection Engine (MADE)**
 - **Builds on core components**
 - **Virtualization**
 - **Synchronization**
 - **Adds Policy Components**
 - **Policy Decision Point (PDP)**
 - **Policy Information Point (PIP)**
 - **Policy Enforcement Point (PEP)**

Architecture





Policy Decision Point

- **Detects events**
- **Obtains policy information from PIP**
- **Makes determination on activity detected**
- **Coordinates response using PEP**



Modeling the Problem

- **Boolean function describing an account and its relationships with other data sources**

$$s_i \equiv s_j(a_1, a_2, \dots, a_j)$$

- $s_j(a_1, a_2, \dots, a_j) = 1$ **implies** $a_1 = 1, a_2 = 1, \dots, a_j = 1$
- **Disjunctive conditions can be included in separate conjunctive states**



Function Conditions

- **Variety of possibilities**
 - **Value (or presence) of a specific attribute on an account**
 - **Membership of an account in a particular group**
 - **Relationship between two objects and their attributes**
 - **Etc.**



Example

- **Unauthorized account creation**

- **Policy:** $S_{new_acct}(a_X, a_Y, a_Z)$

- a_X : presence of account in data source X
 - a_Y : “active” status of the account in data source Y
 - a_Z : presence of specific training records for the account owner in data source Z



Agenda

- Background on Insider Threats
- Background on Directory Services
 - Directory Virtualization
- Enhancing Directory Virtualization
- **Implementation Examples**
- Primary Contributions
- Future Work



Implementation: Hypothetical Organization

- **DVS implemented with Microsoft Visual Studio**
- **Directory Services**
 - Microsoft Active Directory
 - Sun Java System Directory Server
- **Database Services**
 - Microsoft SQL
- **PIP integrated into DVS**
 - Could have been Windows Server Group Policy, etc.
- **Tests**
 - 15000 user objects
 - 4 domain controllers, 1 SQL database



Implementation: Creating Unauthorized User Accounts

- Insider creates new user to access unauthorized resource
- Policy to prevent such action:

$$S_{auth_acct}(a_{auth_acct})$$

- Virtualization service monitors directory service for changes, passes events to synchronization service
- Sync service reports new account to PDP
- PDP requests information from auth_acct,
- Applies information to policy retrieved from PIP
- Because $S_{auth_acct}(a_{auth_acct}) = 0$, PDP notifies PEP of actions to take, based on information from PIP
- PEP takes appropriate action
 - disable account
 - notify admin



Implementation: Creating Unauthorized User Accounts

- **Test**

- Use administrative controls to create a new account in directory service

- **Performance**

- Unauthorized account disabled by MADE within 1 sec of detection



Implementation: Changing Group Membership

- **Object attributes often used to make access control decisions**
 - **Group membership, clearance level, citizenship, organizational status, management level, etc.**



Implementation: Changing Group Membership

- **Policy:**

$$S_{group_membership} \left(a_{auth_admn}, a_{mbr_org}, a_{grp_size} \right)$$

- a_{auth_admn} : **Account is listed as valid domain administrator in *auth_admin***
- a_{mbr_org} : **Account owner is part of specified organization in a separate data source**
- a_{grp_size} : **Three or less accounts can exist in domain administrators group at any one time**



Implementation: Changing Group Membership

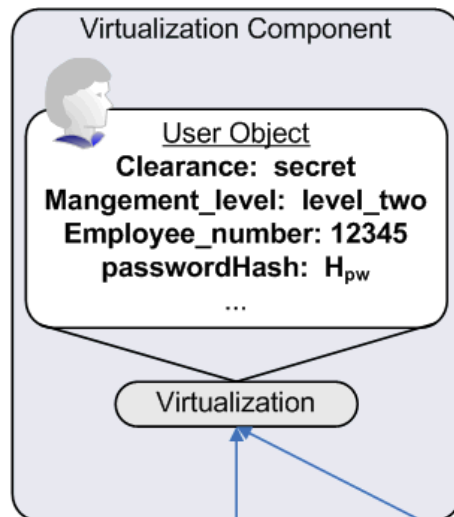
- **Attack:** Adding an unauthorized account to *Domain Administrators* group
- **Detection** similar to previous example
- a_{auth_admin} and a_{mbr_org} are determined to be false
- a_{grp_size} may also be exceeded
- **PEP** notified to take appropriate action
- Interesting to note that policy decision used information from the directory service itself
 - Group size
- **Performance** and results similar to previous case



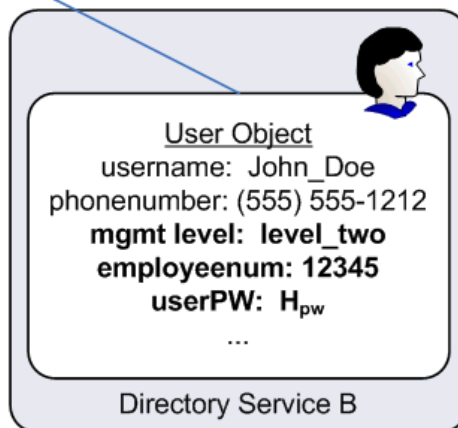
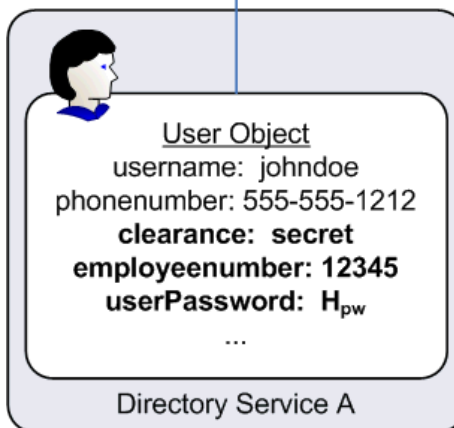
Implementation: Using Cross-Directory Relationships

- Two different directory services
- Users required to have matching accounts in both directories
 - Directory D_A contains mandatory clearance level information
 - Directory D_B contains mandatory management level information
- Administrators are not allowed to administer both directories simultaneously

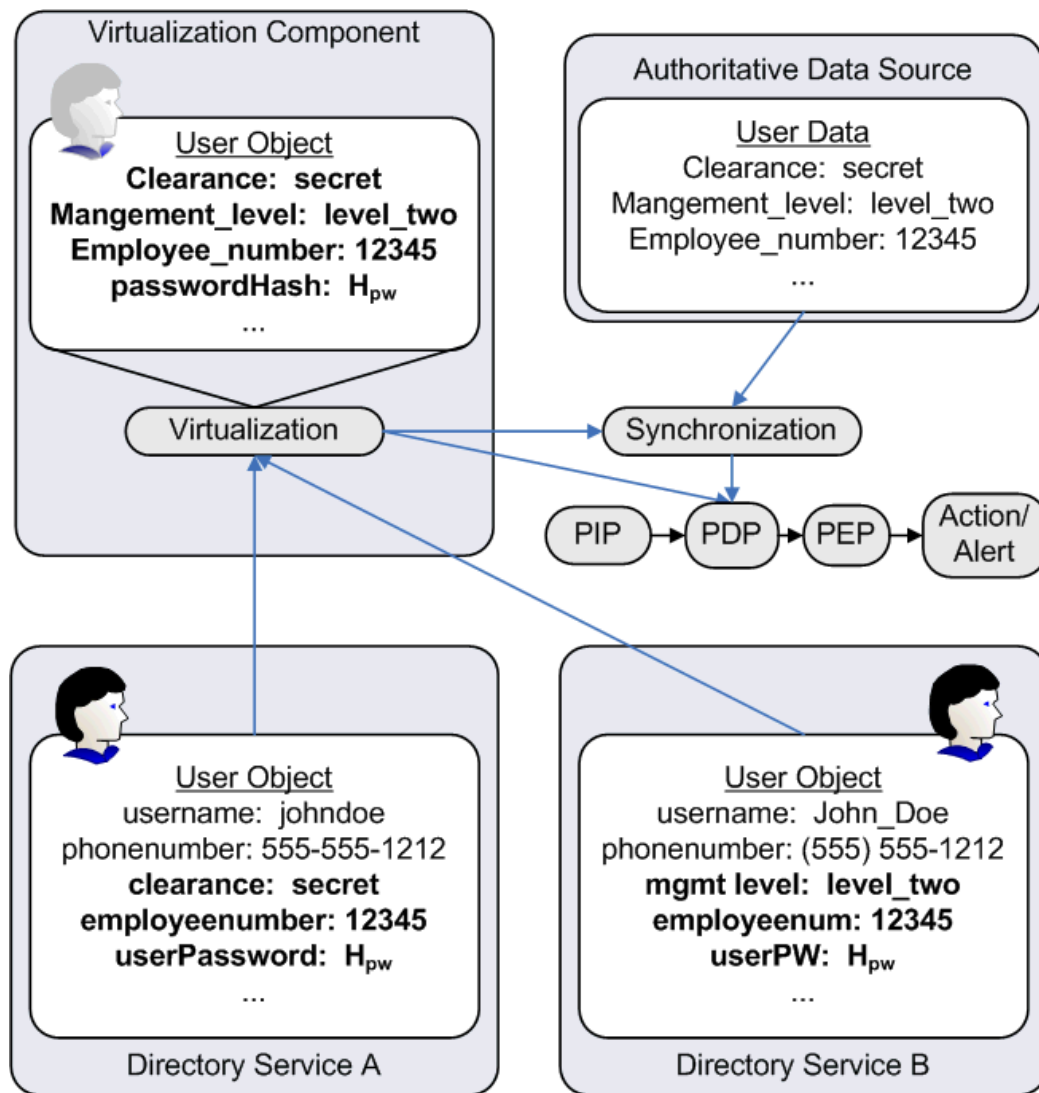
Implementation: Using Cross-Directory Relationships



- Virtualization component combines data from D_A and D_B to form a **single view**
- Can determine authoritative source, or consider a discrepancy a potential attack, and notify PDP



Implementation: Using Cross-Directory Relationships





Implementation: Using Cross-Directory Relationships

- **Policy:**

$$S_{auth_acct} \left(a_{clearance}, a_{mgmt_level}, a_{pwd_match} \right)$$

- $a_{clearance}$: **presence of correct clearance level**
- a_{mgmt_level} : **presence of correct management level**
- a_{pwd_match} : **synchronization of passwords between domains**



Implementation: Using Cross-Directory Relationships

- **Attack:** Insider changes a user password in D_A in order to impersonate that account
- **Performance similar to previous examples**
 - Detection and enforcement within one second of detection



Agenda

- Background on Insider Threats
- Background on Directory Services
 - Directory Virtualization
- Enhancing Directory Virtualization
- Implementation Examples
- **Primary Contributions**
- Future Work



Primary Contributions

- Monitoring for unauthorized changes is not new
- Adding policy engine in this form is somewhat novel
- Real contribution is the ability to monitor multiple systems simultaneously, and present a consolidated **systems view** of key data and relationships.
- Insider may be able to modify individual components, but is less likely to be able to modify all monitored components simultaneously

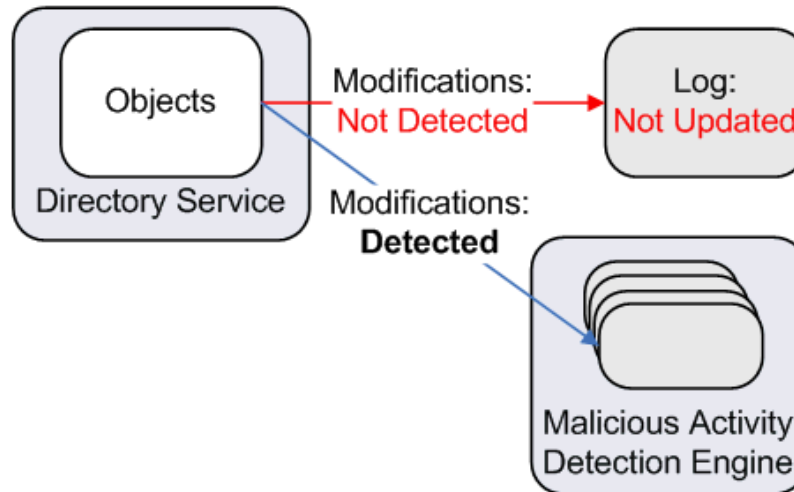
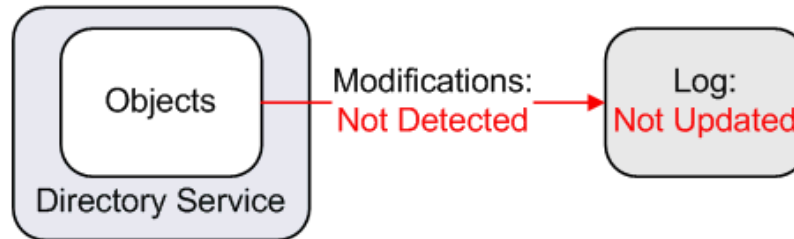
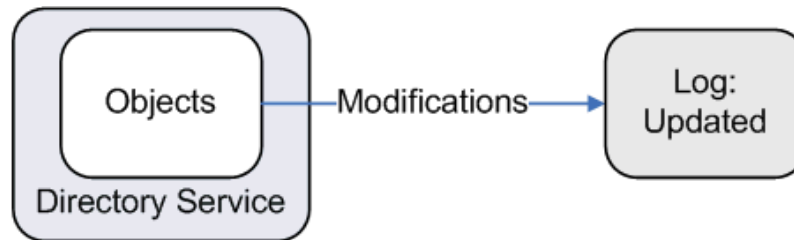


Primary Contributions

- **DVS present a lightweight footprint**
 - Many DVS may be distributed across the network
 - Makes disabling or defeating them more difficult
- **DVS exist outside normal directory services**
 - Enhances the chances of detecting a savvy insider
 - Basically a **covert channel** for directory services



Covert Channel





Attacks on MADE

- **Administrator with control of authoritative data sources**
 - Separation of duties
 - Careful monitoring and auditing
- **Administrator access to DVS**
 - Disable monitoring
 - Makes key changes to specific repositories simultaneously
- **Administrator access to administrative applications**
 - Makes unauthorized changes appear legitimate
- **Administrative access or attacks on PIP or PEP**





Future Work

- **Additional performance testing under real-world scenarios and in production environments**
- **Apply MADE to other access control models**
- **Deployment of many lightweight DVS**
- **Enhanced interaction with PIP**
 - **What other systems can be used as PIP**
 - **How to prevent PIP compromise**





Questions

<http://www.sandia.gov>

<http://scl.cs.nmt.edu>

wrclayc@sandia.gov

doshin@nmt.edu

**This work was partially supported at
the Secure Computing Laboratory at
New Mexico Tech by the grant from
the National Science Foundation
(NSF-IIS-0916875)**

