



Tall and Skinny SAND2011-4037C *QR Factorizations in MapReduce*

David F. Gleich

Paul G. Constantine



*MapReduce 2011 Workshop
8 June 2011*

*Thanks to Mark Hoemmen, Tammy Kolda,
Todd Plantenga, Craig Ulmer, Justin Basilico
@ Sandia and
James Demmel, Gray Ballard, Austin
Benson @ Berkeley for helpful discussions.*

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MapReduce is great for TSQR!

Data A tall and skinny (TS) matrix by rows

Map QR factorization of local rows

Reduce QR factorization of local rows

Demmel et al. showed that this construction works to compute a QR factorization with *minimal communication*



Input **500,000,000-by-100 matrix**

Each record **1-by-100 row**

HDFS Size **423.3 GB**

Time to compute $\|\mathbf{A}\mathbf{e}_i\|$ (the norm of each column) **161 sec.**

Time to compute \mathbf{R} in $\text{qr}(\mathbf{A})$ **387 sec.**

`git clone https://github.com/dgleich/mrtsqr`

On a 64-node Hadoop cluster with 4x2TB, one Core i7-920, 12GB RAM/node

Tall and Skinny QR factorizations

Outline

Implementation in MapReduce

Synthetic experiments for performance

Real-world test: Tinyimages

Simulation Informatics

QR Factorization

Let $\mathbf{A} : m \times n, m \geq n$, real

$$\mathbf{A} = \mathbf{QR}$$

\mathbf{Q} is $m \times n$ orthogonal ($\mathbf{Q}^T \mathbf{Q} = \mathbf{I}$)

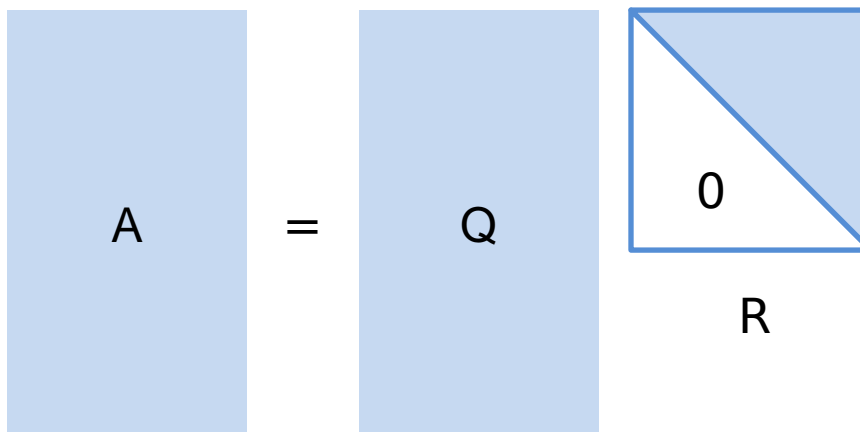
\mathbf{R} is $n \times n$ upper triangular.

Using QR for regression

$\min \|\mathbf{b} - \mathbf{Ax}\|$ is given by the solution of $\mathbf{Rx} = \mathbf{Q}^T \mathbf{b}$

QR is block normalization

“normalize” a vector
usually generalizes to
computing \mathbf{Q} in the QR



A



regression with many samples
block iterative methods
panel factorizations

general linear models
with many samples

all of these applications
need a QR factorization
some only need R !



Use $\mathbf{A}^T \mathbf{A}$? Yields only $O(\sqrt{\epsilon})$ accuracy.

Communication avoiding TSQR

$$A = \begin{bmatrix} A_1 \\ A_2 \\ A_3 \\ A_4 \end{bmatrix} \quad A = \underbrace{\begin{bmatrix} Q_1 & & & \\ & Q_2 & & \\ & & Q_3 & \\ & & & Q_4 \end{bmatrix}}_{8n \times 4n} \underbrace{\begin{bmatrix} R_1 \\ R_2 \\ R_3 \\ R_4 \end{bmatrix}}_{4n \times n}$$

First, do QR factorizations of each local matrix A_i

Second, compute a QR factorization of the new “R”

$$A = \underbrace{\begin{bmatrix} Q_1 & & & \\ & Q_2 & & \\ & & Q_3 & \\ & & & Q_4 \end{bmatrix}}_{8n \times 4n} \underbrace{\begin{bmatrix} \tilde{Q} \end{bmatrix}}_{4n \times n} \underbrace{\begin{bmatrix} \tilde{R} \end{bmatrix}}_{n \times n}$$

=Q

Demmel et al. 2008. Communicating avoiding parallel and sequential QR.

Communication avoiding TSQR

$$A = \begin{bmatrix} A_1 \\ A_2 \\ A_3 \\ A_4 \end{bmatrix} \quad A = \underbrace{\begin{bmatrix} Q_1 & & & \\ & Q_2 & & \\ & & Q_3 & \\ & & & Q_4 \end{bmatrix}}_{8n \times 4n} \underbrace{\begin{bmatrix} R_1 \\ R_2 \\ R_3 \\ R_4 \end{bmatrix}}_{4n \times n} = \underbrace{\begin{bmatrix} Q_5 & \\ & Q_6 \end{bmatrix}}_{4n \times 2n} \underbrace{\begin{bmatrix} R_5 \\ R_6 \end{bmatrix}}_{2n \times n}$$

But we actually
do the second
step recursively!

$$A = \overbrace{\underbrace{\begin{bmatrix} Q_1 & & & \\ & Q_2 & & \\ & & Q_3 & \\ & & & Q_4 \end{bmatrix}}_{8n \times 4n} \underbrace{\begin{bmatrix} Q_5 \\ & Q_6 \end{bmatrix}}_{4n \times 2n}}^{=Q} \underbrace{\begin{bmatrix} Q_7 \\ & R_7 \end{bmatrix}}_{2n \times n}$$

Demmel et al. 2008. *Communicating avoiding parallel and sequential QR*.

Fully serial TSQR

$$A = \begin{bmatrix} A_1 \\ A_2 \\ A_3 \\ A_4 \end{bmatrix}$$

Compute QR of A_1 ,
read A_2 , update QR, ...

$$A_1 = Q_1 R_1; \begin{bmatrix} R_1 \\ A_2 \end{bmatrix} = Q_2 R_2; \begin{bmatrix} R_2 \\ A_3 \end{bmatrix} = Q_3 R_3; \begin{bmatrix} R_3 \\ A_4 \end{bmatrix} = Q_4 R_4$$

$$A = \underbrace{\begin{bmatrix} Q_1 & & & \\ & I_{2n} & & \\ & & I_{2n} & \\ & & & I_{2n} \end{bmatrix}}_{8n \times 7n} \underbrace{\begin{bmatrix} Q_2 & & \\ & I_{2n} & \\ & & I_{2n} \end{bmatrix}}_{7n \times 5n} \underbrace{\begin{bmatrix} Q_3 & \\ & I_{2n} \end{bmatrix}}_{5n \times 3n} \underbrace{Q_4}_{3n \times n} \underbrace{R}_{n \times n}$$

$=Q$

Demmel et al. 2008. Communicating avoiding parallel and sequential QR.

More generally, any sequence of QR factorizations in a tree that yield a single R is a valid QR factorization.

In our MapReduce implementation, we combine these two steps!

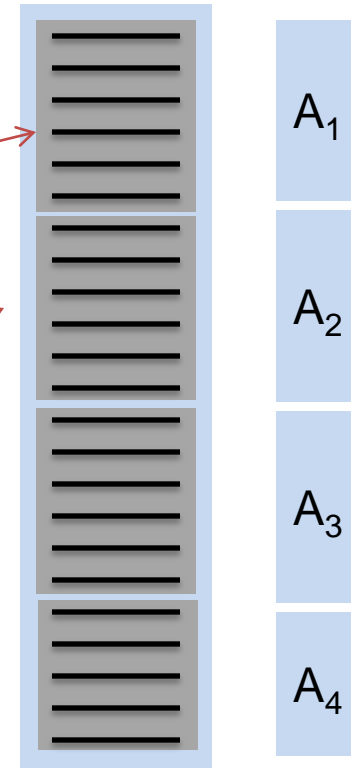
MapReduce matrix storage

$\mathbf{A} : m \times n, m \gg n$

Key is an arbitrary row-id

Value is the $1 \times n$ array for
a row.

Each submatrix \mathbf{A}_i is an
input split.



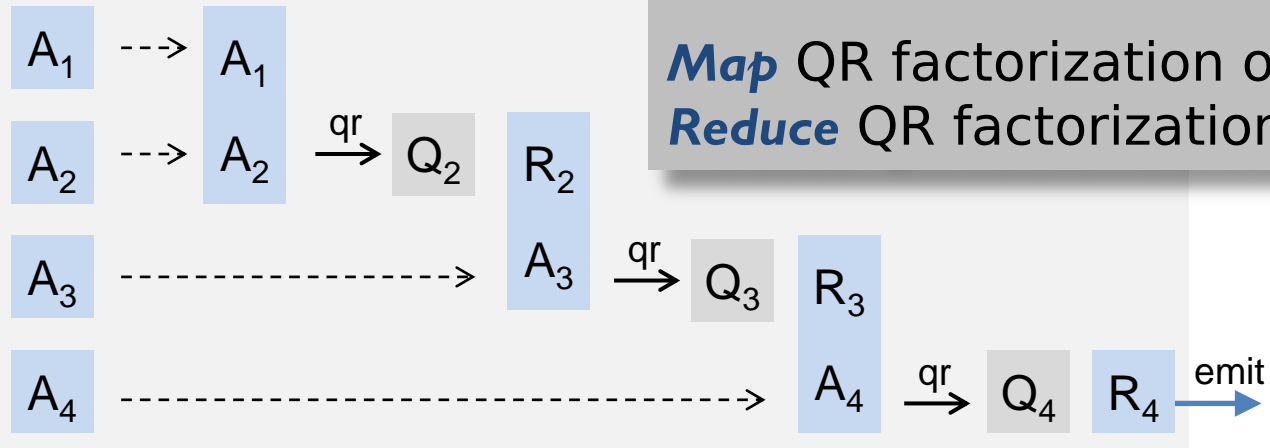
Algorithm

Data Rows of a matrix

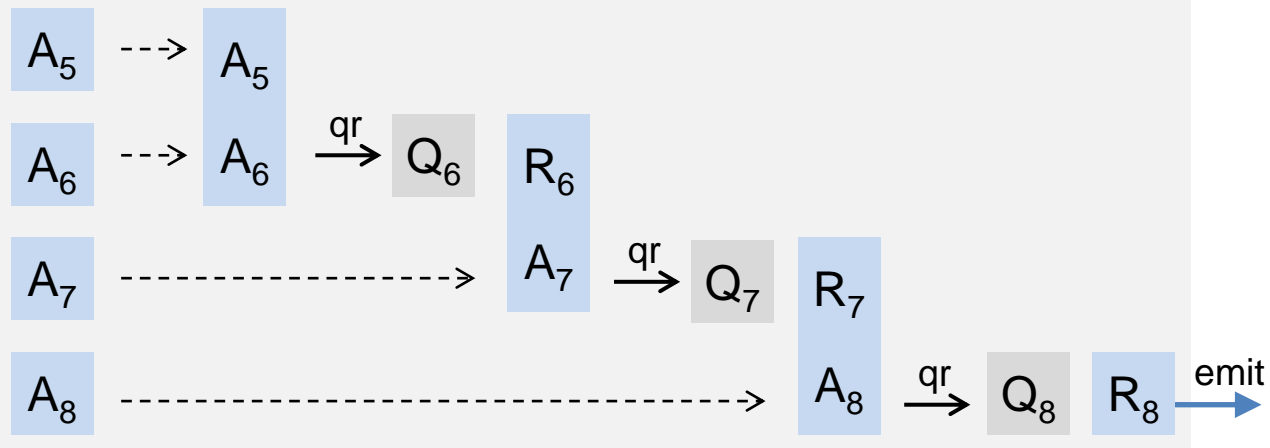
Map QR factorization of rows

Reduce QR factorization of rows

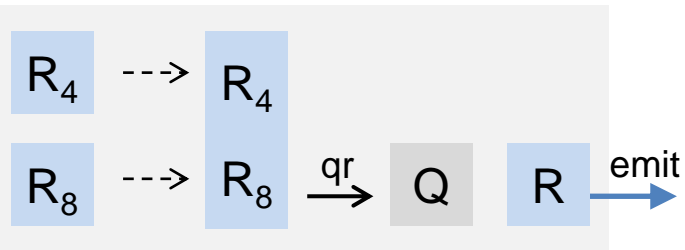
Mapper 1
Serial TSQR



Mapper 2
Serial TSQR



Reducer 1
Serial TSQR



Key Limitation

Computes only **R** and not **Q**

But **$Q = AR^+$** (pseudo-inverse) with another MapReduce iteration.

This approach has poor numerical stability, but we (in collaboration with Jim Demmel and Austin Benson) are working on it. (Multiple iterations do better)

We use random keys for new output – makes it hard to store data to recreate Q in a different fashion.

In hadoop

```
import random, numpy, hadoop

class SerialTSQR:
    def __init__(self,blocksize,isreducer):
        self.bsize=blocksize
        self.data = []
        if isreducer: self.__call__ = self.reducer
        else: self.__call__ = self.mapper

    def compress(self):
        R = numpy.linalg.qr(
            numpy.array(self.data),'r')
        # reset data and re-initialize to R
        self.data = []
        for row in R:
            self.data.append([float(v) for v in row])

    def collect(self,key,value):
        self.data.append(value)
        if len(self.data)>self.bsize*len(self.data[0]):
            self.compress()

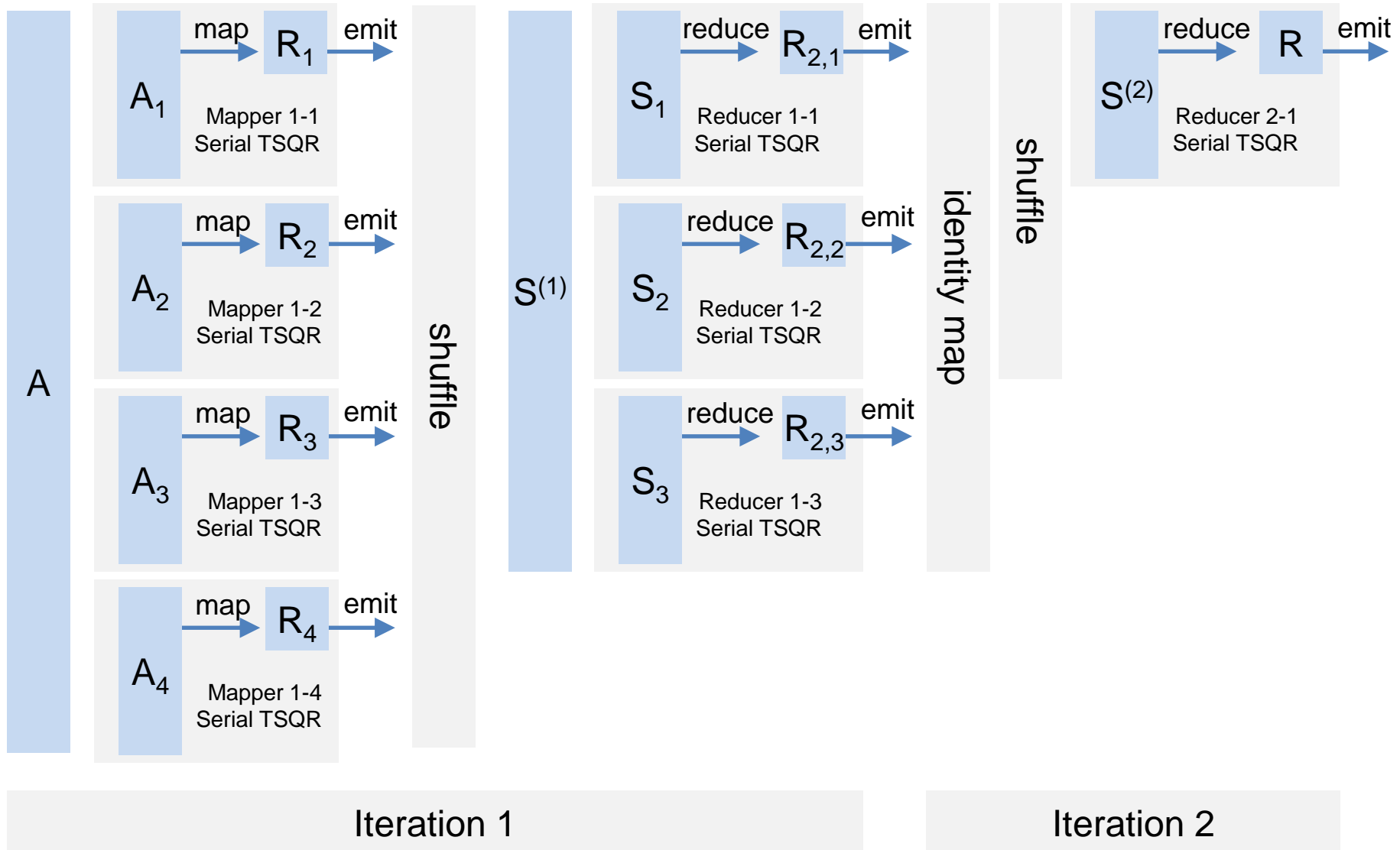
    def close(self):
        self.compress()
        for row in self.data:
            key = random.randint(0,2000000000)
            yield key, row

    def mapper(self,key,value):
        self.collect(key,value)

    def reducer(self,key,values):
        for value in values: self.mapper(key,value)

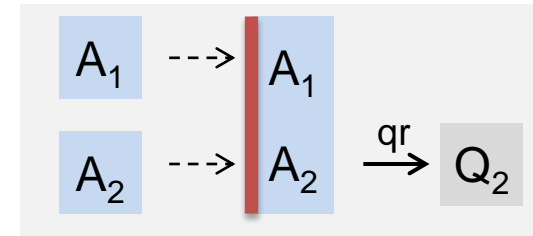
if __name__=='__main__':
    mapper = SerialTSQR(blocksize=3,isreducer=False)
    reducer = SerialTSQR(blocksize=3,isreducer=True)
    hadoop.run(mapper, reducer)
```

Too many maps? Add an iteration!

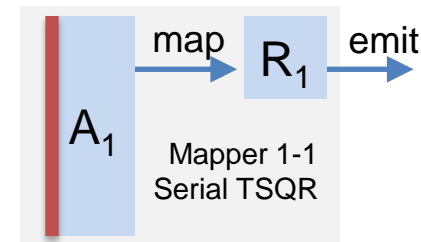


mrtsqr – summary of parameters

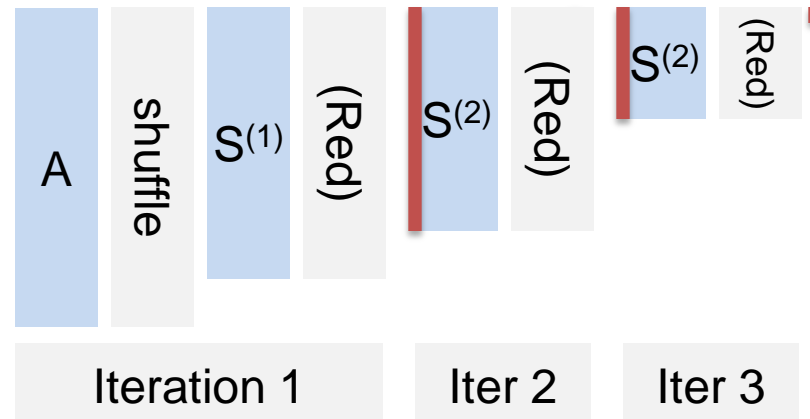
Blocksize How many rows to read before computing a QR factorization, expressed as a multiple of the number of columns (See paper)



Splitsize The size of each local matrix



Reduction tree
The number of reducers and iterations to use



Hadoop streaming frameworks

Synthetic data test 100,000,000-by-500 matrix (~500GB)

Codes implemented in MapReduce streaming

Matrix stored as TypedBytes lists of doubles

Python frameworks use Numpy+Atlas

Custom C++ TypedBytes reader/writer with Atlas

New non-streaming Java implementation too

	<i>Iter 1 QR (secs.)</i>	<i>Iter 1 Total (secs.)</i>	<i>Iter 2 Total (secs.)</i>	<i>Overall Total (secs.)</i>
Dumbo	67725	960	217	1177
Hadoopy	70909	612	118	730
C++	15809	350	37	387
Java		436	66	502

C++ in streaming beats a native Java implementation.

All timing results from the Hadoop job tracker

Varying split size Synthetic Data

<i>Cols.</i>	<i>Iters.</i>	<i>Split (MB)</i>	<i>Maps</i>	<i>Secs.</i>
50	1	64	8000	388
–	–	256	2000	184
–	–	512	1000	149
–	2	64	8000	425
–	–	256	2000	220
–	–	512	1000	191
1000	1	512	1000	666
–	2	64	6000	590
–	–	256	2000	432
–	–	512	1000	337

Increasing split size improves performance (accounts for Hadoop data movement)

Increasing iterations helps for problems with many columns.

(1000 columns with 64-MB split size overloaded the single reducer.)

Tinyimages PCA

1000 pixels

80,000,000 images

A

Zero mean rows

X

TSQR

R

SVD

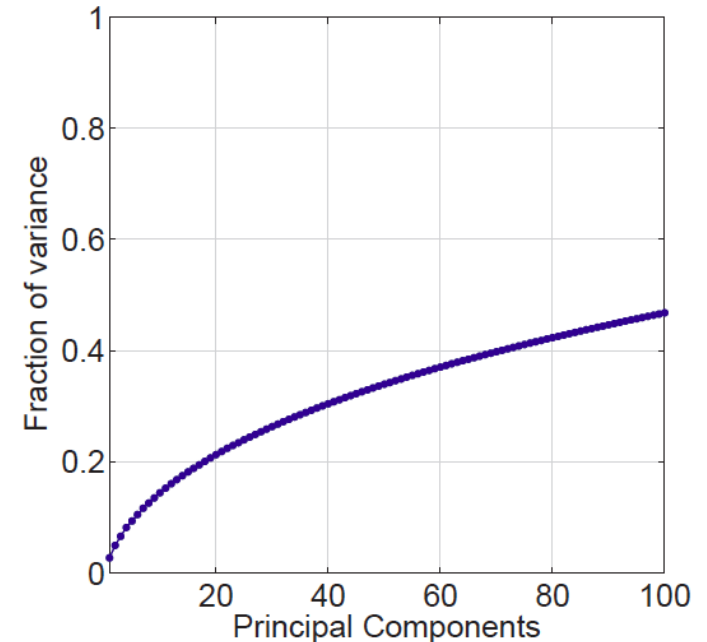
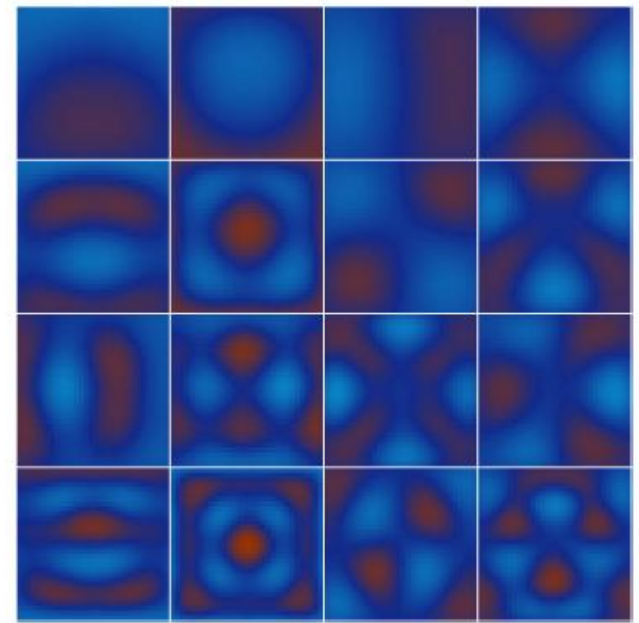
Σ

V

First 16 columns of V as images

(principal components)

Top 100 singular values



MapReduce

Post Processing

Tiny images is a 300GB database of images from 10,000 google queries. TSQR took about 30 minutes using dumbos framework

The problem Simulation ain't cheap!

21st Century Science in a nutshell

Experiments are not practical / feasible.

Simulate things instead.

But do we trust the simulations?

We're trying!

21st Century Science in a nutshell?

Simulations are expensive.

Let data provide a surrogate.

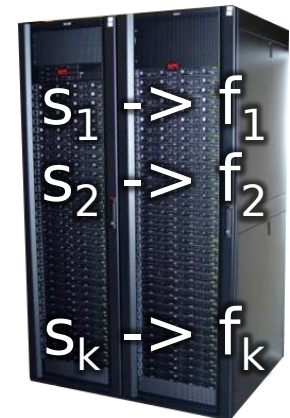
Input
Parameters



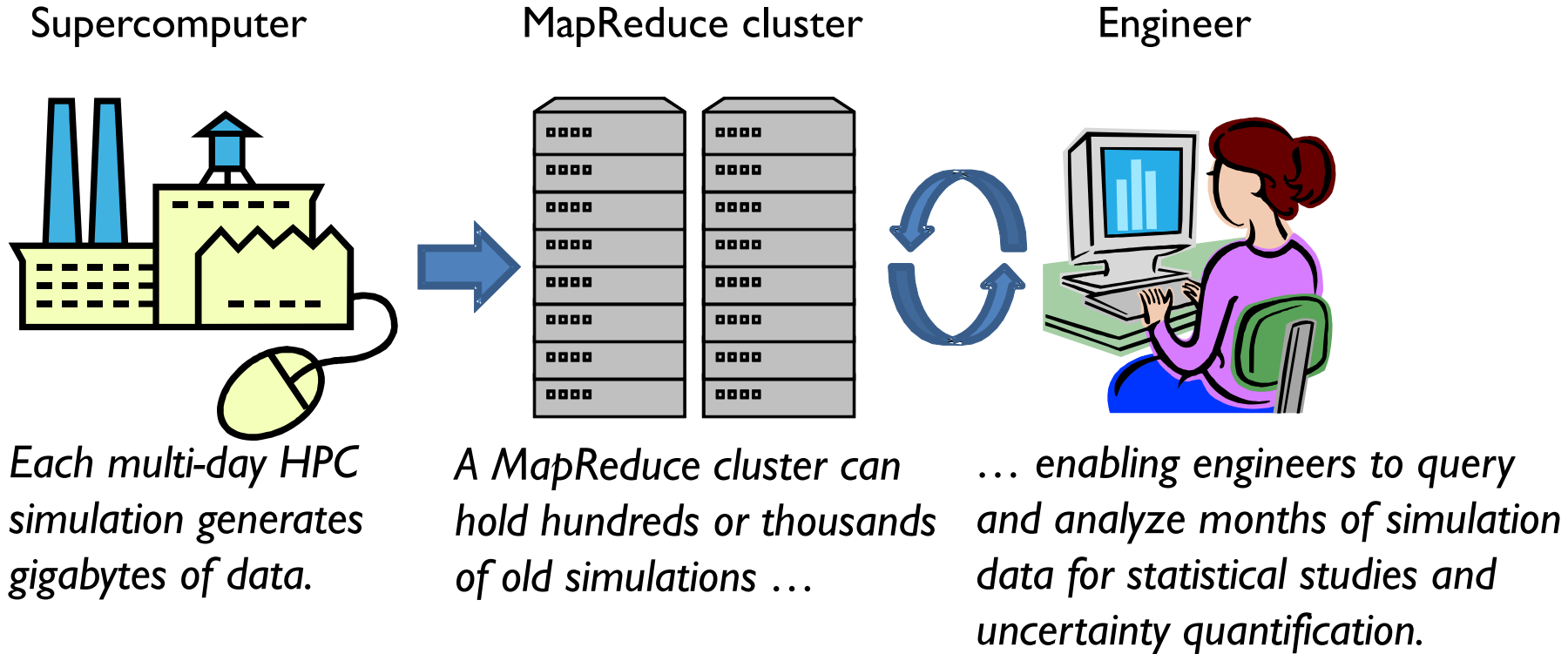
Time history
of simulation



The Database

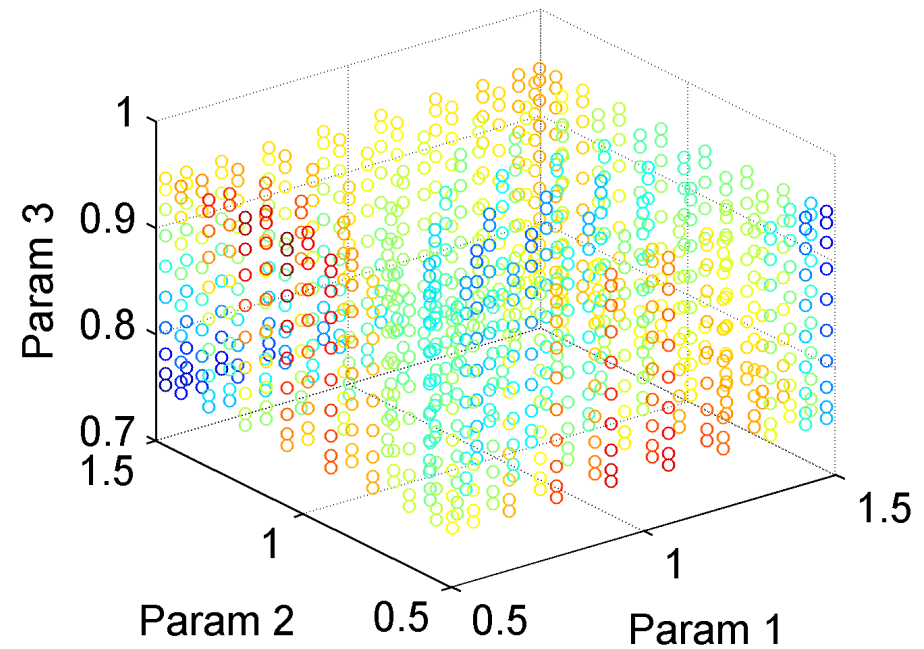
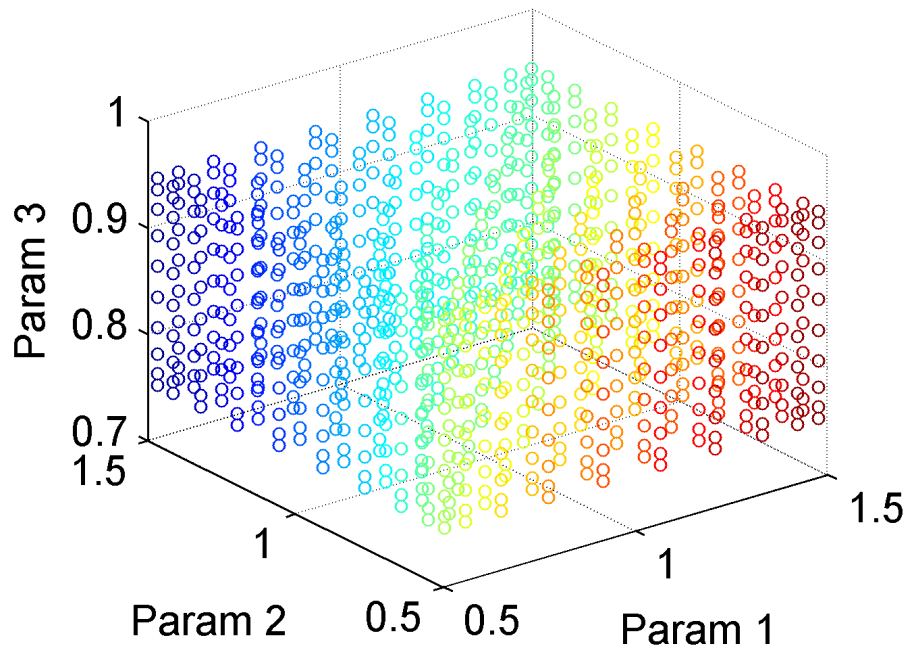


The idea and vision: Store the runs!



Data driven Green's functions

500GB of simulation data → TSSVD on 6,400,000 x 1000 matrix



The simulation varies most in the corners of the parameter space.

These analyses will help understand which parameters matter, and where to continue running the simulations.

Future work & Desiderata

Round off analysis of MR job
for **Q**

Compare against MPI based
TSQR *What's a fair way?*

Compare against random
sampling

Non-random keys, better
partitioners?

A mathematical reduce

Serialization for numeric
arrays *Roll my own?*

Google mrtsqr gleich

```
git clone https://github.com/dgleich/mrtsqr
```

Why not *the normal equations*?

Given **A**

compute **B** = **A**^T**A**

then **B** = **R**^T**R** with a Cholesky factorization

And it's the same **R** as in the QR!

But, you only get $O(\sqrt{\epsilon})$ accuracy vs. $O(\epsilon)$ accuracy for QR.
(ϵ is the machine precision)

Tinyimages Regression

We first solved a regression problem by trying to predict the sum of red-pixel values in each image as a linear combination of the gray values in each image. Formally, if r_i is the sum of the red components in all pixels of image i , and $G_{i,j}$ is the gray value of the j th pixel in image i , then we wanted to find $\min \sum_i (r_i - \sum_j G_{i,j} s_j)^2$. There is no particular importance to this regression problem, we use it merely as a demonstration.

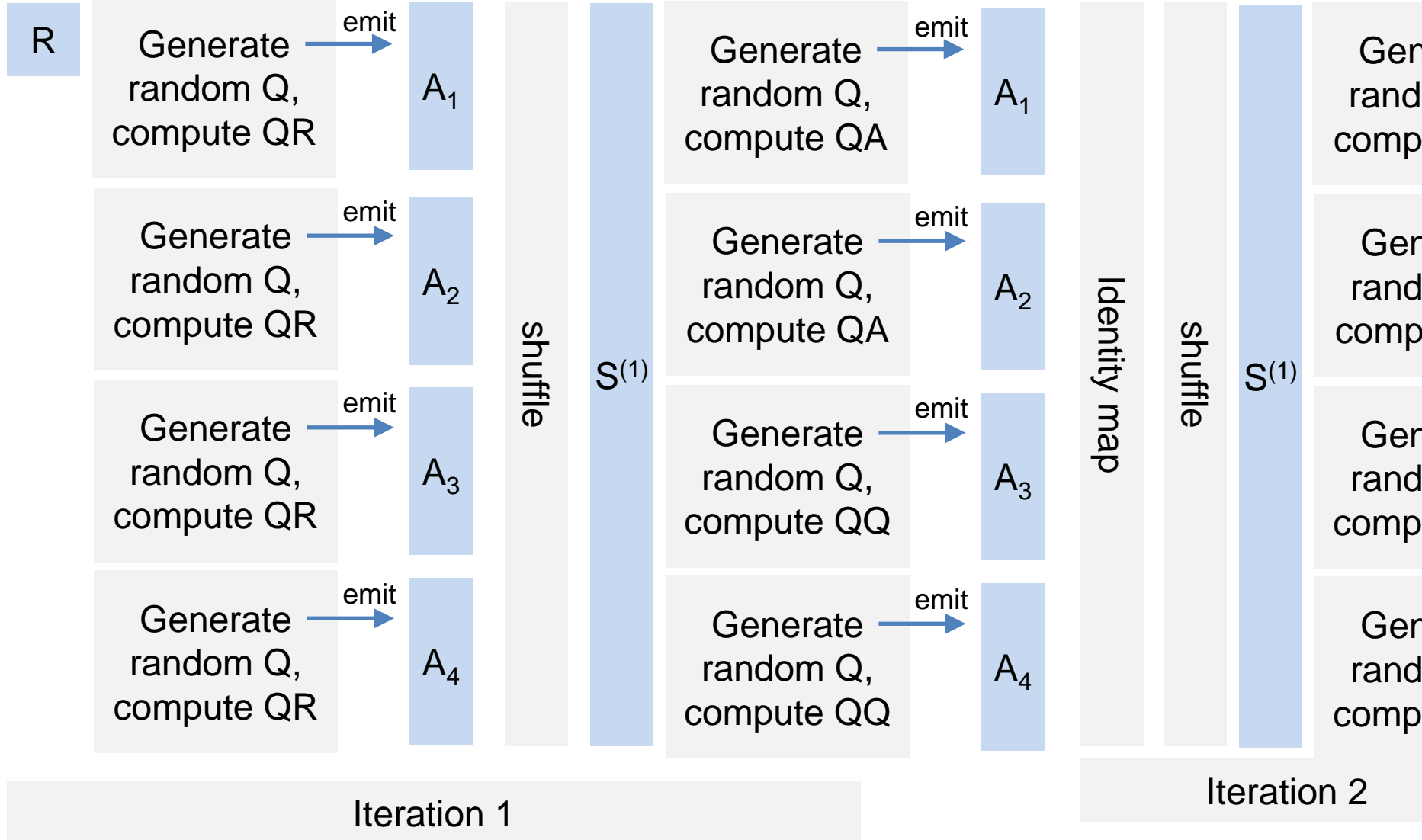


Took about 30 minutes

Synthetic Problems

Synthetic

Fixed



Varying Blocksize Synthetic Data

Cols.	Blks.	Iter. 1 Maps	Secs.	Iter. 2 Secs.
50	2	8000	424	21
—	3	—	399	19
—	5	—	408	19
—	10	—	401	19
—	20	—	396	20
—	50	—	406	18
—	100	—	380	19
—	200	—	395	19
100	2	7000	410	21
—	3	—	384	21
—	5	—	390	22
—	10	—	372	22
—	20	—	374	22
1000	2	6000	493	199
—	3	—	432	169
—	5	—	422	154
—	10	—	430	202
—	20	—	434	202

Using the C++ framework

Varying *splitsize* Synthetic Data

Cols.	Iters.	Split (MB)	Maps	Secs.
50	1	64	8000	388
—	—	256	2000	184
—	—	512	1000	149
—	2	64	8000	425
—	—	256	2000	220
—	—	512	1000	191
1000	1	512	1000	666
—	2	64	6000	590
—	—	256	2000	432
—	—	512	1000	337

Using the C++ framework

Comparing Frameworks *Synthetic Data*

	Iter 1				Total
	Map		Red.		
	Secs.	QR (s.)	Secs.	QR (s.)	
Dumbo	911	67725	884	2160	960
hadoopy	581	70909	565	2263	612
C++	326	15809	328	485	350
Java-MTJ					436
	Iter 2				Overall
	Map	Red.		Total	
	Secs.	Secs.	QR (s.)	Secs.	
Dumbo	5	214	80	217	1177
hadoopy	5	112	81	118	730
C++	5	34	15	37	387
Java-MTJ				66	502

All values from Hadoop Job tracker. 100,000,000-by-500 matrix.