

LA-UR-11-*1100221*

*Approved for public release;  
distribution is unlimited.*

<i>Title:</i>	Graph Anomalies in Cyber Communication
<i>Author(s):</i>	Scott A. Vander Wiel Curtis B. Storlie Gary Sandine Aric Hagberg Michael Fisk
<i>Intended for:</i>	2011 INFORMS Computing Society Conference, Monterey, California



Los Alamos National Laboratory, an affirmative action/equal opportunity employer, is operated by the Los Alamos National Security, LLC for the National Nuclear Security Administration of the U.S. Department of Energy under contract DE-AC52-06NA25396. By acceptance of this article, the publisher recognizes that the U.S. Government retains a nonexclusive, royalty-free license to publish or reproduce the published form of this contribution, or to allow others to do so, for U.S. Government purposes. Los Alamos National Laboratory requests that the publisher identify this article as work performed under the auspices of the U.S. Department of Energy. Los Alamos National Laboratory strongly supports academic freedom and a researcher's right to publish; as an institution, however, the Laboratory does not endorse the viewpoint of a publication or guarantee its technical correctness.

## ABSTRACT

### **Graph Anomalies in Cyber Communication**

Scott A. Vander Wiel

Curtis B. Storlie

Gary Sandine

Aric Hagberg

Michael Fisk

2011 INFORMS Computing Society Conference,  
Monterey, California

Enterprises monitor cyber traffic for viruses, intruders and stolen information. Detection methods look for known signatures of malicious traffic or search for anomalies with respect to a nominal reference model. Traditional anomaly detection focuses on aggregate traffic at central nodes or on user-level monitoring. More recently, however, traffic is being viewed more holistically as a dynamic communication graph. Attention to the graph nature of the traffic has expanded the types of anomalies that are being sought. We give an overview of several cyber data streams collected at Los Alamos National Laboratory and discuss current work in modeling the graph dynamics of traffic over the network. We consider global properties and local properties within the communication graph. A method for monitoring relative entropy on multiple correlated properties is discussed in detail.

## Outline

---

# Graph Anomalies in Cyber Communication

Scott Vander Wiel<sup>a</sup>, Curtis Storlie<sup>a</sup>,  
Gary Sandine<sup>b</sup>, Aric Hagberg<sup>b</sup>  
and Michael Fisk<sup>c</sup>

<sup>a</sup> Statistical Sciences Group  
<sup>b</sup> Applied Mathematics and Plasma Physics  
<sup>c</sup> Information Sciences Group

Los Alamos National Laboratory

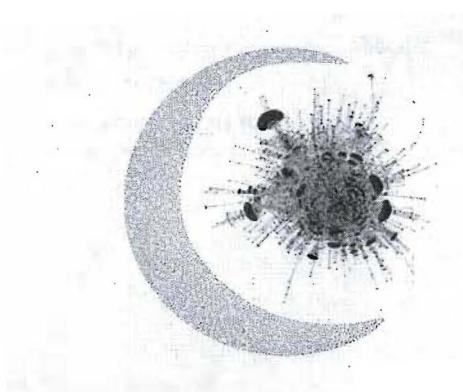
January 10, 2011

Motivation

Temporal Communication Graphs

Modeling Dynamic Properties

Detecting Anomalies



January 27, 2010  
**ExxonMobil, ConocoPhillips and Marathon Oil**  
executives were unwittingly duped by unsolicited e-mails carrying data-extracting malware ... This new form of corporate and, quite possibly, nation-sponsored espionage utilized custom spyware that is virtually undetectable by antivirus software applications used by the vast majority of large companies around the globe.  
[internetnews.com](http://internetnews.com)



January 15, 2010  
**Juniper Networks and Symantec** ... are investigating a widespread cyber-espionage incident that has hit dozens of technology companies, including **Google** and **Adobe**.

Sources familiar with the situation say that 34 companies, most of them large Fortune 500 names, were hit by a sophisticated cyber attack, first uncovered by Google last month. The attackers used a previously unknown "zero-day" attack on Internet Explorer, and possibly other techniques, to break into company networks and steal sensitive information.

InfoWorld

2009  
We regularly face attempts by others to gain unauthorized access ... by masquerading as authorized users or surreptitious introduction of software. These attempts ... are sometimes successful. ... We seek to detect and investigate these security incidents and to prevent their recurrence, but in some cases we might be unaware of an incident or its magnitude and effects.

*Intel report to SEC*

3/27

## Attack Detection and Response

**Cyber Attack:** misuse of a networked system such as

- penetration (*intrusion, exploitation*),
- remote command and control
- exfiltration of data
- denial of availability or integrity

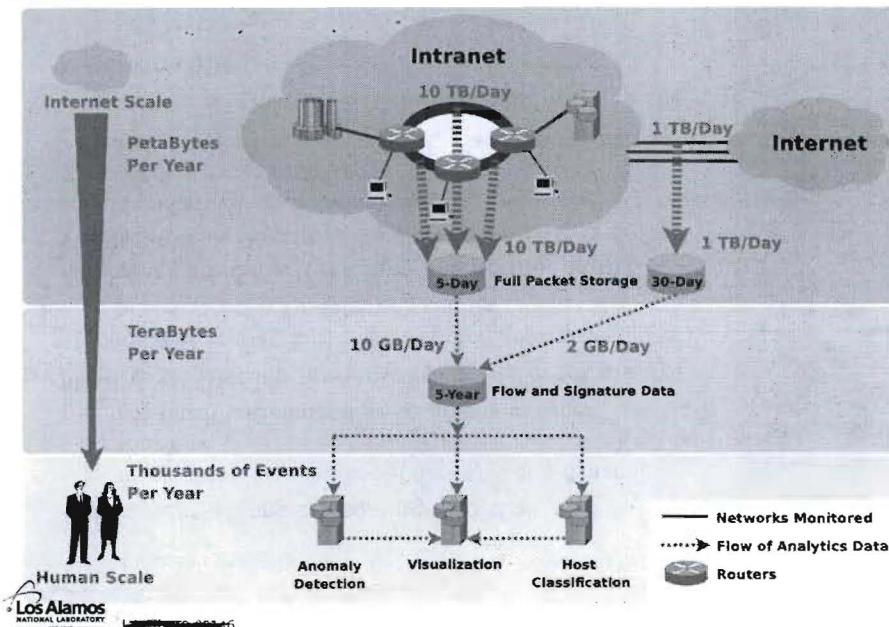
**Goals:** use observed sensor data to

- detect known attack methods and tools
- **detect unexplained patterns that could be attacks**
- prioritize responses based on likelihood of being an attack

New efforts to detect today's subtle cyber adversaries

- focus on invariants of attacker objectives rather than artifacts of specific attack technologies
- structural detection rather than simple rate-based detection

## Scale of Data Collection at Los Alamos



5/27

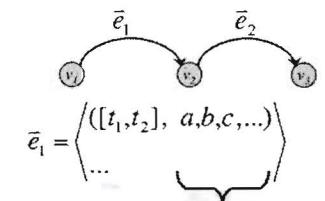
## Temporal Communication Graphs

Many cyber data sets can (and should) be described as graphs

- vertices are hosts, users, etc.
- directed edges are communications
- events from heterogeneous sensors can be combined in one graph

A graph construction supports traditional local analysis while enabling new analysis

- exploitation is often a temporal path through the network
- long paths alter structural characteristics of the graph



### Attributes:

- protocol: `tcp`, `udp`, ...
- packet count, bytes, ...
- type: `chat`, `web`, `ssh`, ...
- probability of edge
- user type: `scientist`, `admin`, ...



6/27

## Properties of Cyber Graphs

Cyber event graphs are coming into their own

- much graph work on social networks and dynamic topology
- not much on temporal event graphs in cyber

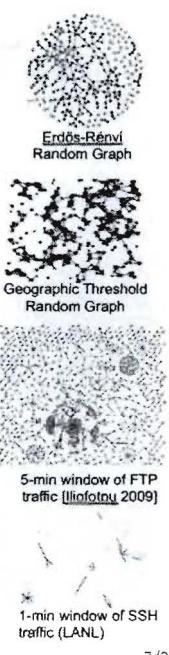
Need to understand properties of cyber event graphs

- improve relevance of random graph generative models
- provide a basis for detecting change

Anomalies are departures from the normal diurnal evolution of

- local activity on a node or edge
- distribution of subgraphs
- global properties (components, diameter, density, size)

Hypothesis: a nosy intruder cannot avoid altering the graph structure



7/27

## Subgraph Statistics

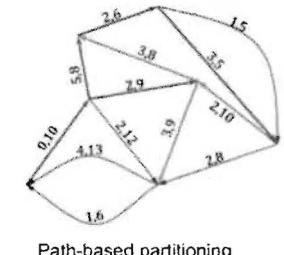
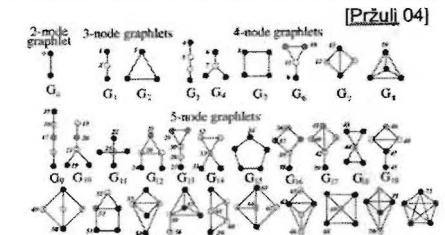
Enumerate a set of subgraphs.  
Some options:

- connected components
- graphlets
- temporal walks (paths)
  - telescoping
  - consecutive
  - overlapping

Measure subgraph features

- order
- diameter
- density

The distribution of features over subgraphs is a statistic of the full graph and evolves over time.



Path-based partitioning

8/27

## SSH Graph from LANL's Network



9/27

## Netflow-Type Data

Netflow is one type of data sampled from LANL's internal network by the routers. A Netflow record contains the following.

1. Source IP address
2. Source port
3. Destination IP address
4. Destination port
5. IP protocol
6. Ingress interface ID
7. IP Type of Service

In this talk we focus on SSH sessions (port=22) and build a graph for each ten-minute time interval in the month of November using the edges with connection times in the given interval

## Summaries of SSH Graphs

Summarize each *connected component* by

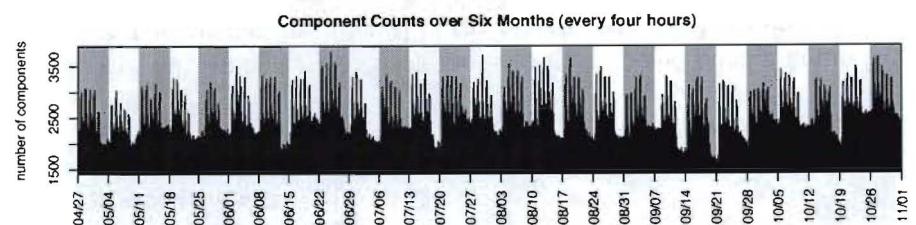
order:  $O = \text{number of nodes}$

diameter:  $D = \text{greatest distance between any pair of vertices}$

Each ten-minute bin of time provides a set of  $(O, D)$  pairs from disconnected components

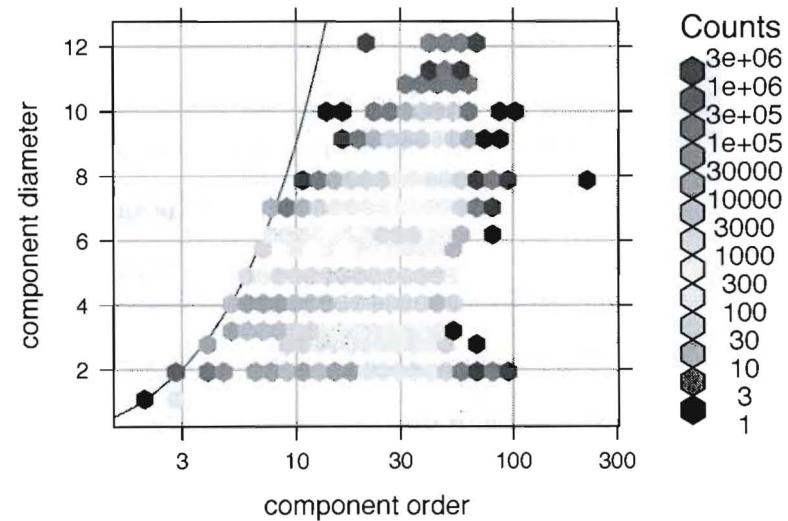
Future work: summarize using *graphlets* (arbitrary subsets of nodes and all edges between them)

## Patterns over Time

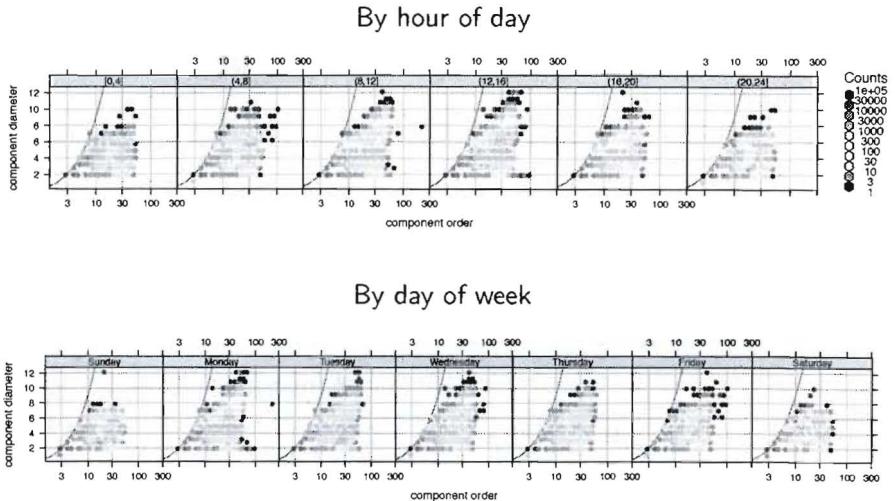


Numbers of components follow predictable daily and weekly patterns.

(O,D) Distribution Aggregated to a Full Month



Daily and weekly trends in (O,D)



## A Dynamic Model for Multivariate Count Data

$$\mathbf{x} = (\mathbf{O}, \mathbf{D})$$

$\mathcal{X}$  = set of possible  $(\mathbf{O}, \mathbf{D})$  pairs

$y_t(\mathbf{x})$  = count of subgraphs with  $(\mathbf{O}, \mathbf{D})$  equal  $\mathbf{x}$

Observe  $\{y_t(\mathbf{x}) : \mathbf{x} \in \mathcal{X}\}$  over a sequence of times  $t = 1, 2, \dots$

Goal: monitor counts over time and alarm if observed counts depart from established patterns.

Model the counts at time  $t$  as

$$\begin{aligned} \{y_t(\mathbf{x}) : \mathbf{x} \in \mathcal{X}\} | M_t &\sim \text{Multinomial}(M_t, \{p_t(\mathbf{x}) : \mathbf{x} \in \mathcal{X}\}), \\ M_t &\sim \text{Poisson}(\lambda_t). \end{aligned}$$

$M_t$  is the total count over the table.

Predict  $\lambda_t$  and  $p_t(\mathbf{x})$  from data up to  $t - 1$ .



## Predict $\lambda_t$ by Kernel Filtering Previous Counts, $M_s$

$$\hat{\lambda}_{t|t-1} = \frac{\sum_{d=0}^{\infty} \sum_{s=1}^{t-1} K_{\lambda}(t, s, d, \mathbf{h}_{\lambda}) M_s}{\sum_{d=0}^{\infty} \sum_{s=1}^{t-1} K_{\lambda}(t, s, d, \mathbf{h}_{\lambda})},$$

The kernel function decomposes as

$$K_{\lambda}(t, s, d, \mathbf{h}_{\lambda}) = K_D(d, h_{\lambda, d}) K_B(t, s, d, h_{\lambda, b})$$

$K_D$  reaches back to data on previous days

$$K_D(d, h_{\lambda, d}) = (1 - h_{\lambda, d})^d$$

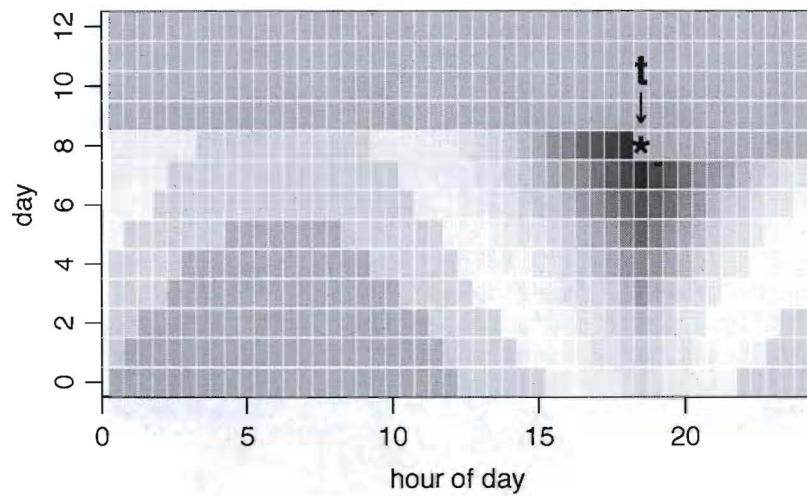
$K_B$  reaches back to data on previous time bins

$$K_B(t, s, d, h_{\lambda, b}) = (1 - h_{\lambda, b})^{|t - d * T - s|} \quad (T = 144 \text{ bins/day})$$

$\mathbf{h}_{\lambda} = (h_{\lambda, d}, h_{\lambda, b})$  are tuning parameters chosen by cross-validation.

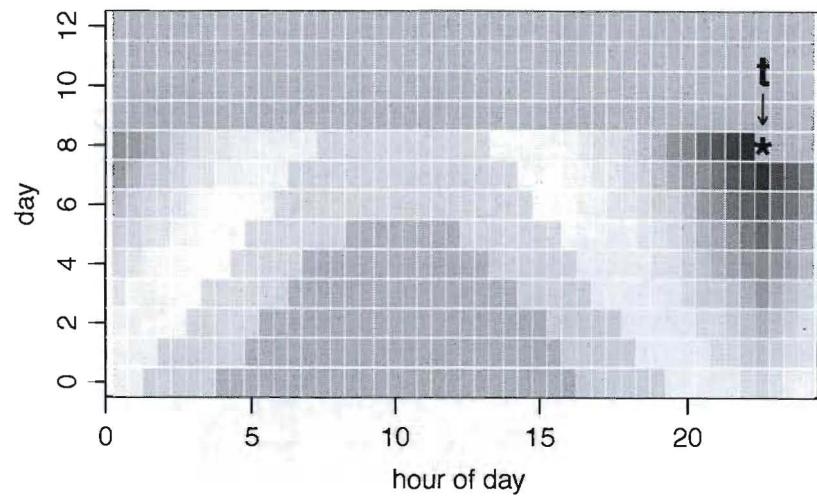


The Kernel for Predicting  $\lambda_t$  from  $M_1, \dots, M_{t-1}$



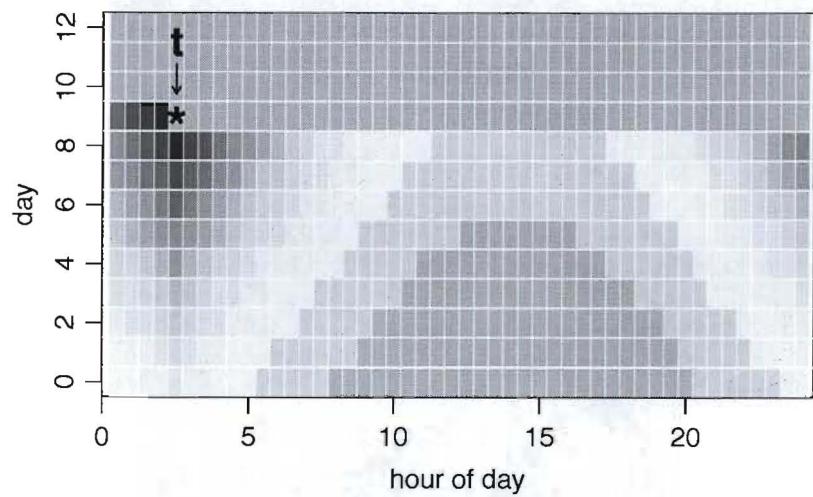
$\mathbf{h}_\lambda = (0.17, 0.13)$  — equivalent to 11 days  $\times$  14 half-hour bins

The Kernel for Predicting  $\lambda_t$  from  $M_1, \dots, M_{t-1}$



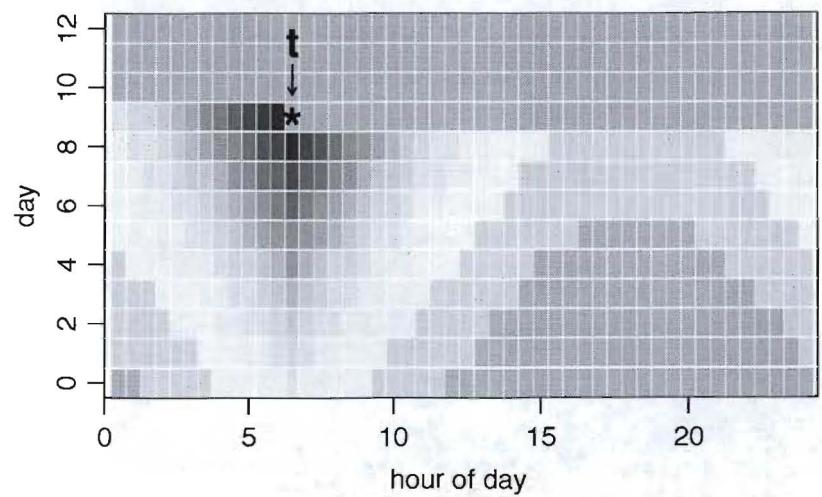
$\mathbf{h}_\lambda = (0.17, 0.13)$  — equivalent to 11 days  $\times$  14 half-hour bins

The Kernel for Predicting  $\lambda_t$  from  $M_1, \dots, M_{t-1}$



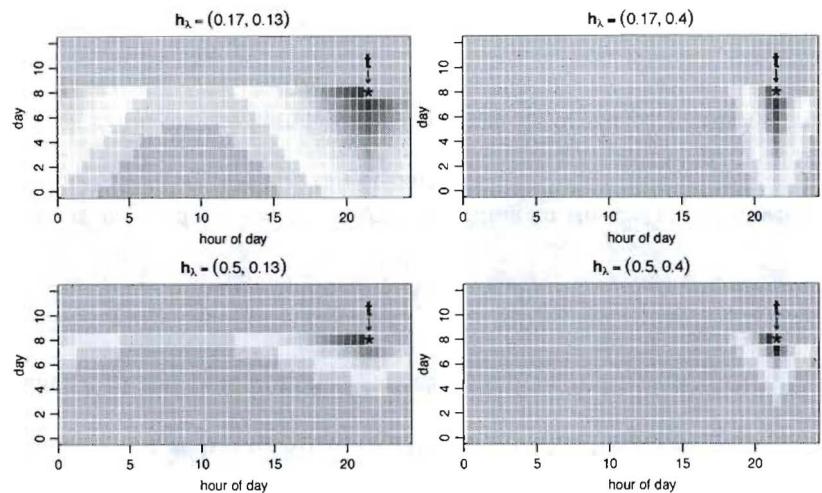
$\mathbf{h}_\lambda = (0.17, 0.13)$  — equivalent to 11 days  $\times$  14 half-hour bins

The Kernel for Predicting  $\lambda_t$  from  $M_1, \dots, M_{t-1}$



$\mathbf{h}_\lambda = (0.17, 0.13)$  — equivalent to 11 days  $\times$  14 half-hour bins

## Adjusting the Tuning Parameters



$h_{\lambda,d}$  adjusts height;  $h_{\lambda,b}$  adjusts width

## Predict $p_t(\mathbf{x})$ by Kernel Filtering

Filter previous empirical (O,D) distributions

$$\{\tilde{p}_s(\mathbf{x}) \equiv y_s(\mathbf{x})/M_s : \mathbf{x} \in \mathcal{X}\} \quad (s < t)$$

to predict  $p_t(\mathbf{x})$  as

$$\hat{p}_{t|t-1}(\mathbf{x}) = \frac{\sum_{d=0}^{\infty} \sum_{s=1}^{t-1} \sum_{\mathbf{w} \in \mathcal{X}} K_P(t, s, d, \mathbf{x}, \mathbf{w}, \mathbf{h}_P) \tilde{p}_s(\mathbf{w})}{\sum_{d=0}^{\infty} \sum_{s=1}^{t-1} \sum_{\mathbf{w} \in \mathcal{X}} K_P(t, s, d, \mathbf{x}, \mathbf{w}, \mathbf{h}_P)},$$

Choose tuning parameters  $\mathbf{h}_P = [h_{P,d}, h_{P,b}, h_{P,1}, \dots, h_{P,q}, h_{P,\delta}]$  by cross validation.

## Predict $p_t(\mathbf{x})$ by Kernel Filtering (cont.)

The kernel decomposes as

$$K_P(t, s, d, \mathbf{x}, \mathbf{w}, \mathbf{h}_P) = K_D(d, h_{P,d}) K_B(t, s, d, h_{P,b}) K_W(\mathbf{x}, \mathbf{w}, \mathbf{h}_P) M_t^{-1}$$

with day ( $K_D$ ) and time-bin ( $K_B$ ) kernels as before and now a *within* (O,D) kernel

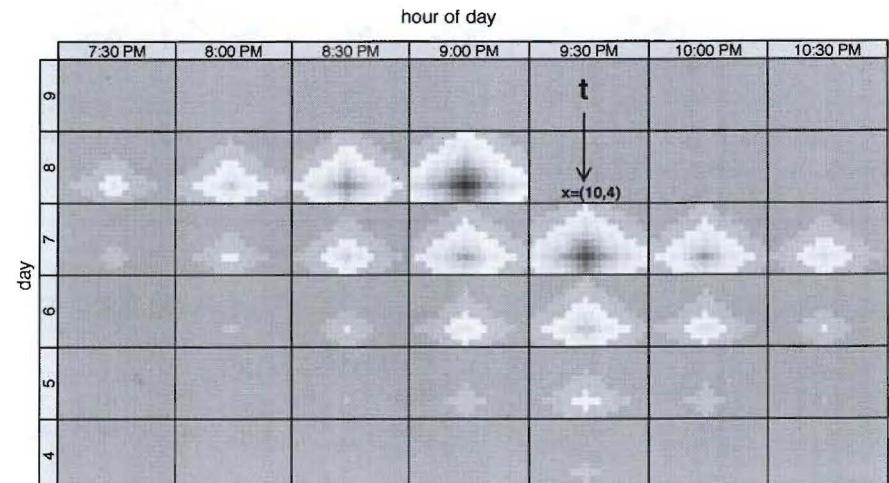
$$K_W(\mathbf{x}, \mathbf{w}, \mathbf{h}_P) = \prod_{j=1}^2 (1 - h_{P,j} h_{\delta}^{x_j})^{|x_j - w_j|}.$$

Parameters  $h_{P,j}$  reduce kernel weights according to the distance between a given (O,D) pair ( $\mathbf{w}$ ) and the target pair ( $\mathbf{x}$ )

Parameter  $h_{\delta}$  makes the kernel wider for larger  $x_j$ , where probabilities are likely smaller.

$M_t^{-1}$  weights by inverse variance.

## Kernel for Estimating $p_t(10, 3)$



## Anomaly Detection

At time  $t - 1$ , a predictive distribution for the next set of counts is

$$y_t(\mathbf{x}) \sim \text{Poisson}(\hat{\lambda}_{t|t-1} \hat{p}_{t|t-1}(\mathbf{x})).$$

These are independent over  $\mathbf{x} \in \mathcal{X}$  under the multinomial model.

We use exceedance probabilities to detect unusually large subgraph counts relative to predictions:

$$U_t(\mathbf{x}) \equiv \Pr(\text{Poisson} \geq y_t(\mathbf{x}))$$

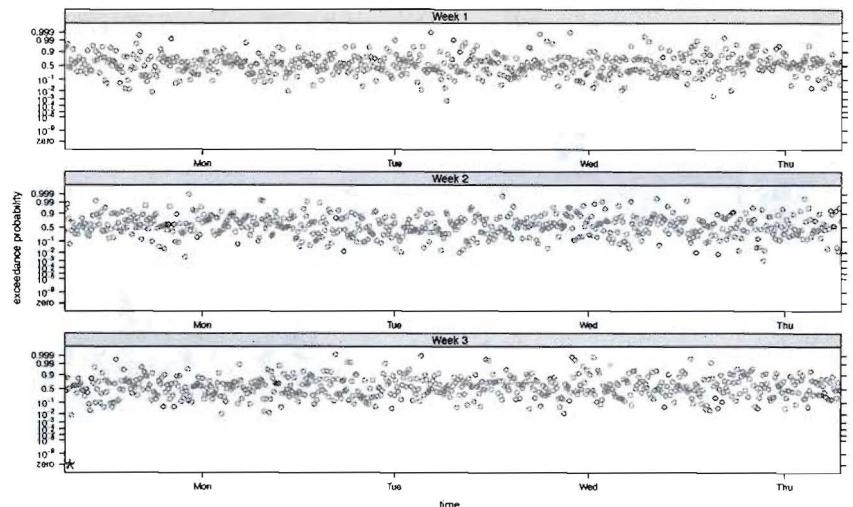
The joint exceedance probability

$$U_t = \prod_{\mathbf{x} \in \mathcal{X}} U_t(\mathbf{x})$$

has a known distribution (negative log-gamma if randomization is used), providing a reference to calculate p-values for joint exceedance.

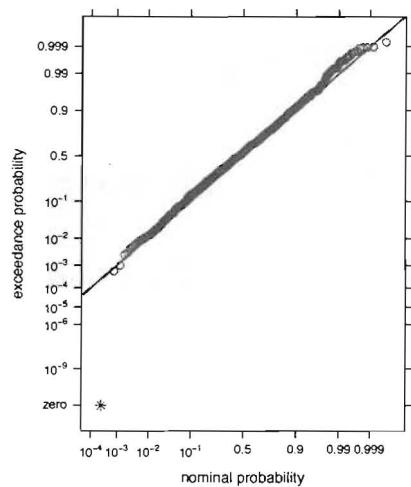
## P-Values for Joint Exceedance

Three weeks of monitoring SSH graphs every ten-minutes.



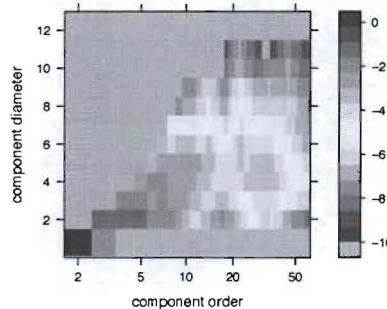
## Distribution of P-Values for Joint Exceedance

A QQ plot shows we have a reasonable reference model for monitoring.

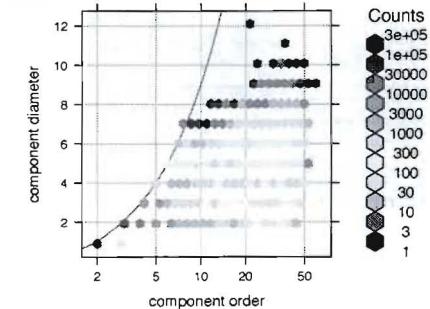


## The Anomaly: Compare $\hat{p}_t(x)$ to Data

$\log_{10}$  estimated probabilities



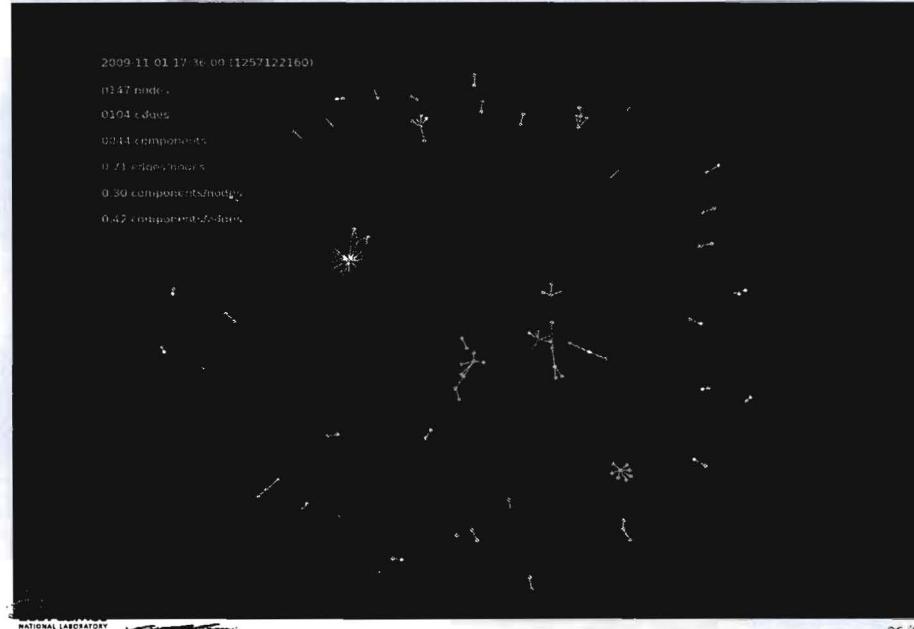
Histogram for 4 to 8 PM



One subgraph with  $(O,D) = (21, 12)$  is highly unusual at this time of day (but not at other times).

More broadly distributed anomalies can also be detected.

SSH graph from LANL's Network During the Anomaly



SSH graph from LANL's Network During the Anomaly



SSH graph from LANL's Network During the Anomaly



26/27

SSH graph from LANL's Network During the Anomaly



26/27

SSH graph from LANL's Network During the Anomaly



26/27

SSH graph from LANL's Network During the Anomaly



26/27

SSH graph from LANL's Network During the Anomaly



SSH graph from LANL's Network During the Anomaly

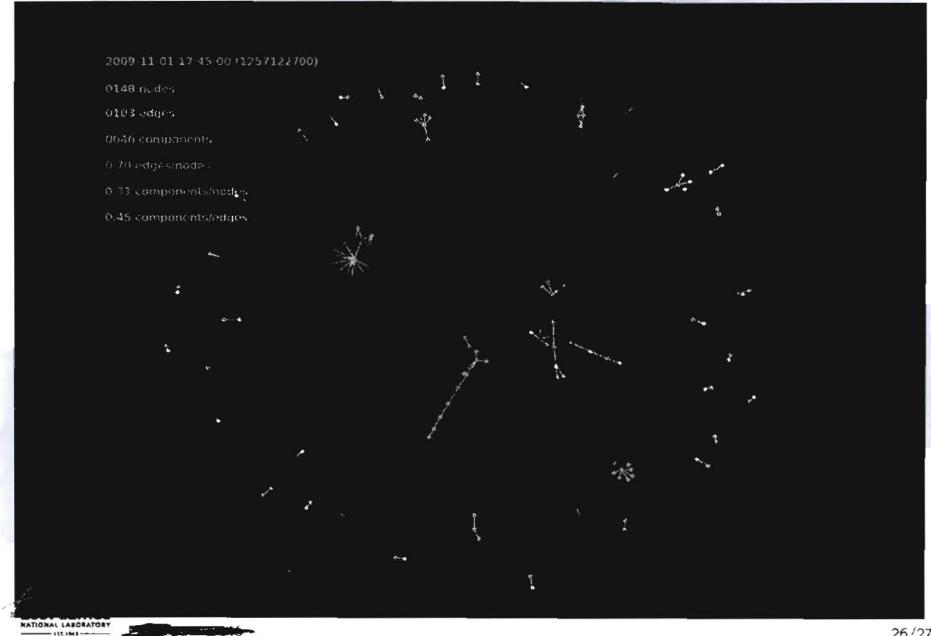


SSH graph from LANL's Network During the Anomaly



26/27

SSH graph from LANL's Network During the Anomaly



26/27

## SSH graph from LANL's Network During the Anomaly



## Conclusions & Further Work

---

- Monitoring subgraph properties can signal changes in network behavior.
- The dynamic kernel filter effectively smooths data from recent times in the target day from other recent days.
- The approach can be extended to incorporate
  - additional subgraph properties (eg., number of edges)
  - time constraints on paths (eg., telescoping times)
  - local communication patterns (eg., busy edges vs. quiet ones)
  - additional services (SSH is only illustrative)